



Katsushika Hokusai (葛飾 北斎)

Enoshima in Sagami Province (相州江の島)

*from Thirty-six views of Mount Fuji*

# An information theorist's tour of differential privacy

Anand D. Sarwate, Rutgers University

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# Some thanks and credits



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**Shahab Asoodeh (McMaster)**  
**Flavio Calmon (Harvard)**  
**Oliver Kosut (Arizona State)**  
**Lalitha Sankar (Arizona State)**  
**Mario Diaz (UNAM) - in memoriam**

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- Describe some of these three connections for those less familiar
- Suggest some questions for discussion later?



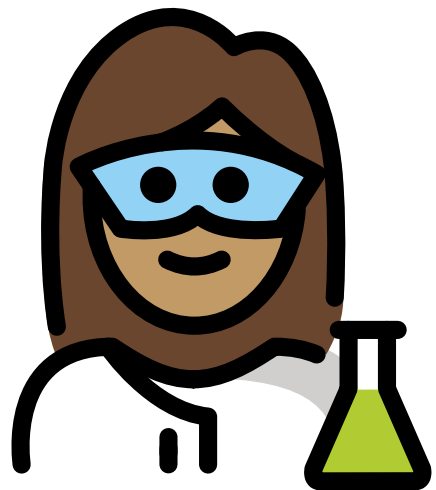
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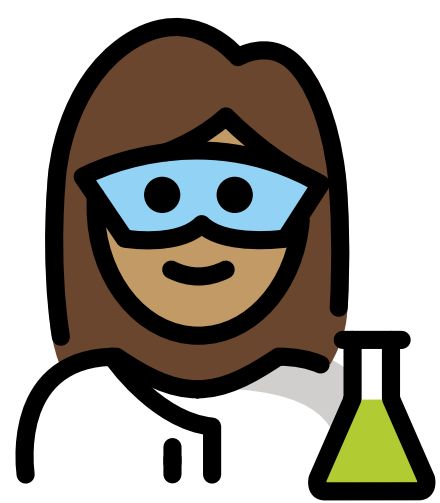
Sasha



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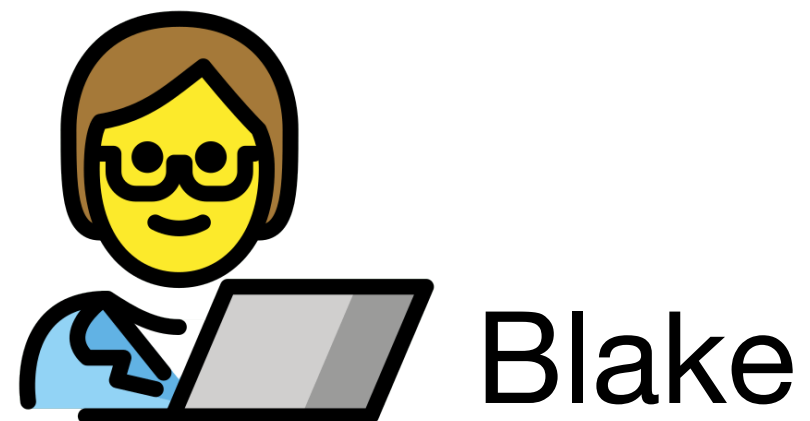
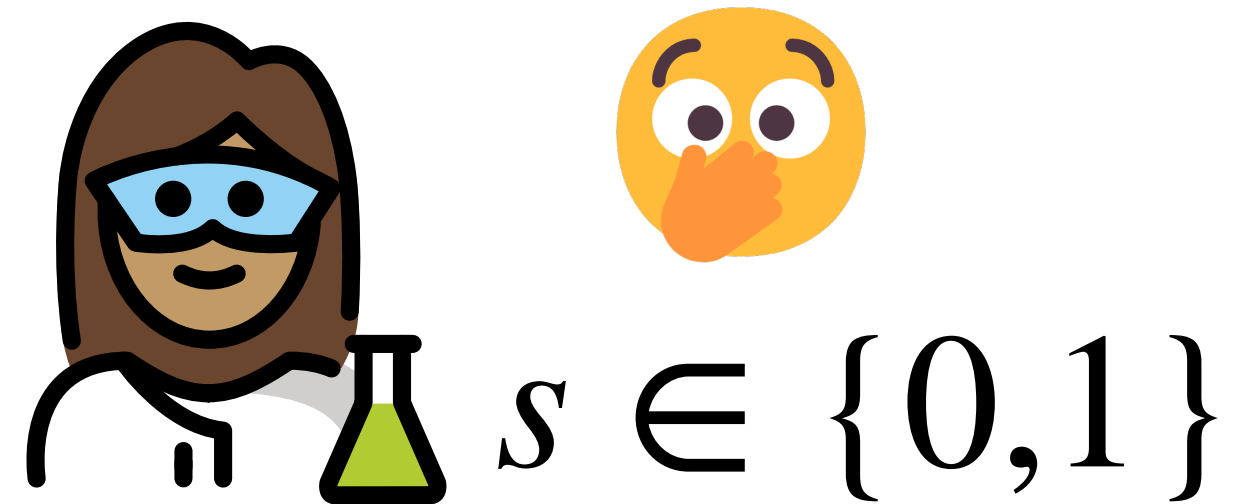
$s \in \{0,1\}$



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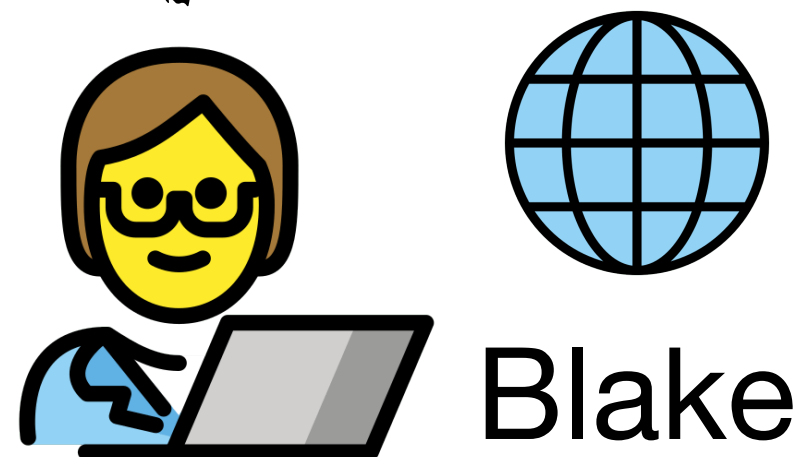
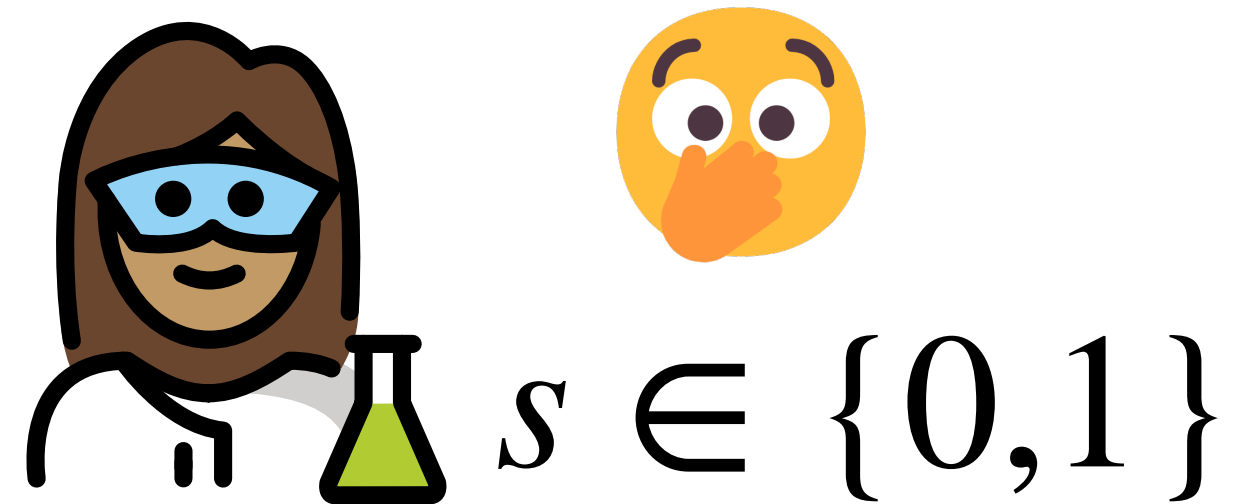
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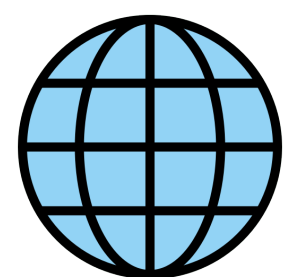
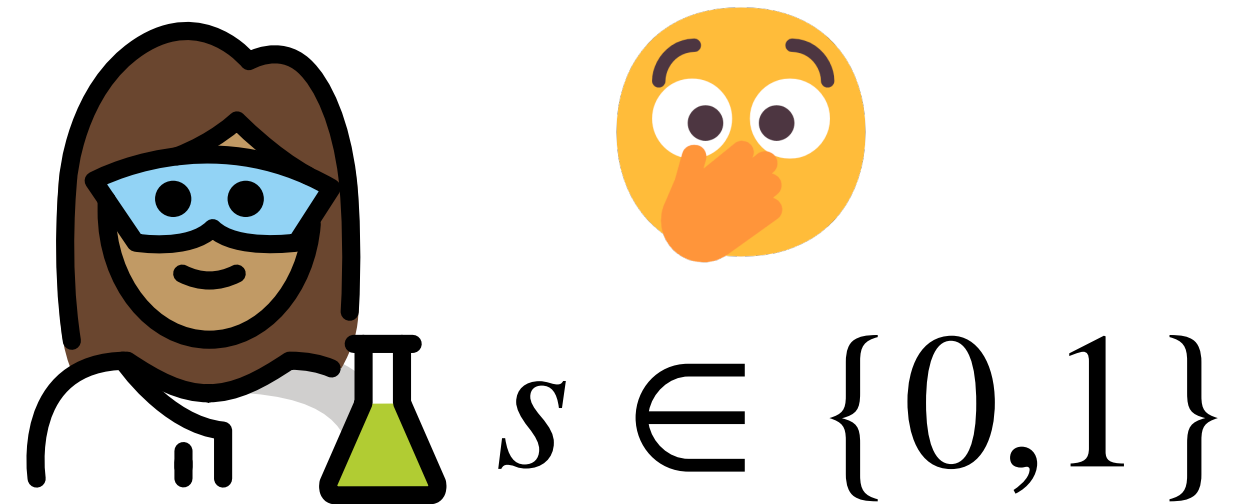
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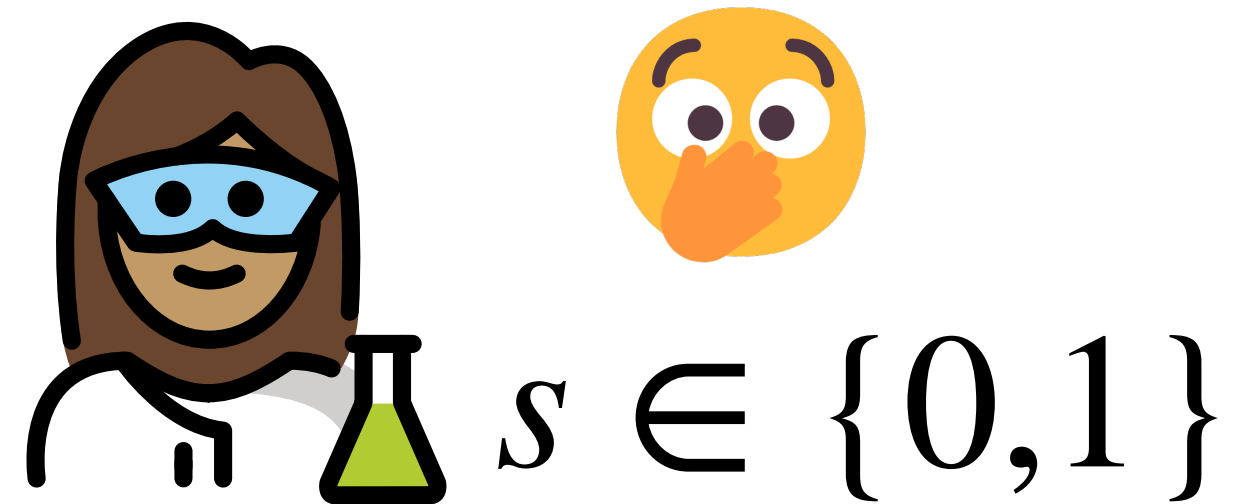
Blake



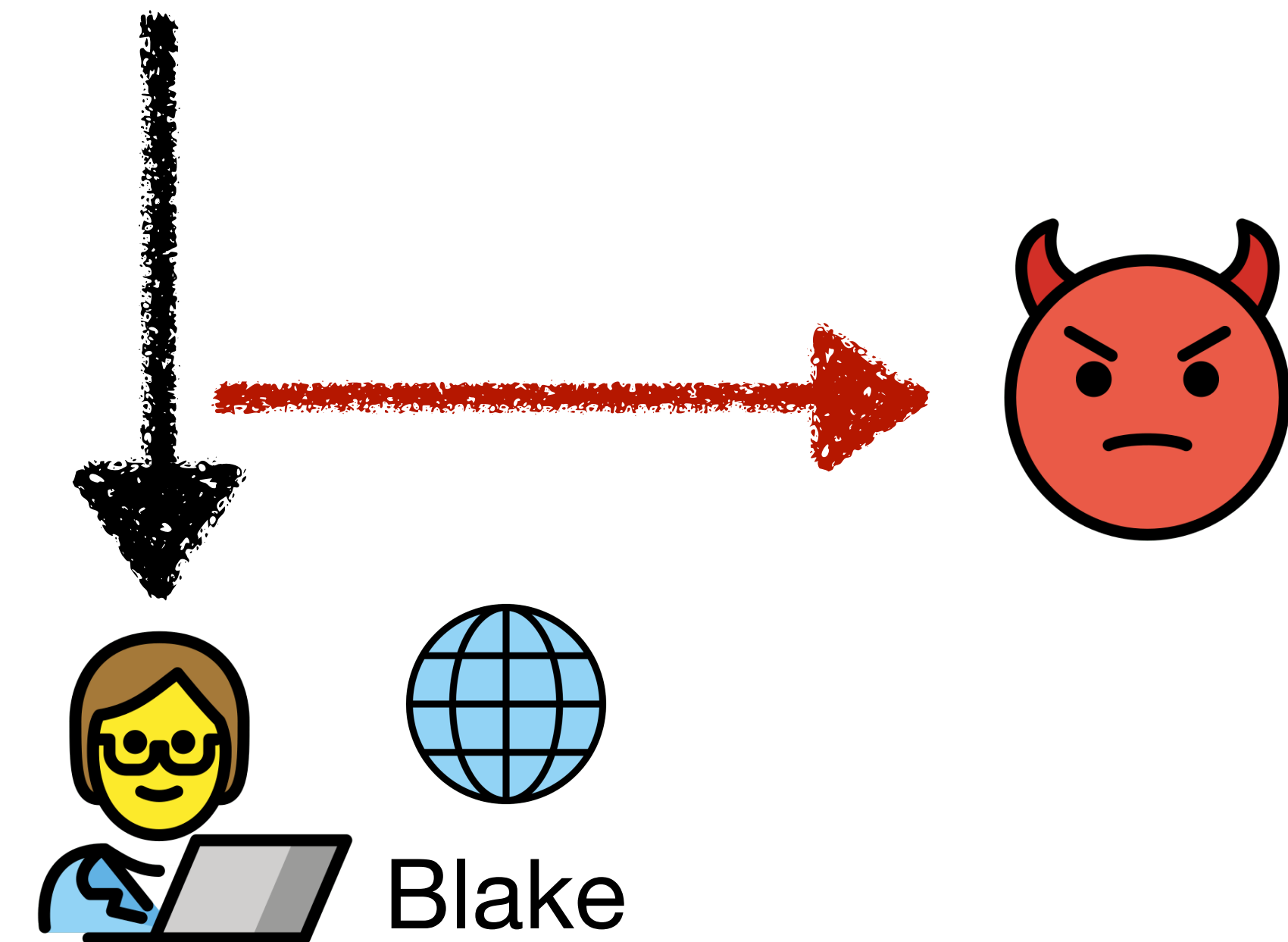
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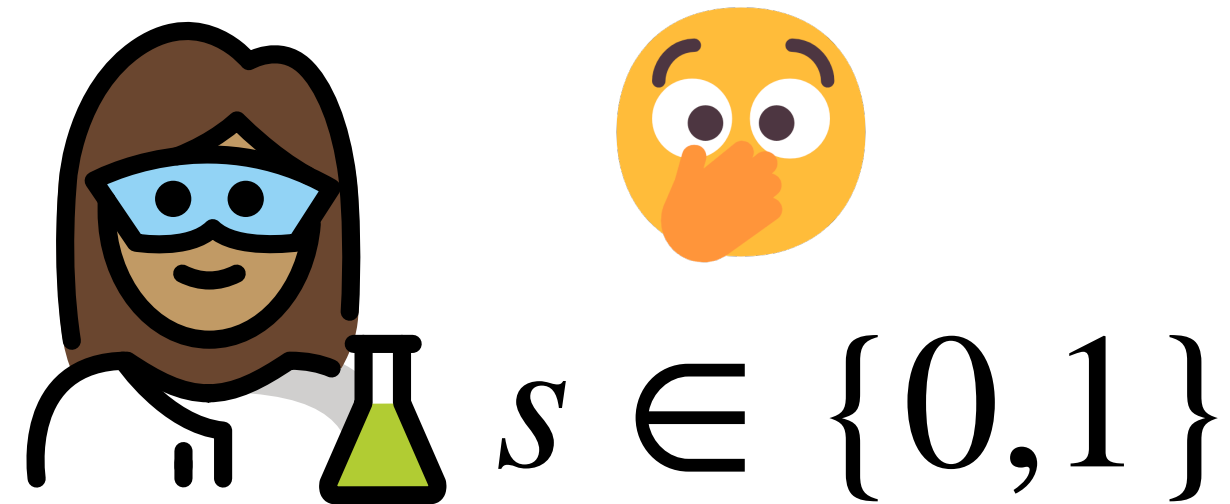
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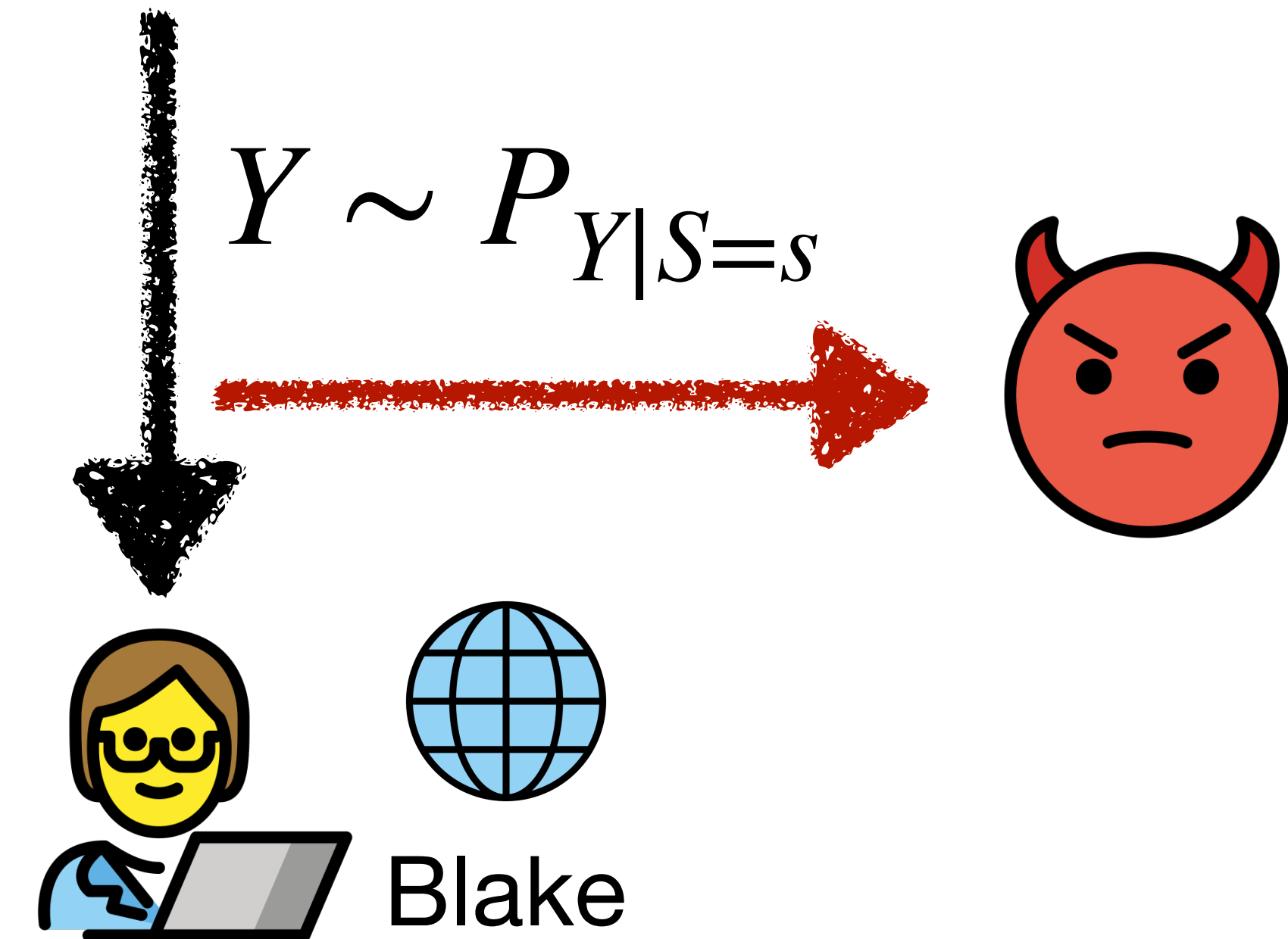
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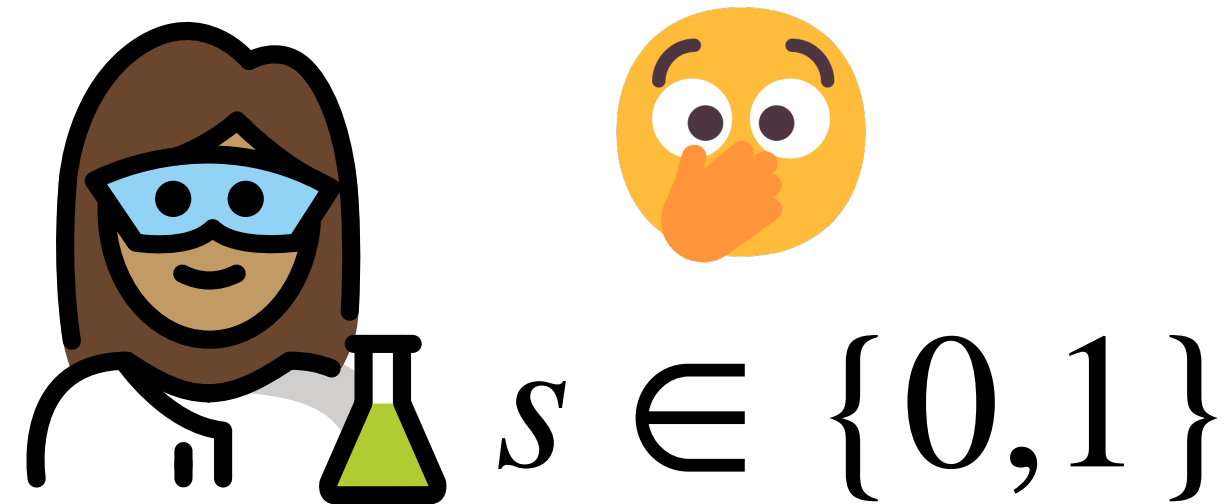


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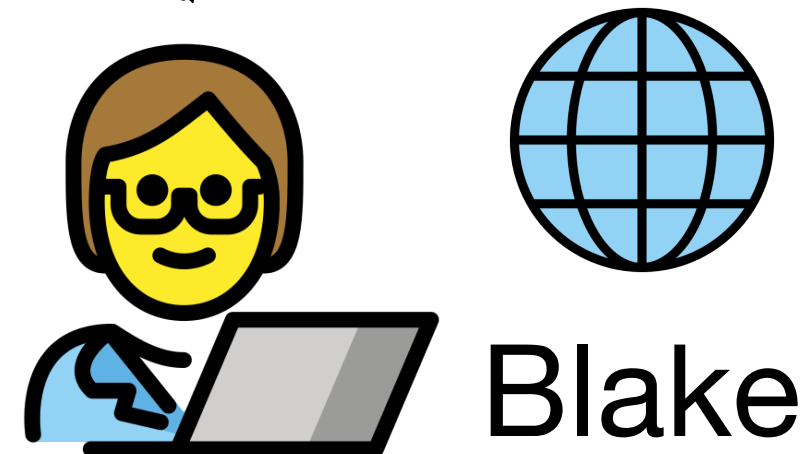
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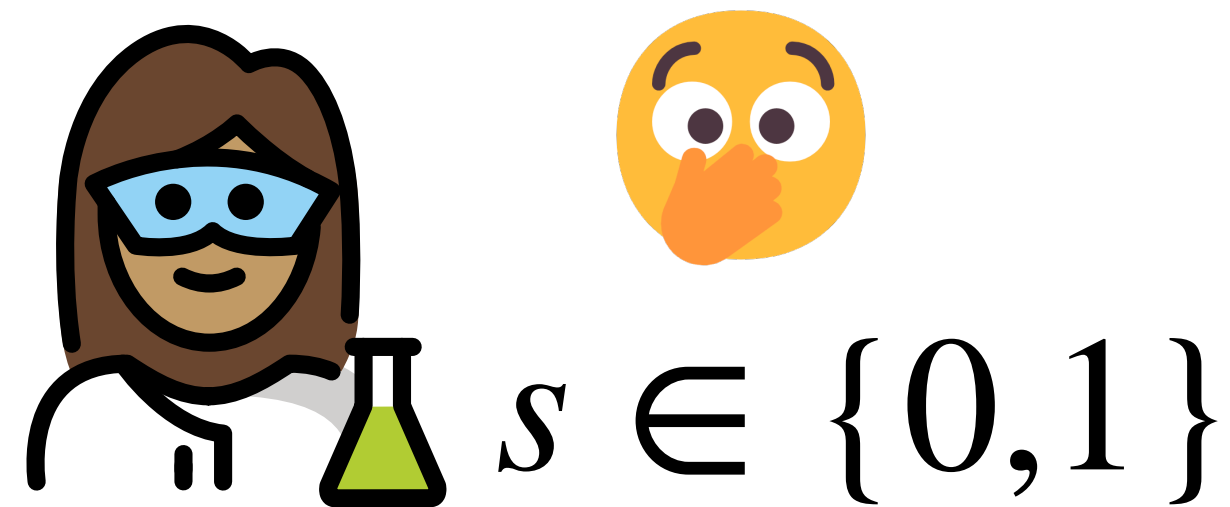
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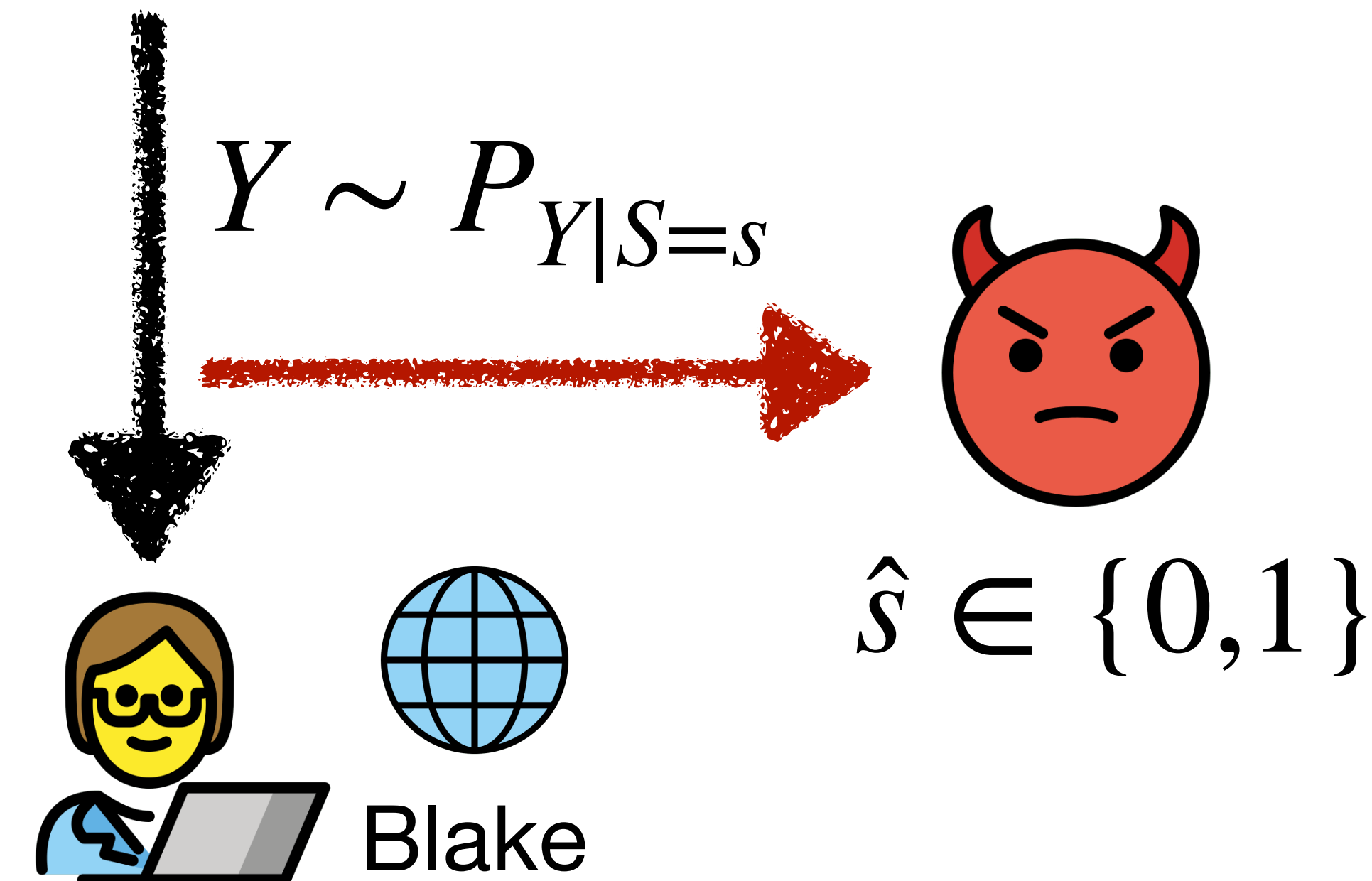
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The **privacy question** is a **hypothesis testing question**:

$$\mathcal{H}_0: Y \sim P_{Y|S=0}$$

$$\mathcal{H}_1: Y \sim P_{Y|S=1}$$





**The Lake of Hakone in  
Sagami Province**

**相州箱根湖水**

**Sōshū Hakone Kosui**

**Vista 1**

**hypothesis testing**



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## Example

$$\mathcal{H}_0: Y = 0 + Z \sim \mathcal{N}(0, \sigma^2)$$

$$\mathcal{H}_1: Y = 1 + Z \sim \mathcal{N}(1, \sigma^2)$$



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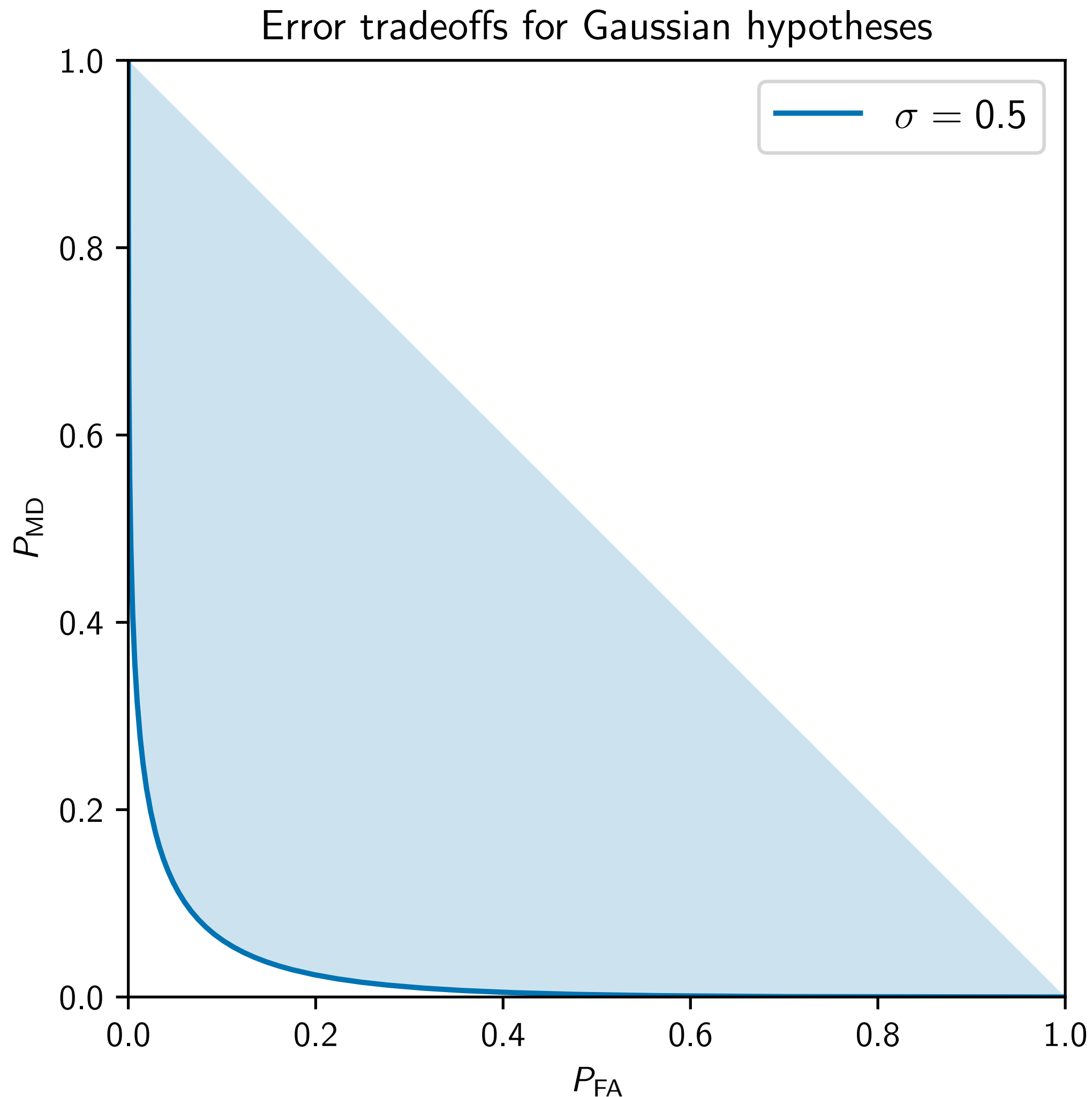
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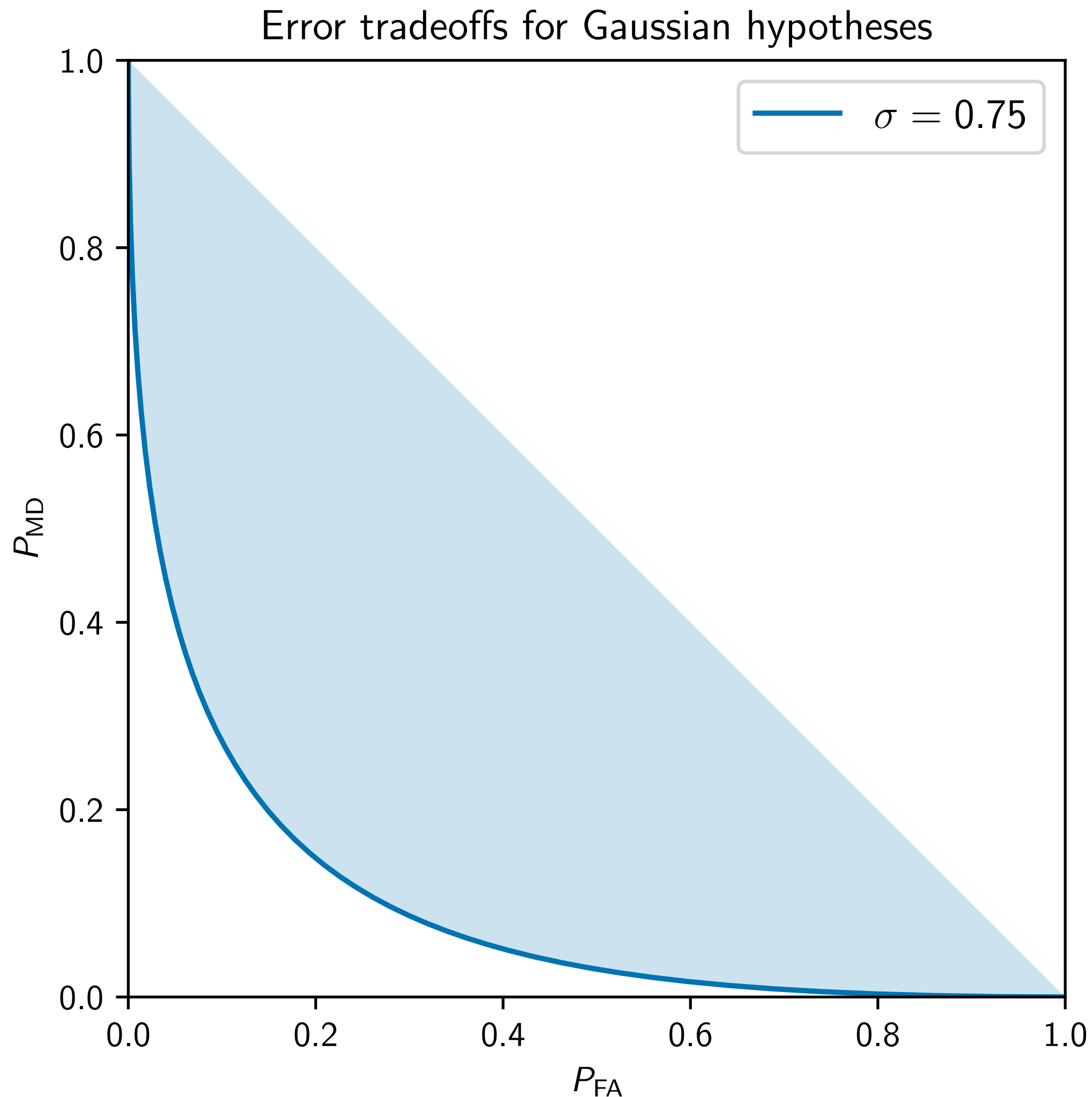
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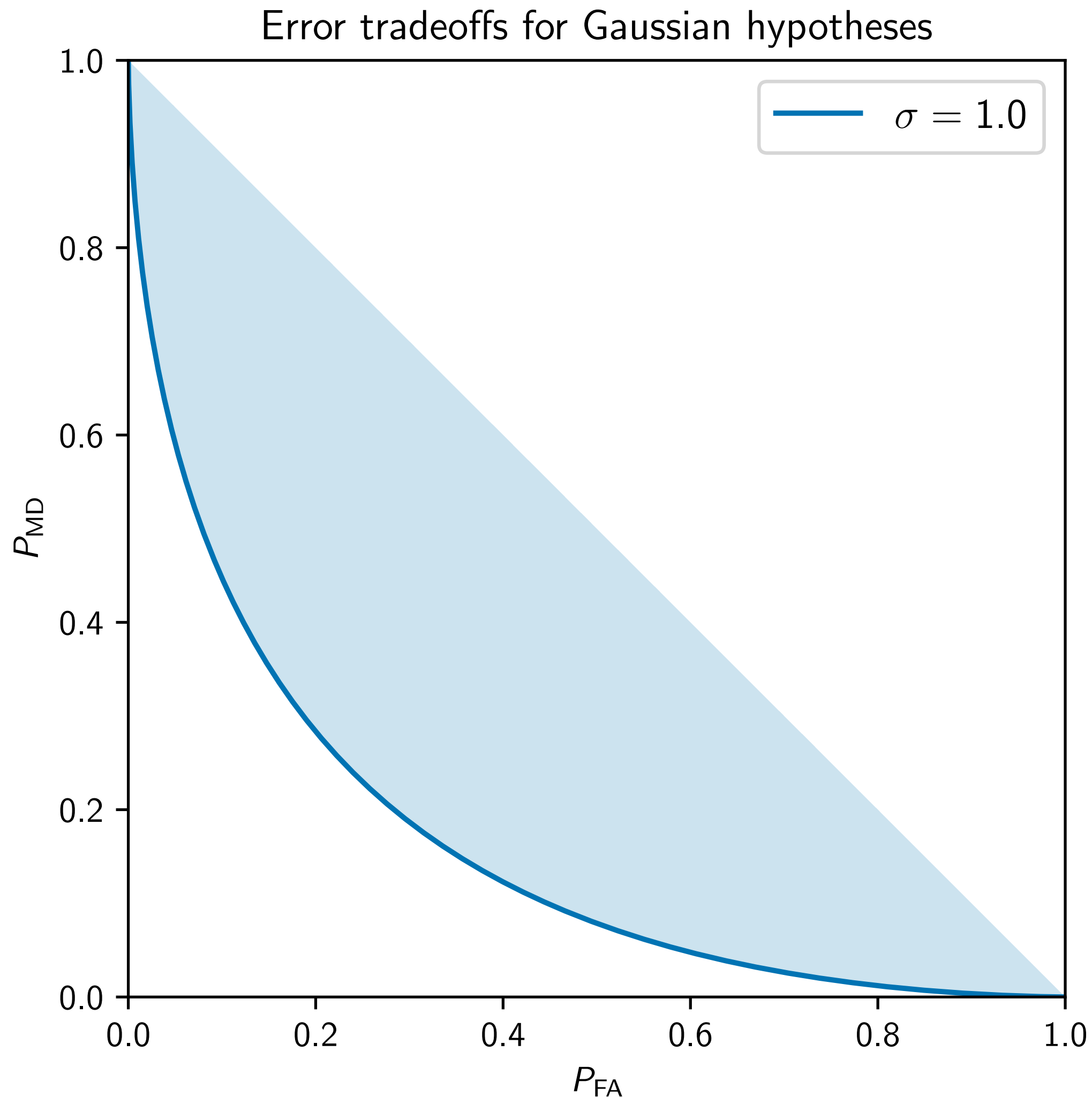
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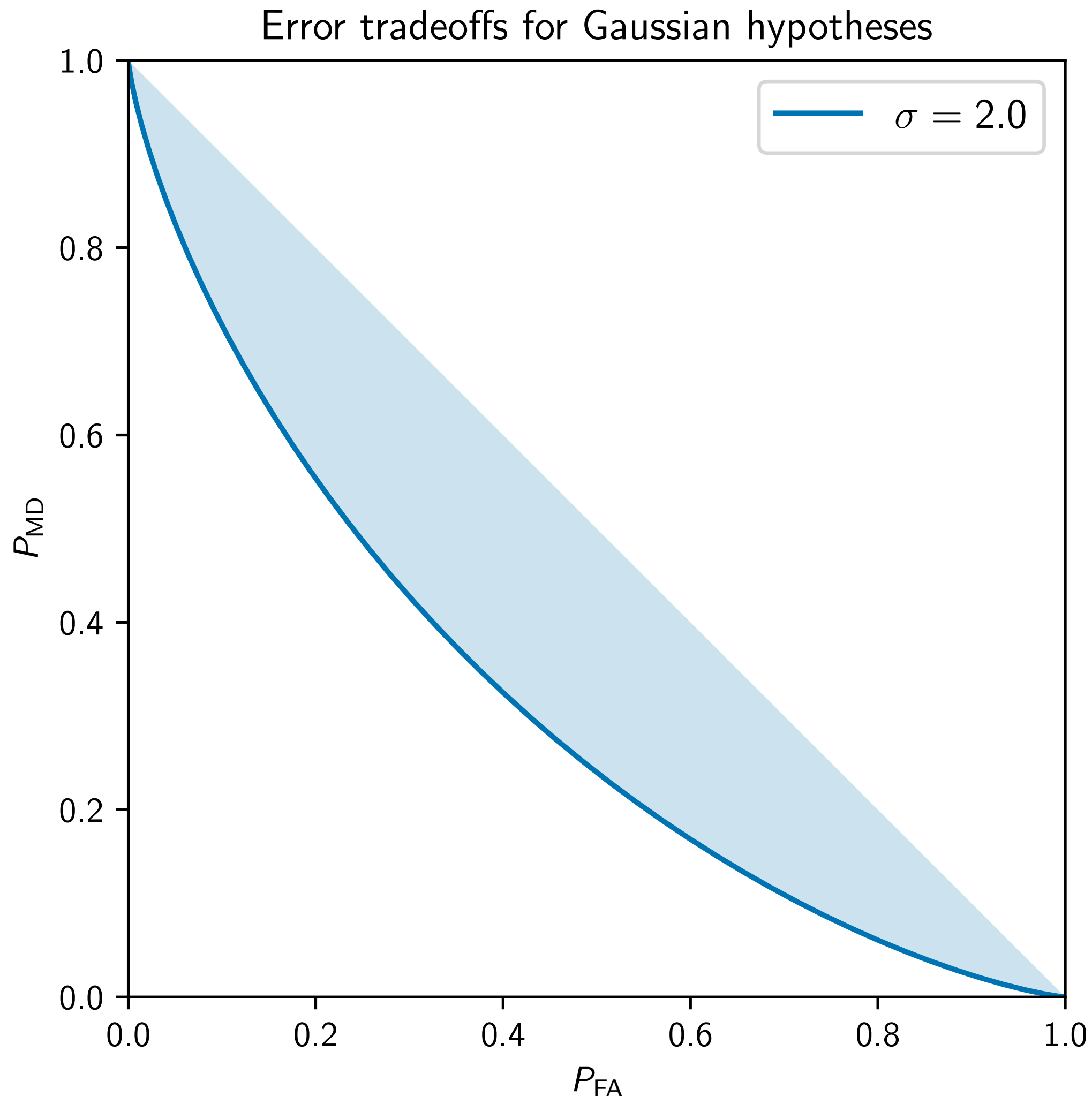
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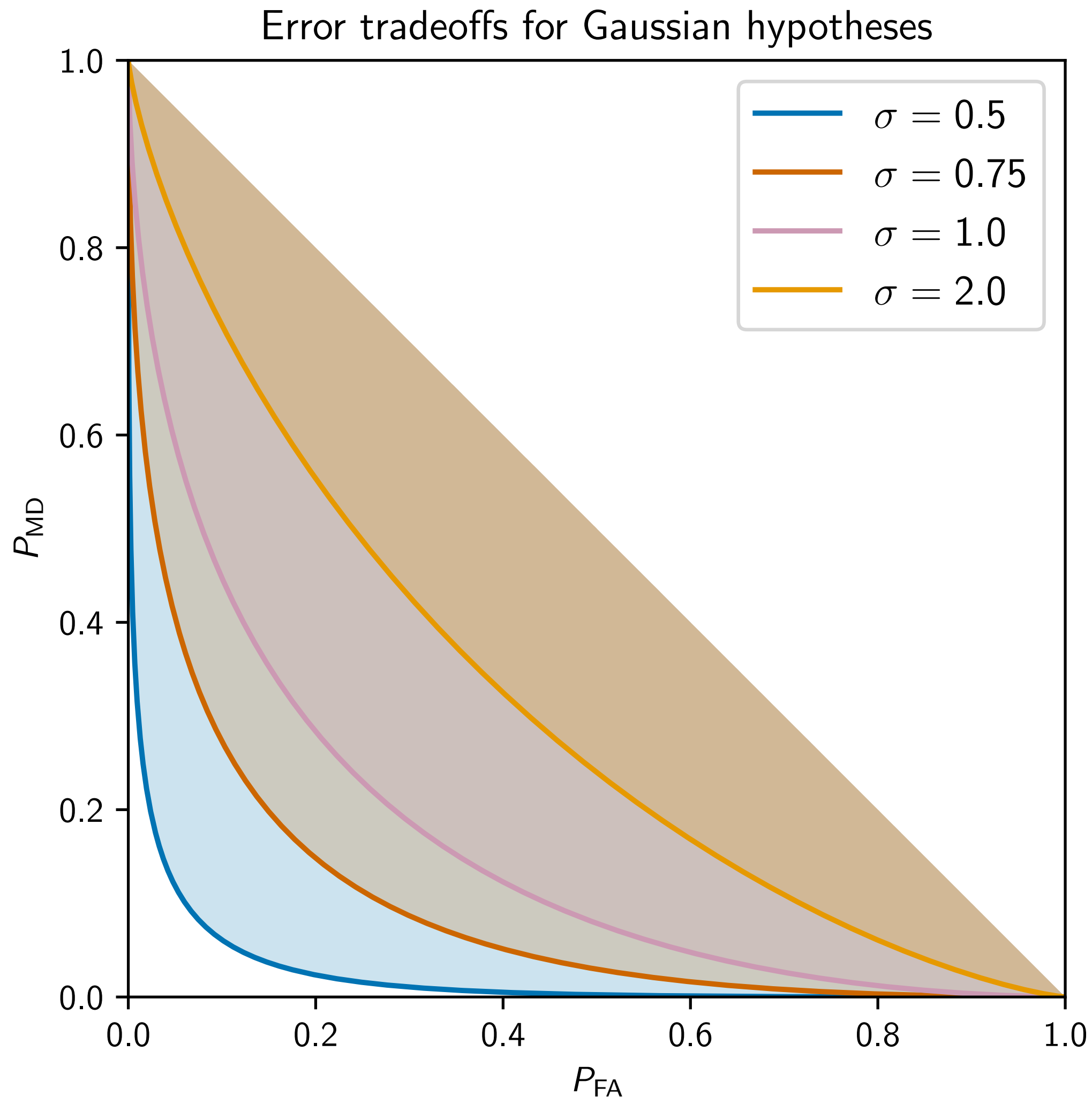
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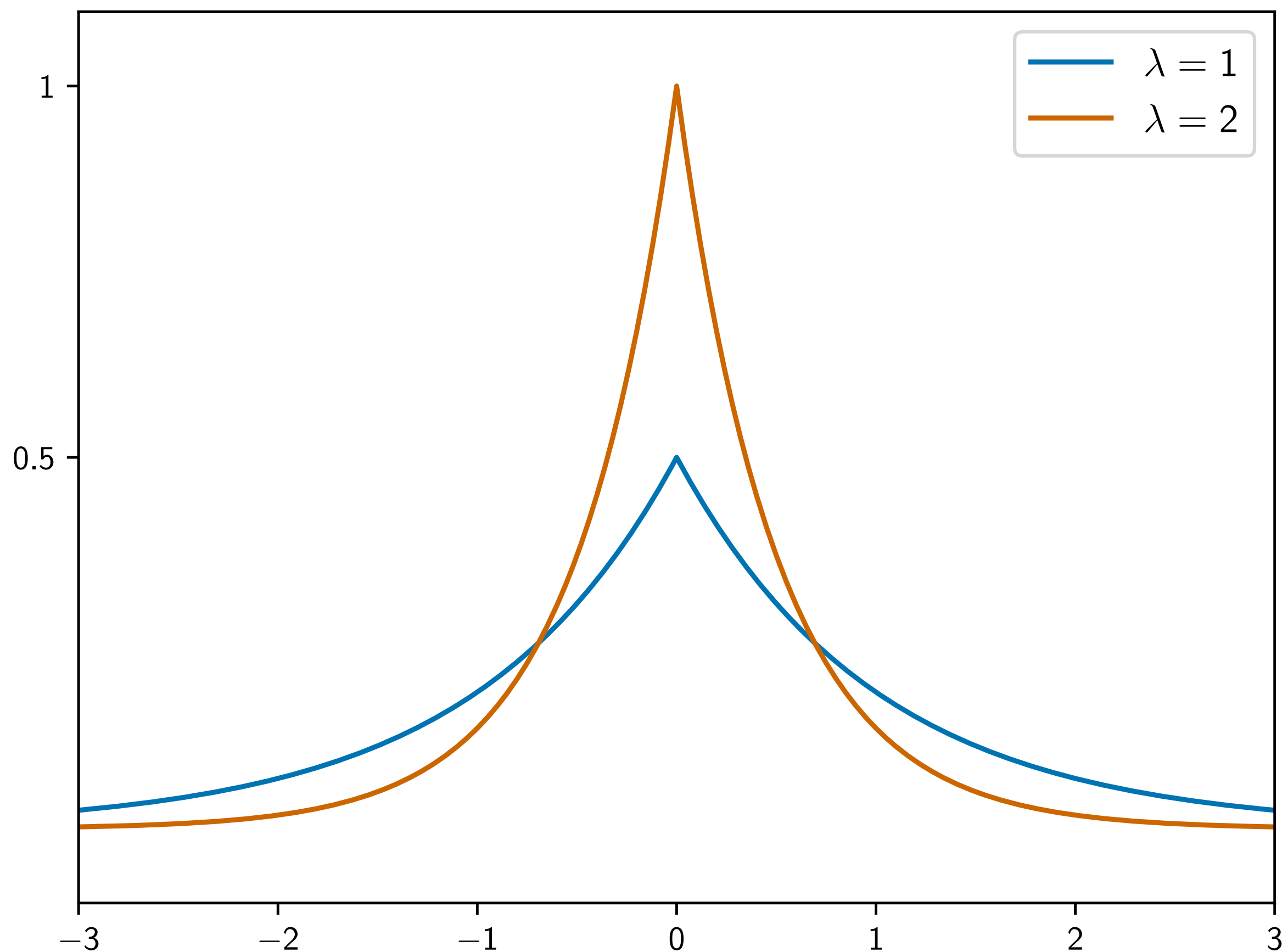
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Examples of Laplace distributions



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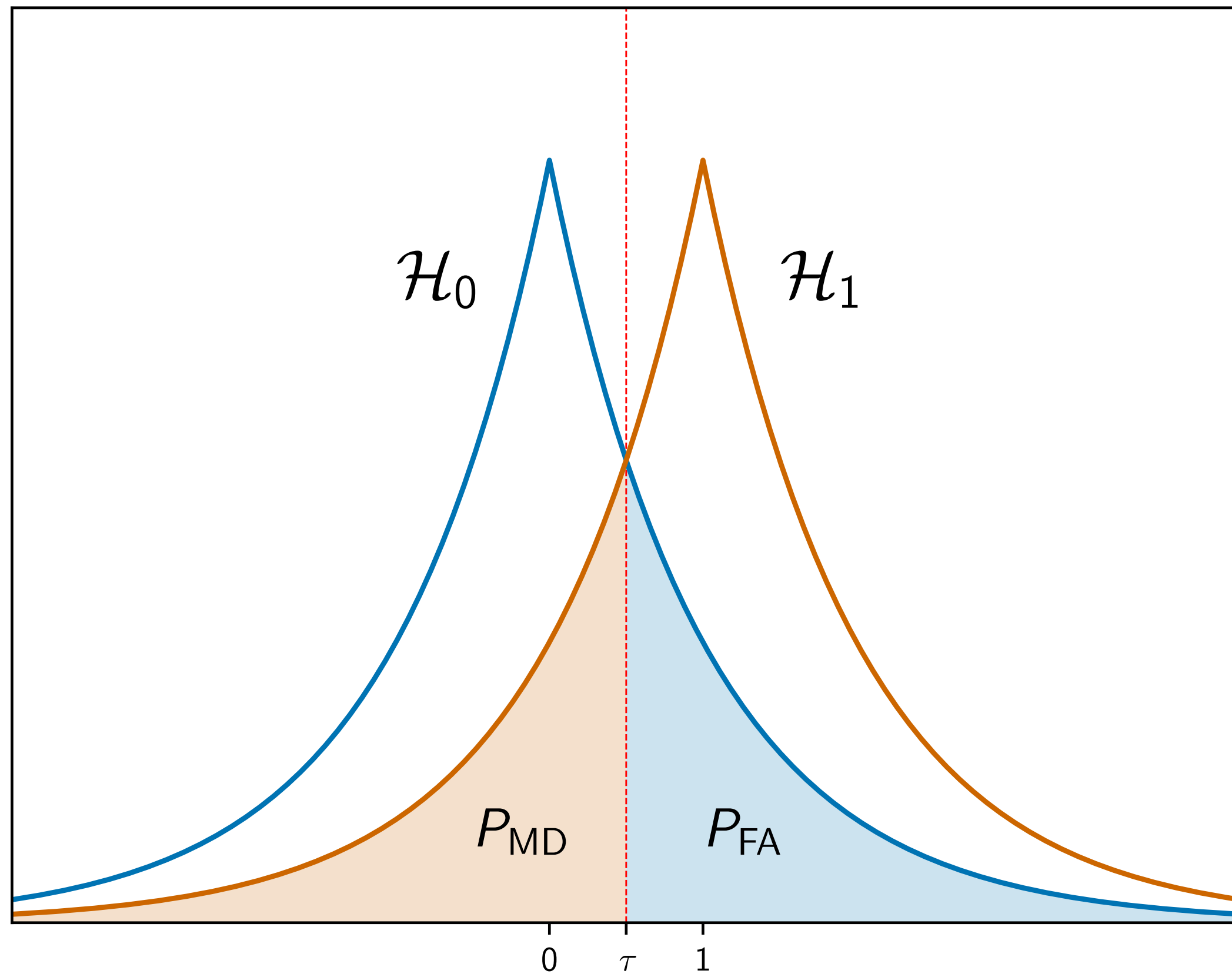
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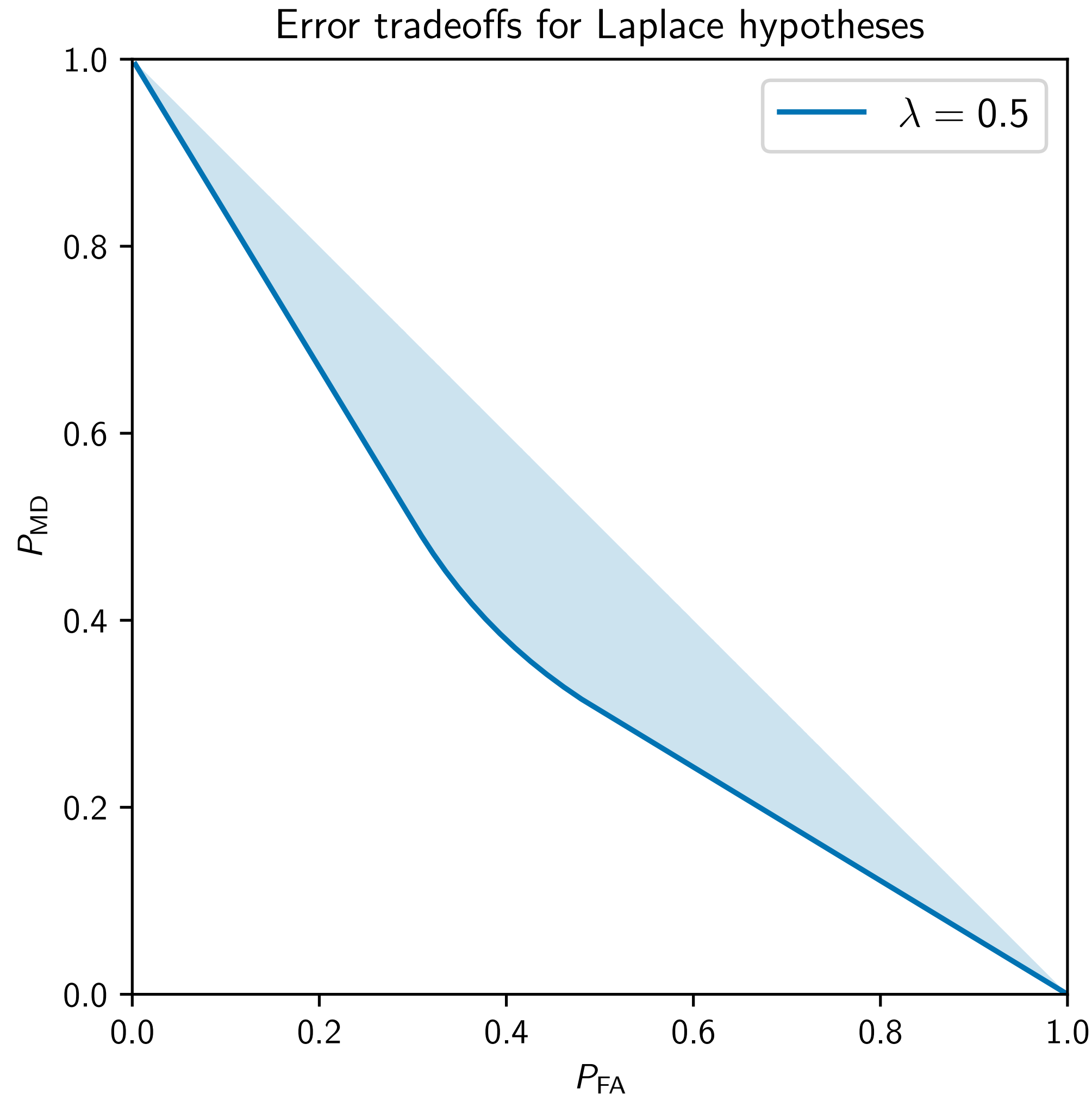
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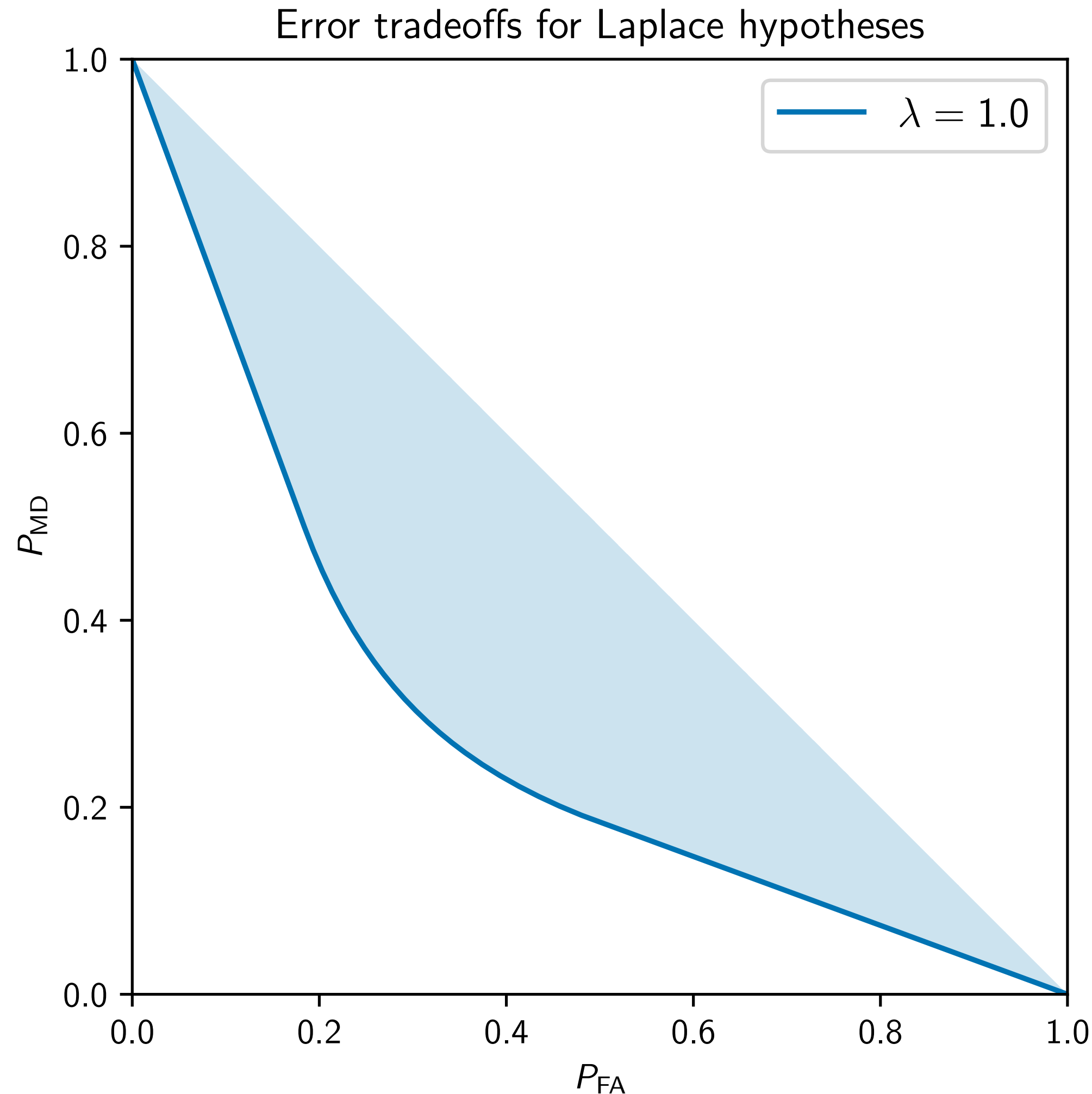
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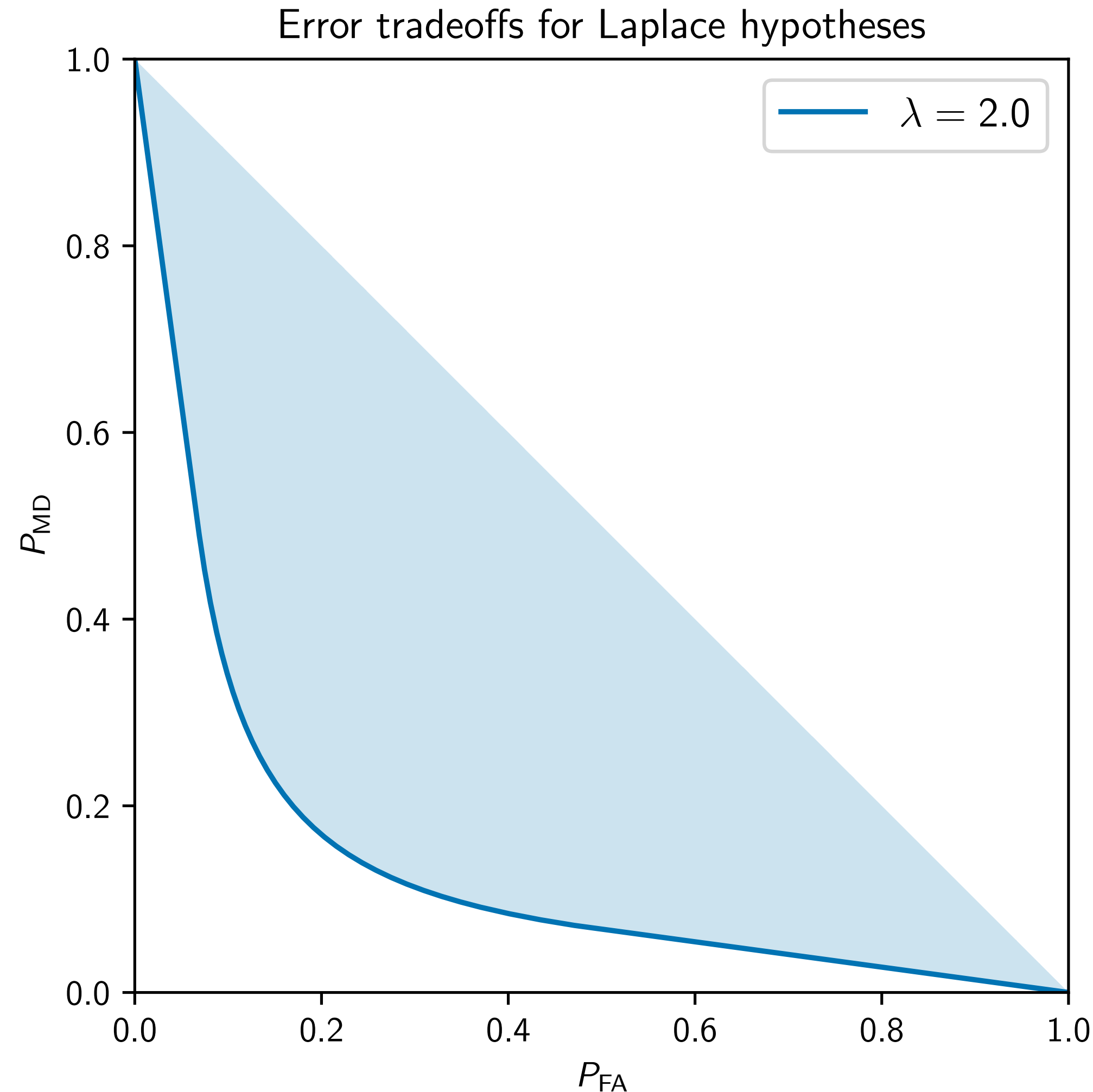
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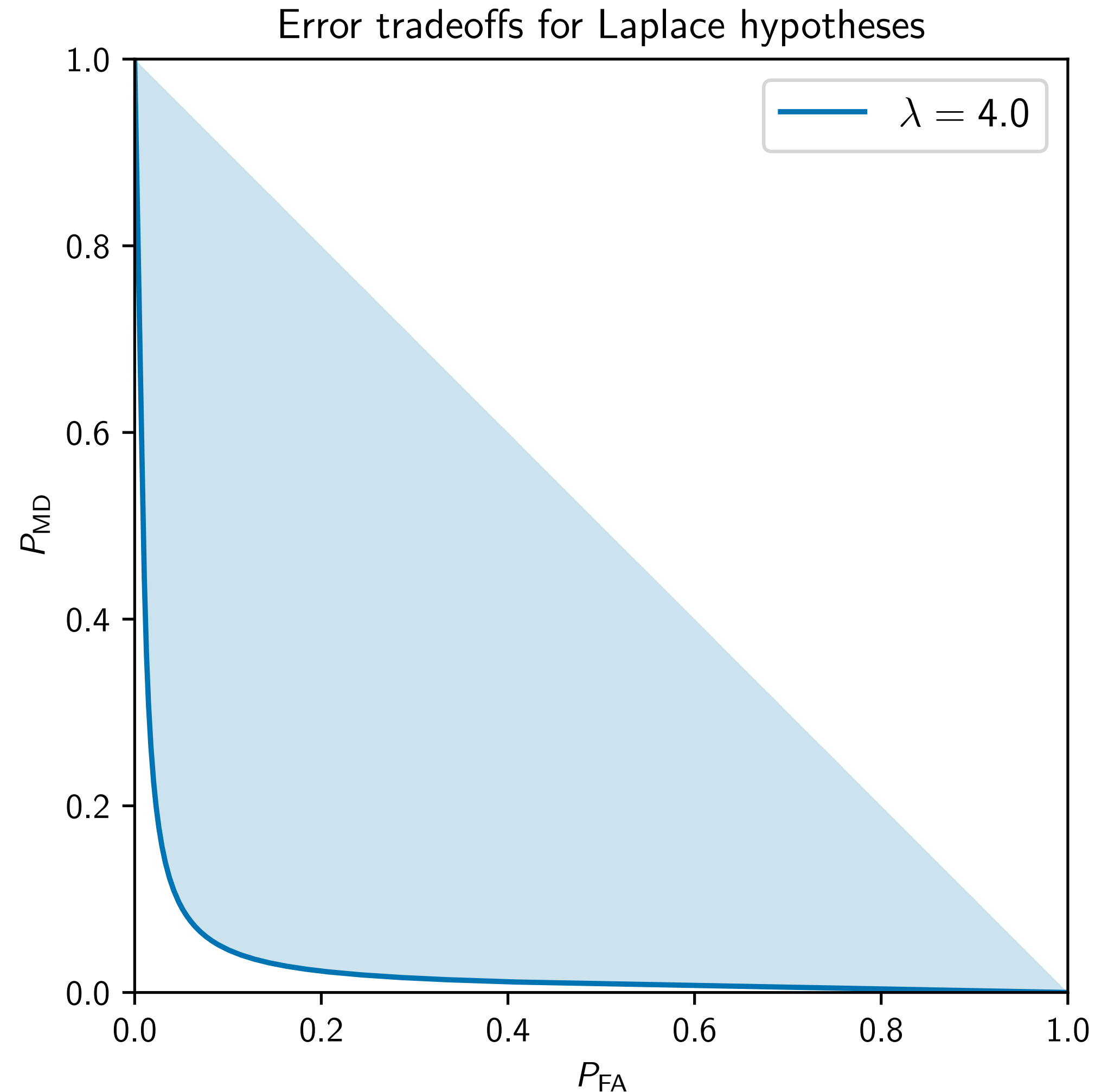
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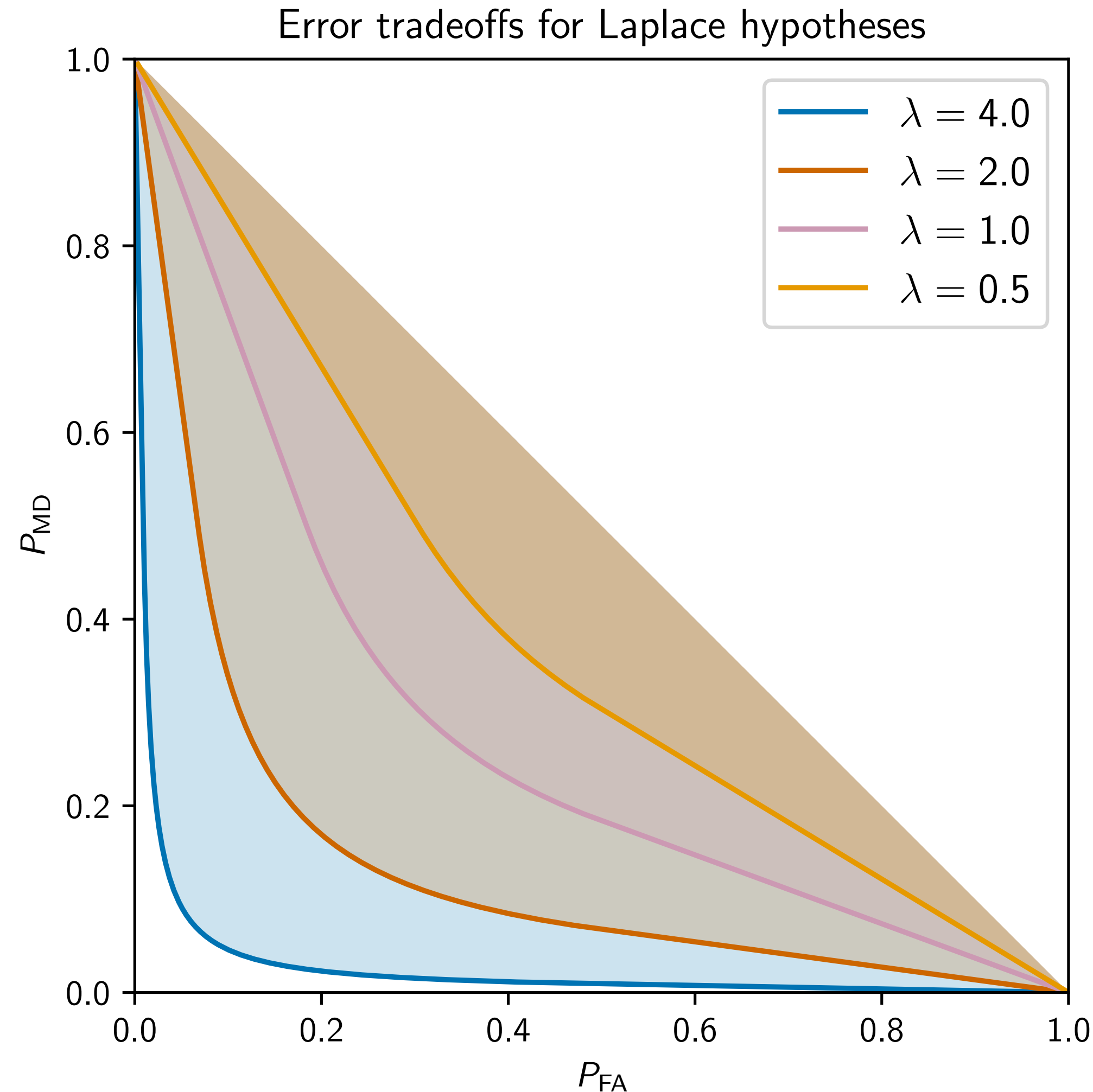
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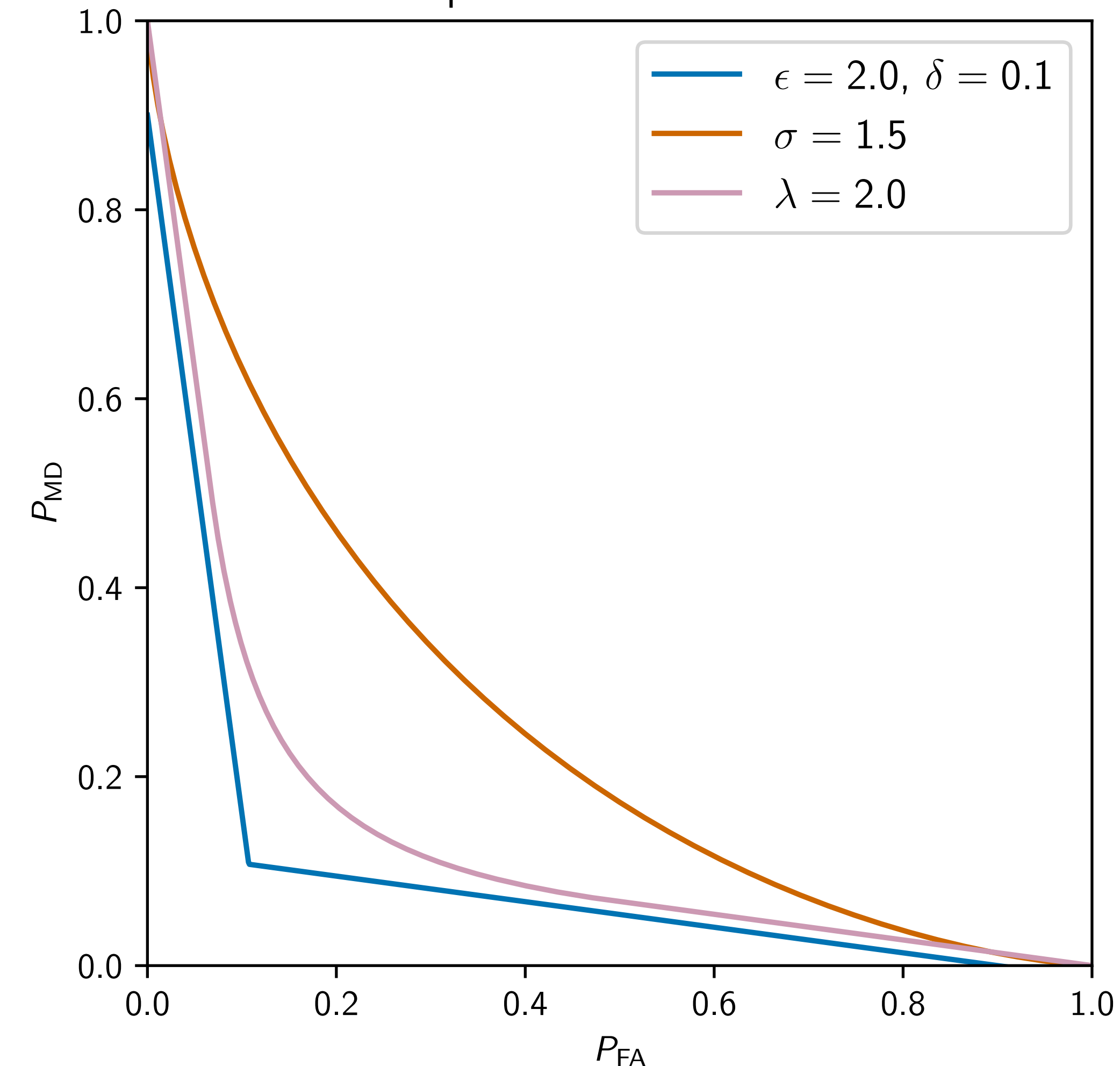
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Comparison of error tradeoffs



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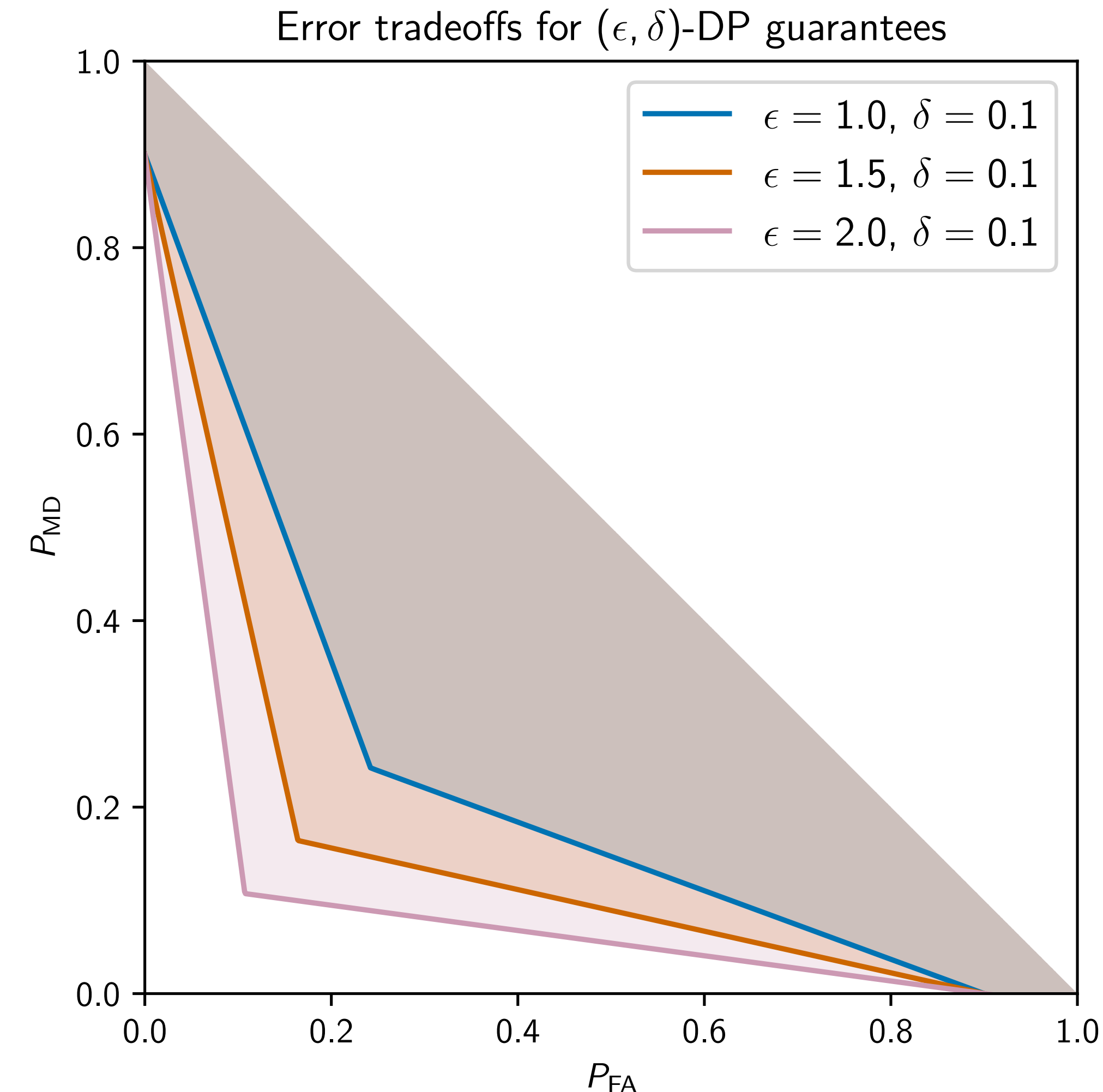
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Wasserman, Zhou (2010) Kairouz, Oh, Vishwanath (2015)

# Privacy versus testing

**We get to design the test!**



# Privacy versus testing

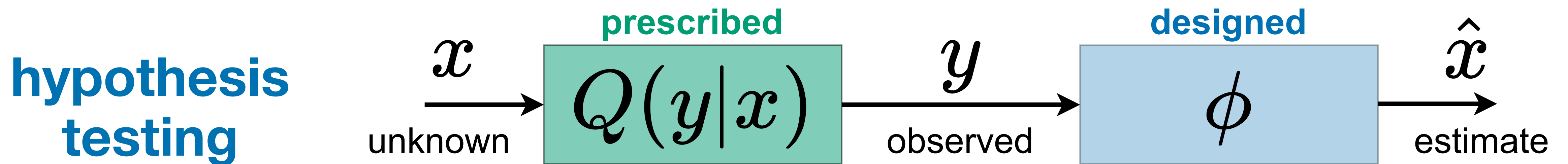
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The key difference between **hypothesis testing** (as we usually encounter it) and **(differential) privacy** is that we get to design the **likelihoods** but not the **test**!

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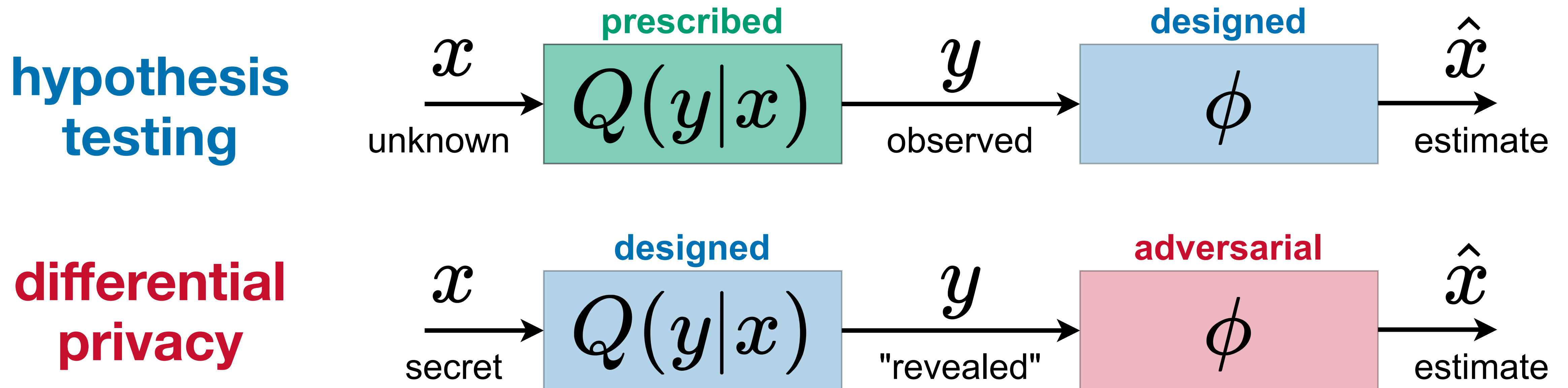
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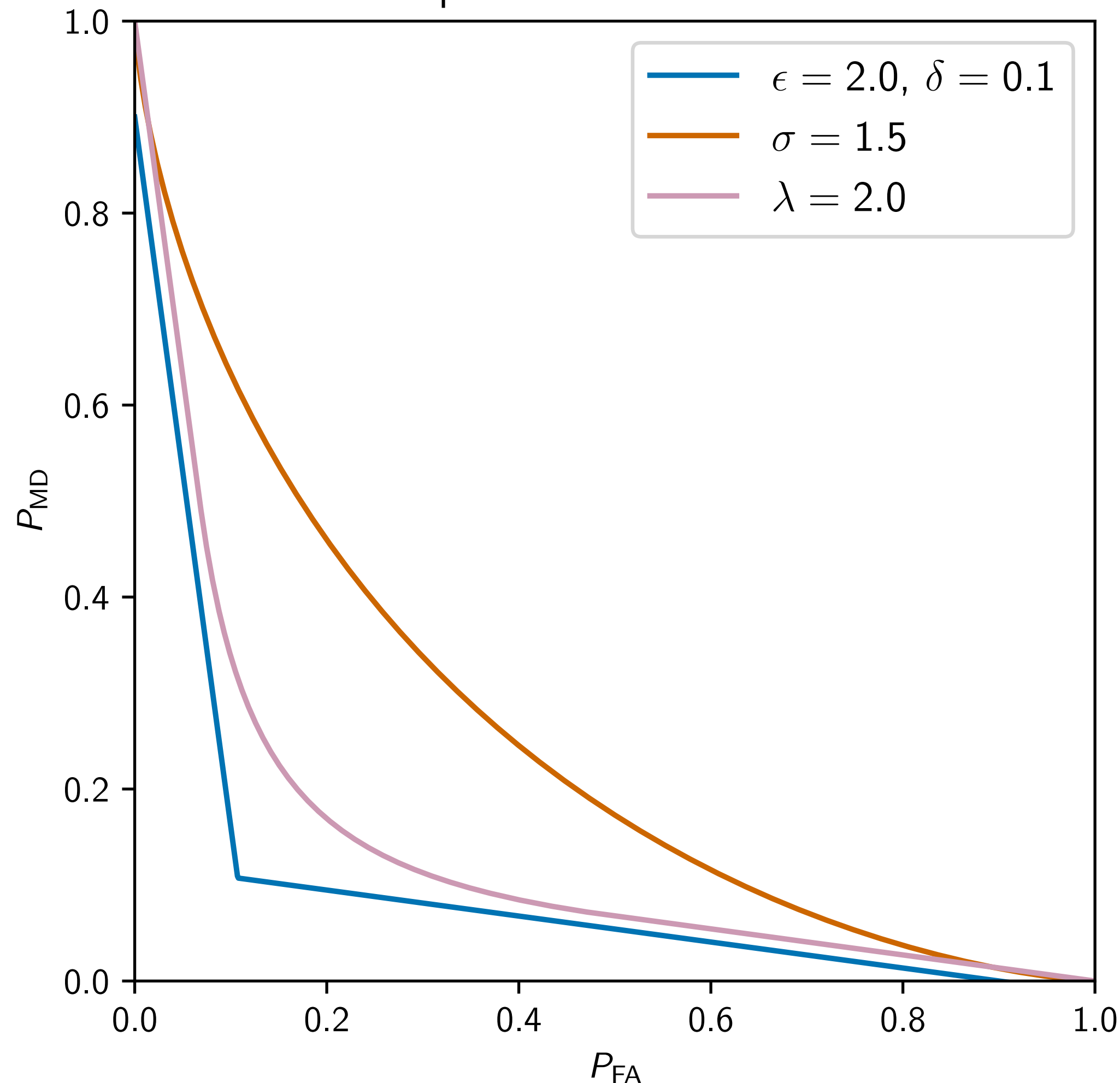
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Sunset Across Ryōgoku  
Bridge from  
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御厩川岸より西國橋夕陽  
見

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Vista 2

differential privacy the normal way



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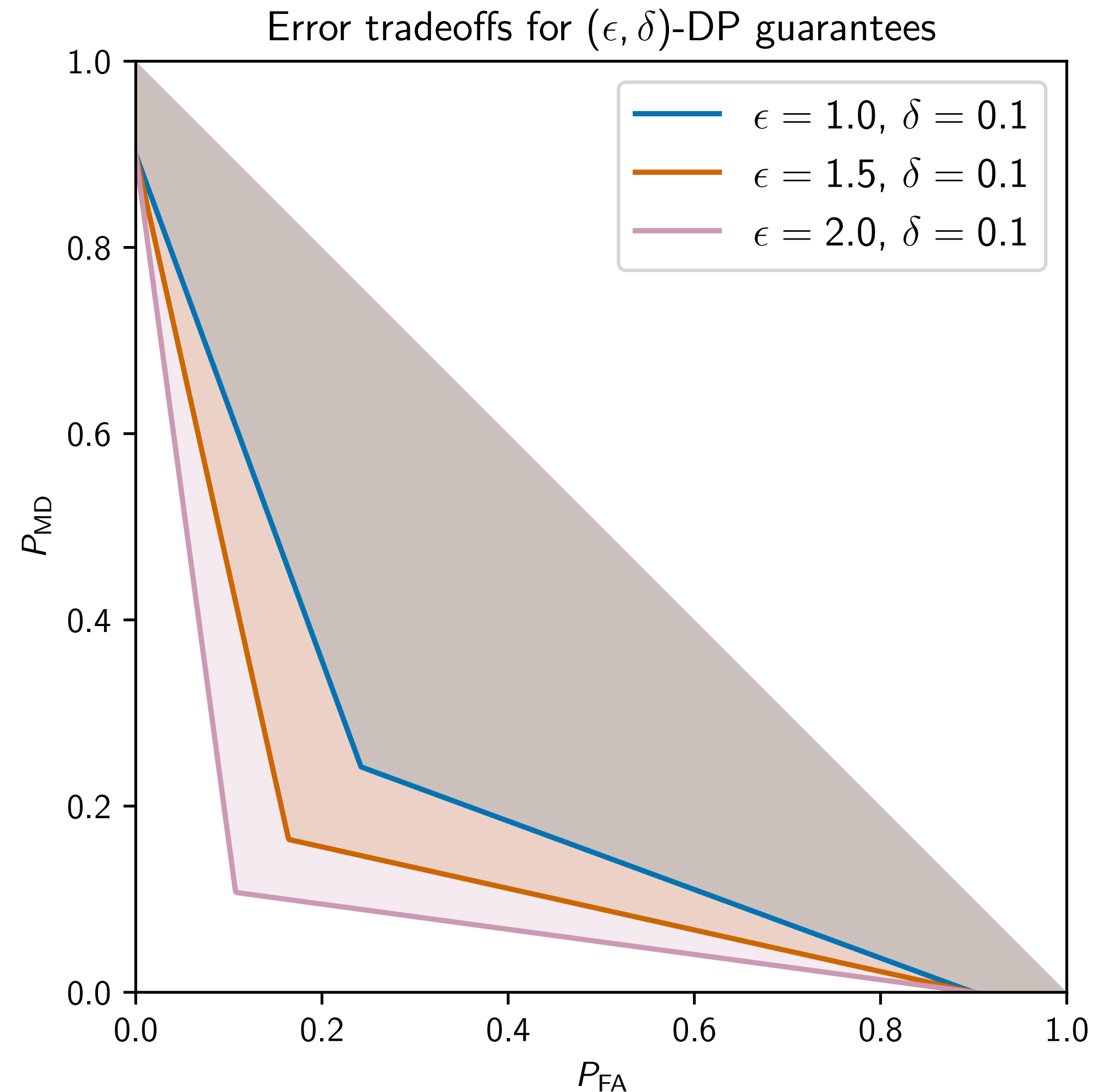
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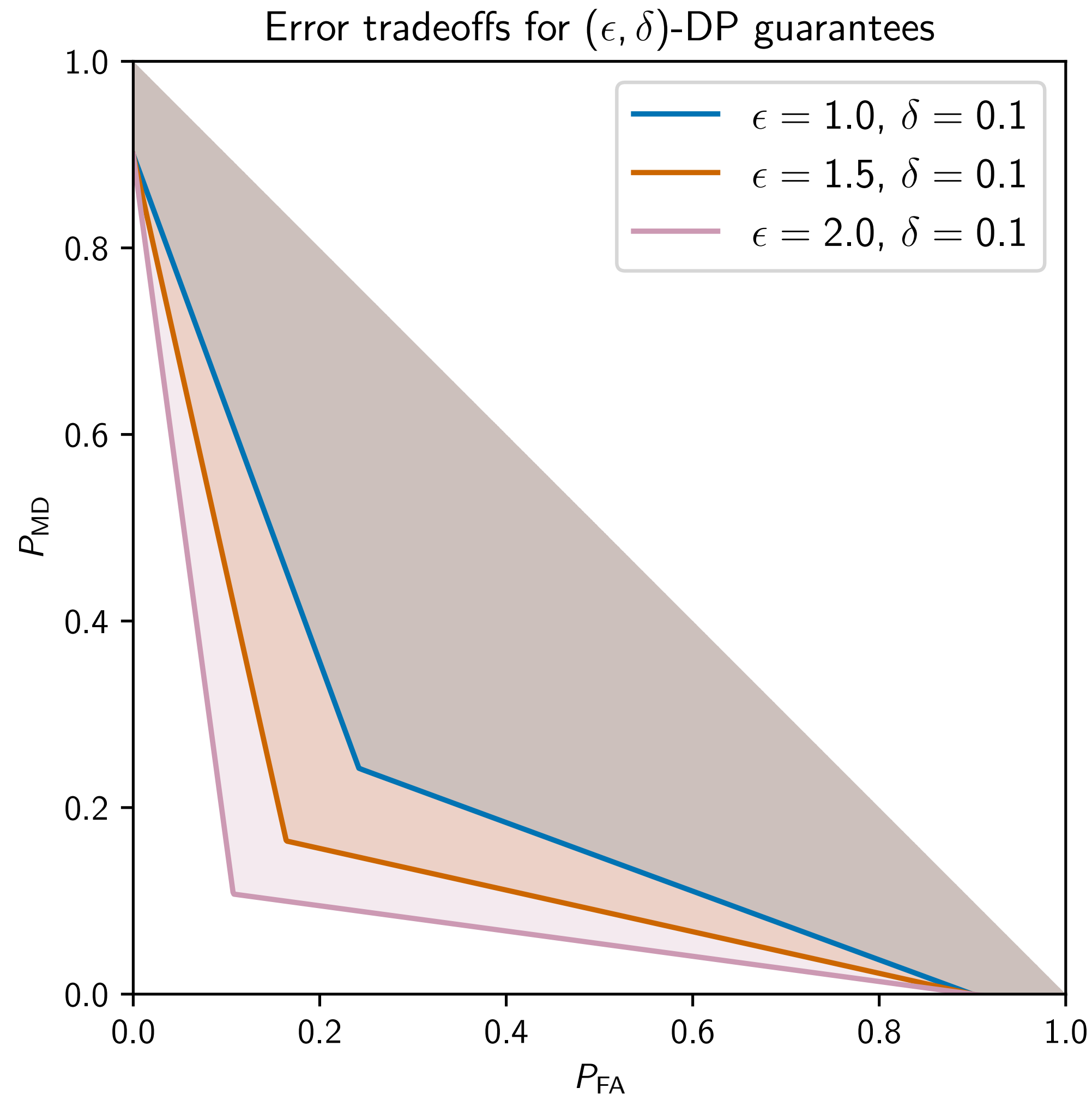
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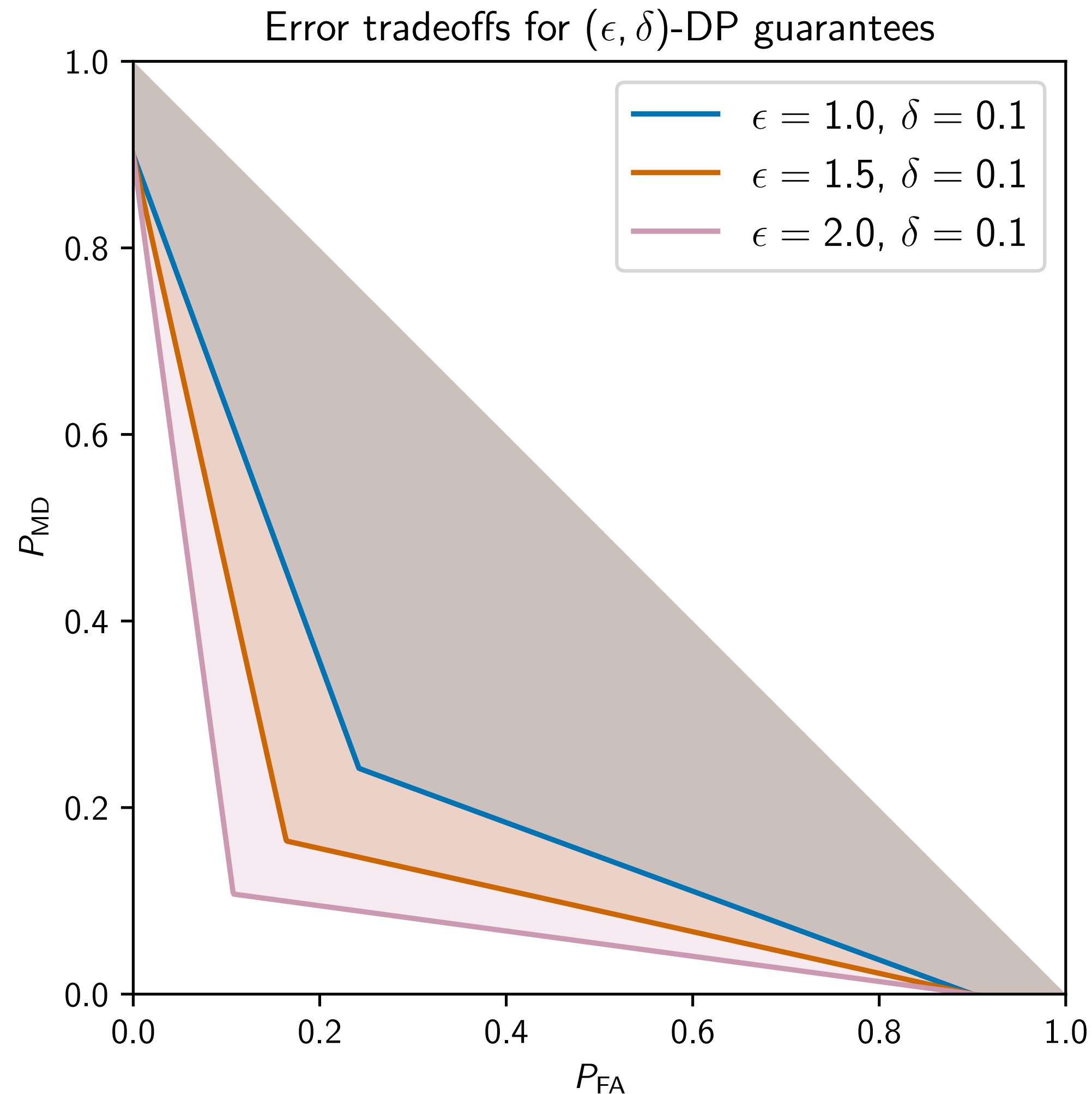
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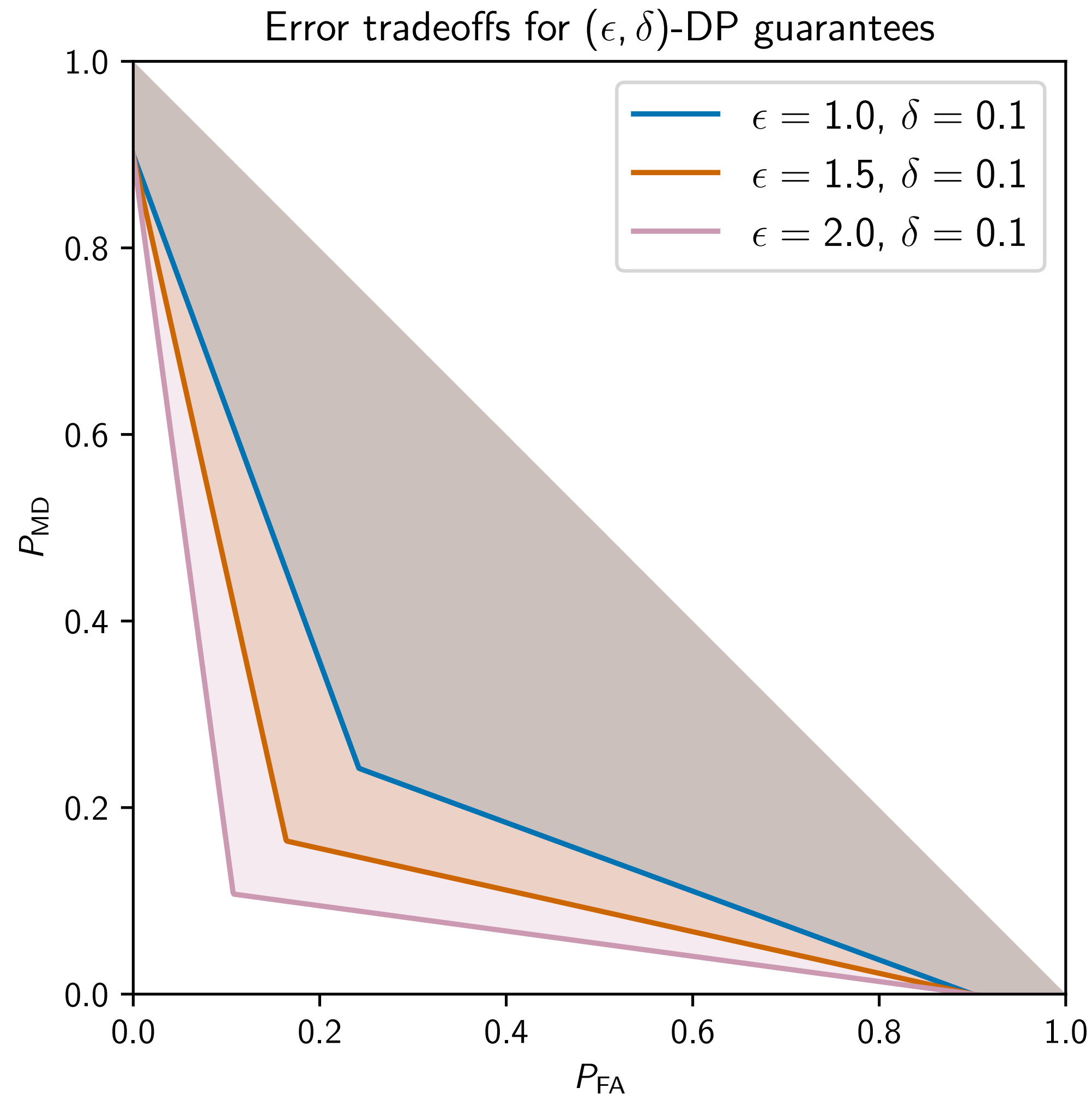


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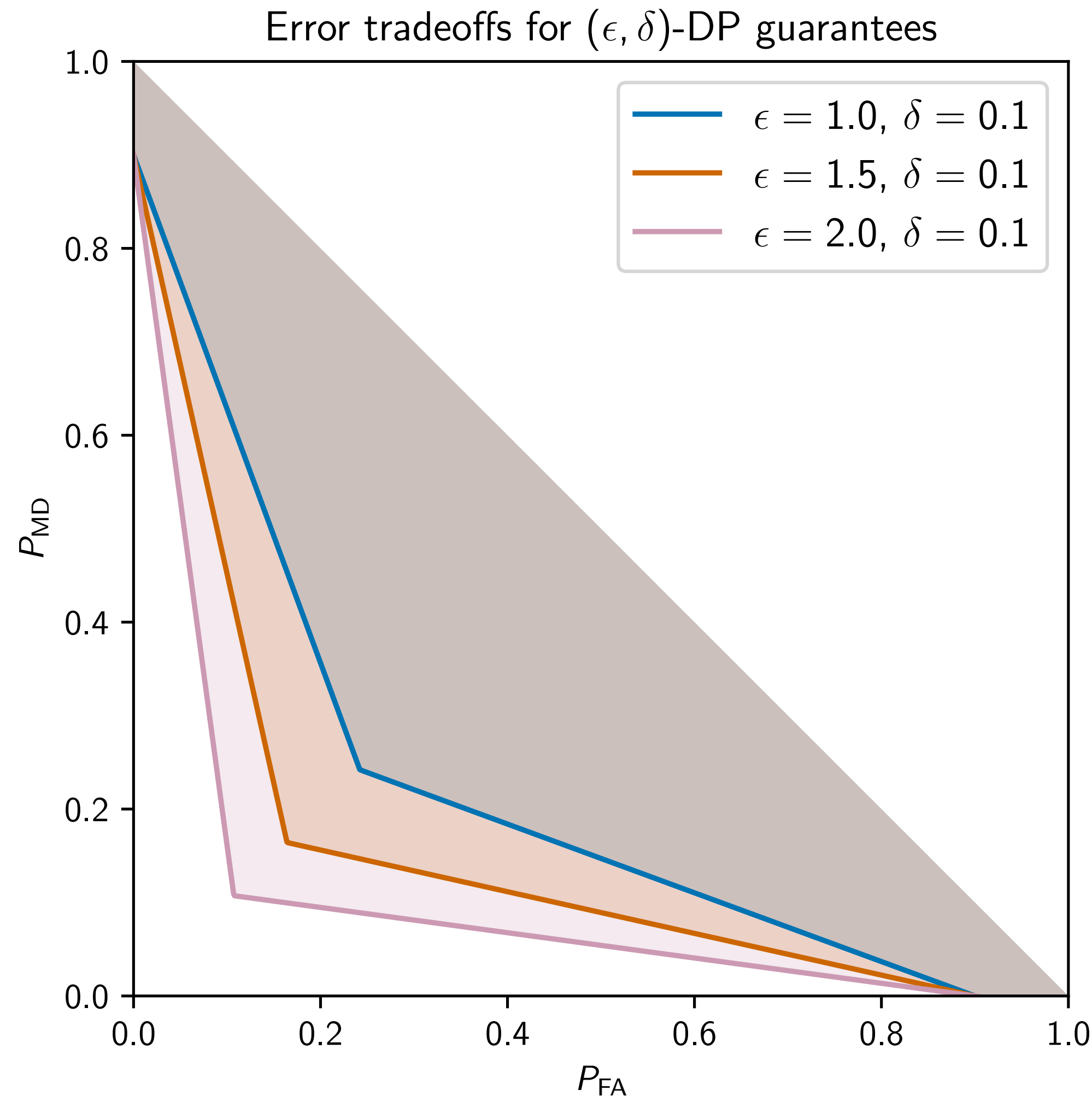
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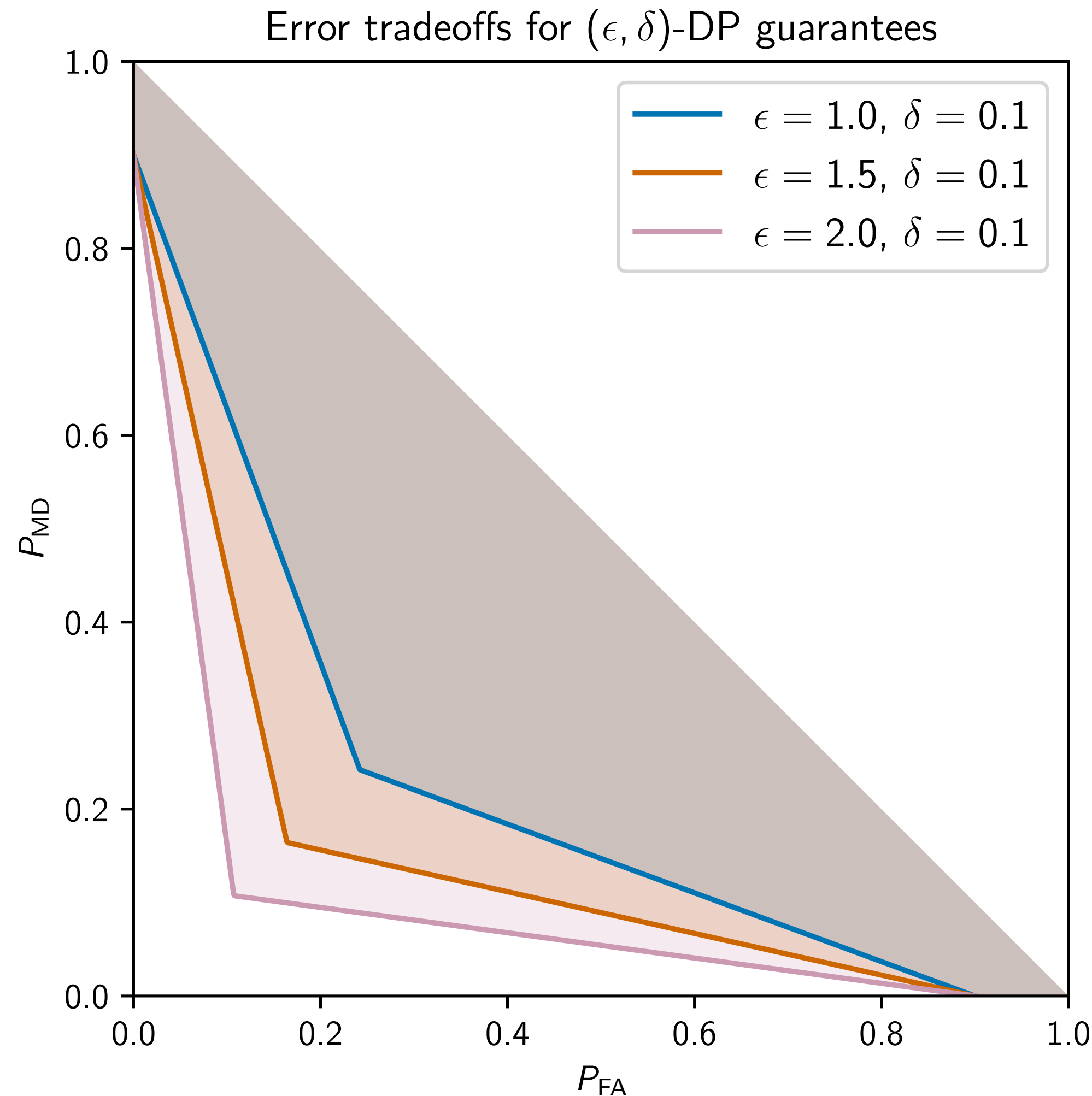
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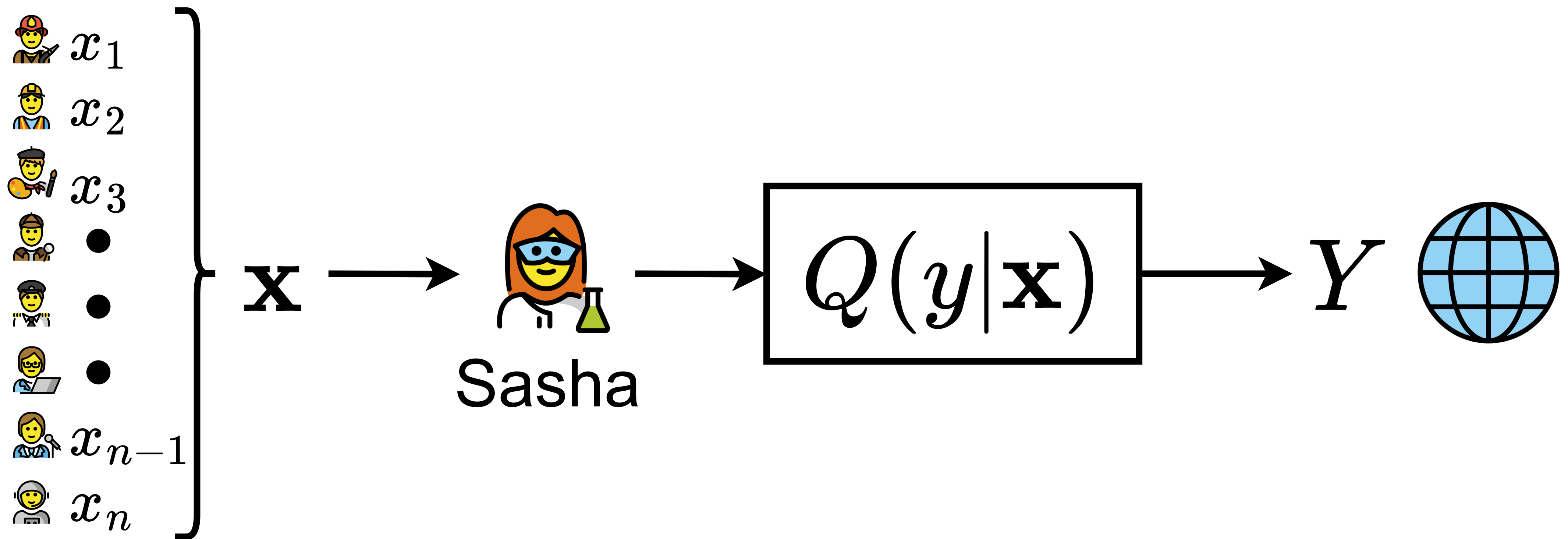
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When can we do this? When **neighboring data sets make similar output distributions**.

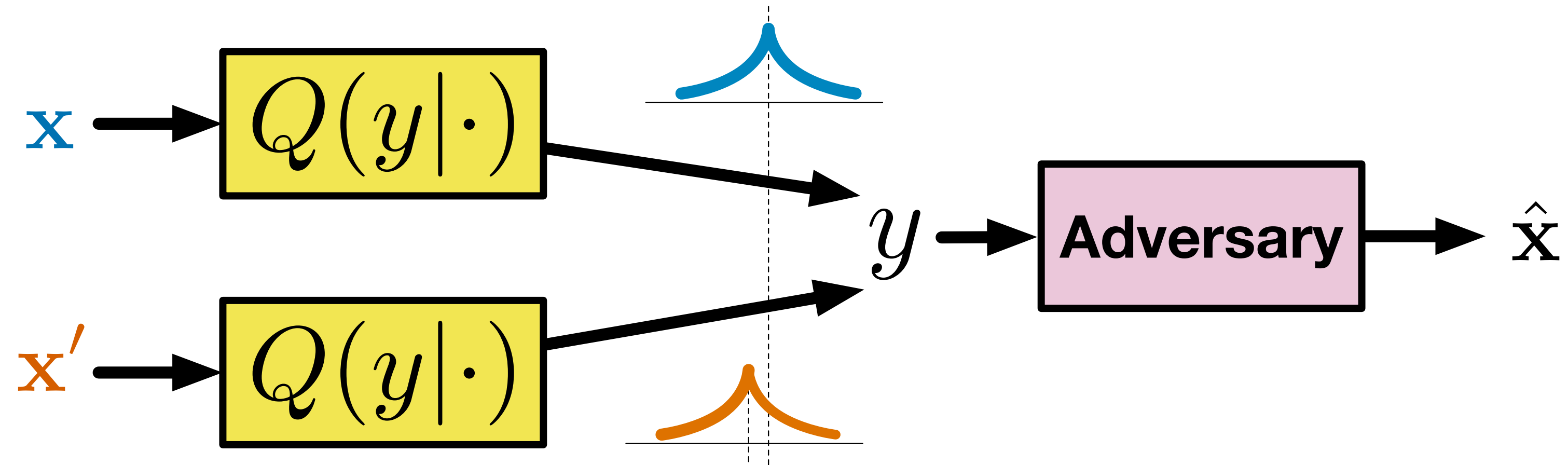
# In a snapshot

Replacing a single bit with a database



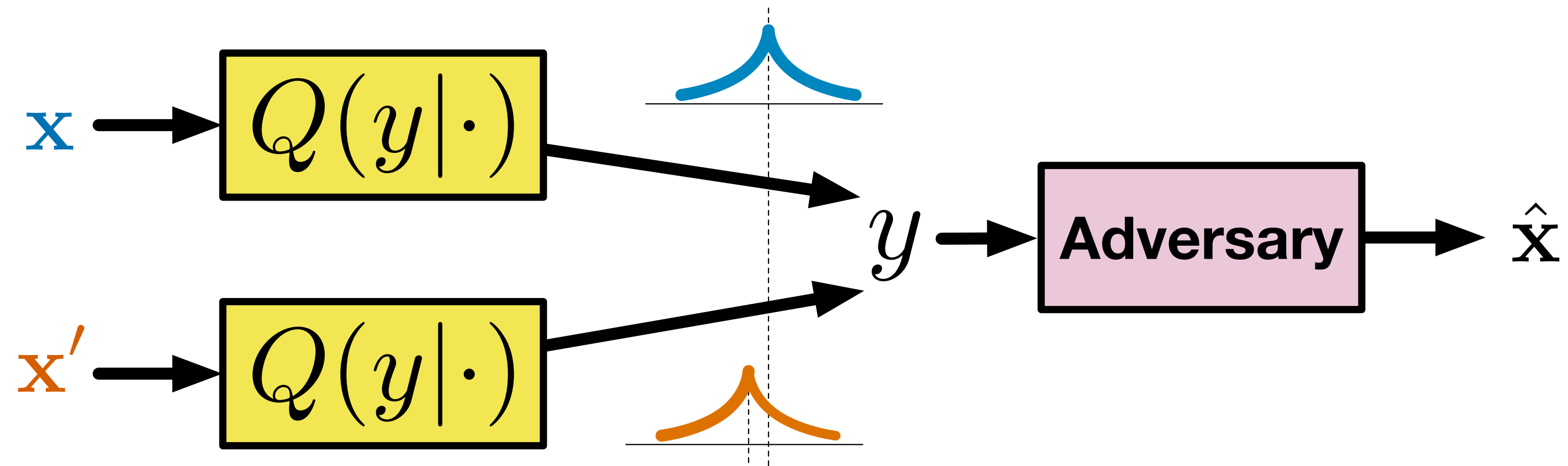
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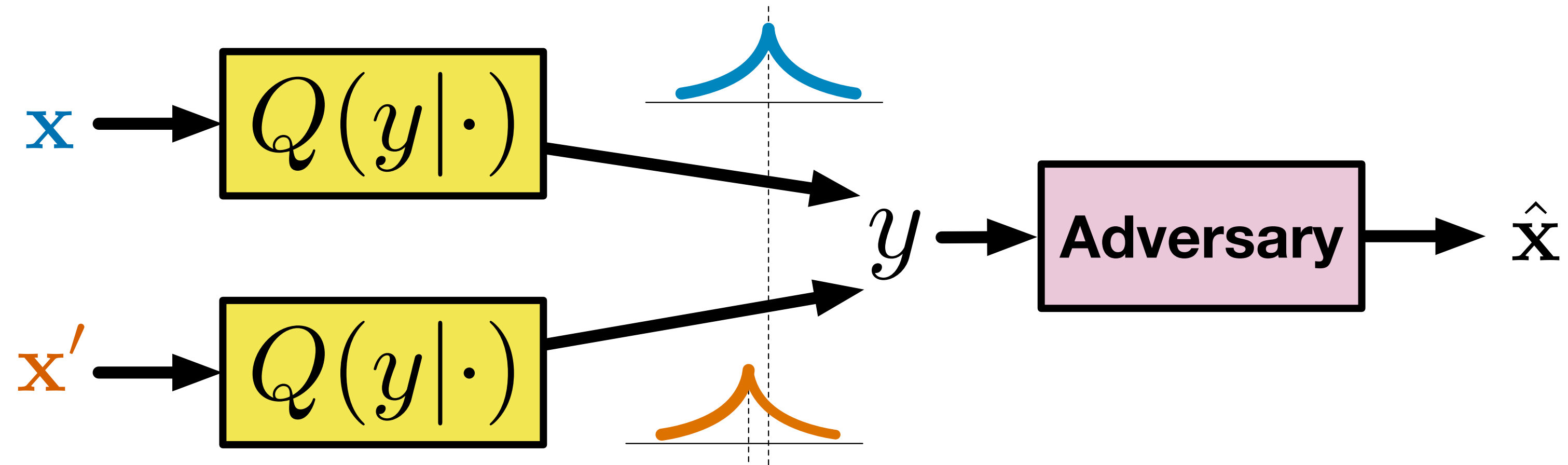
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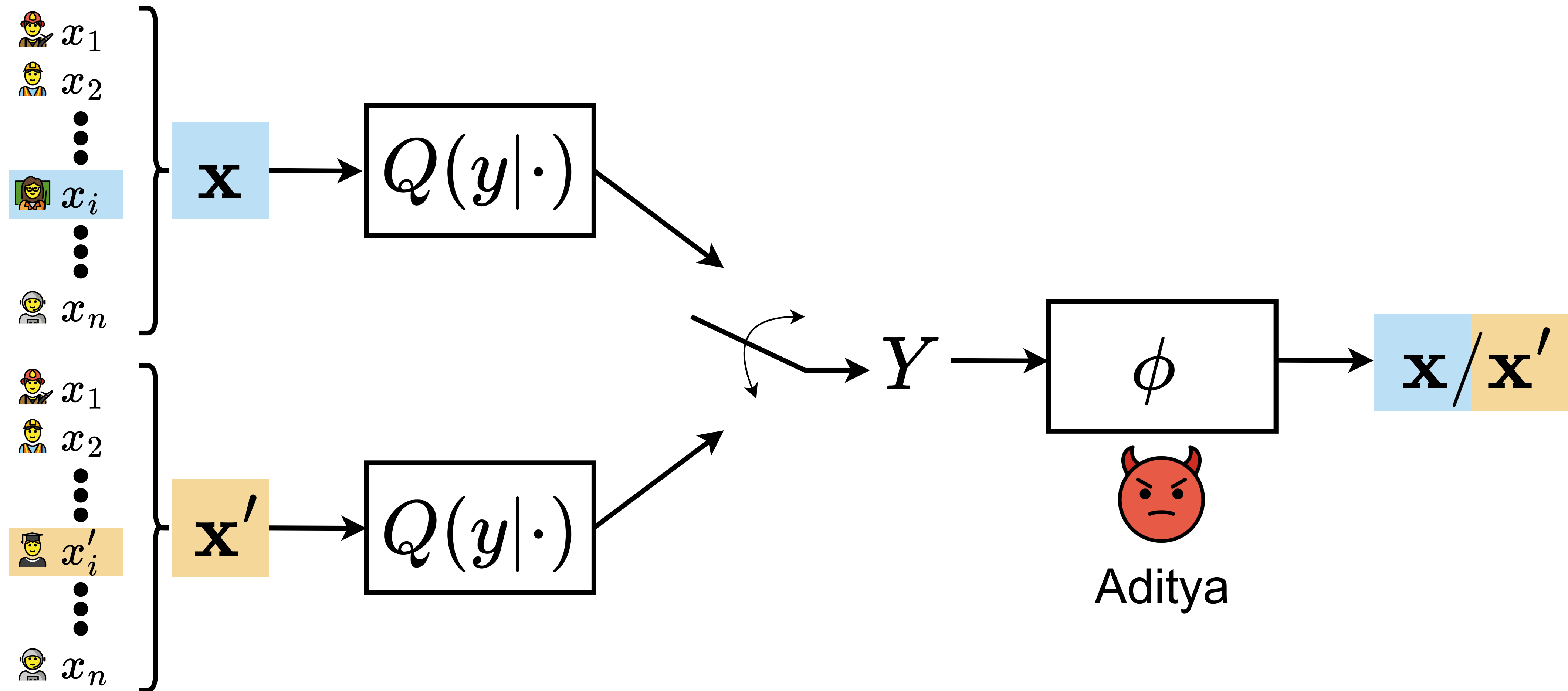
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[Dwork-Kenthapadi-McSherry-Mironov-Naor 2006]

[Wasserman-Zhou 2010]

# Neighboring datasets in a picture

The adversary's hypothesis test





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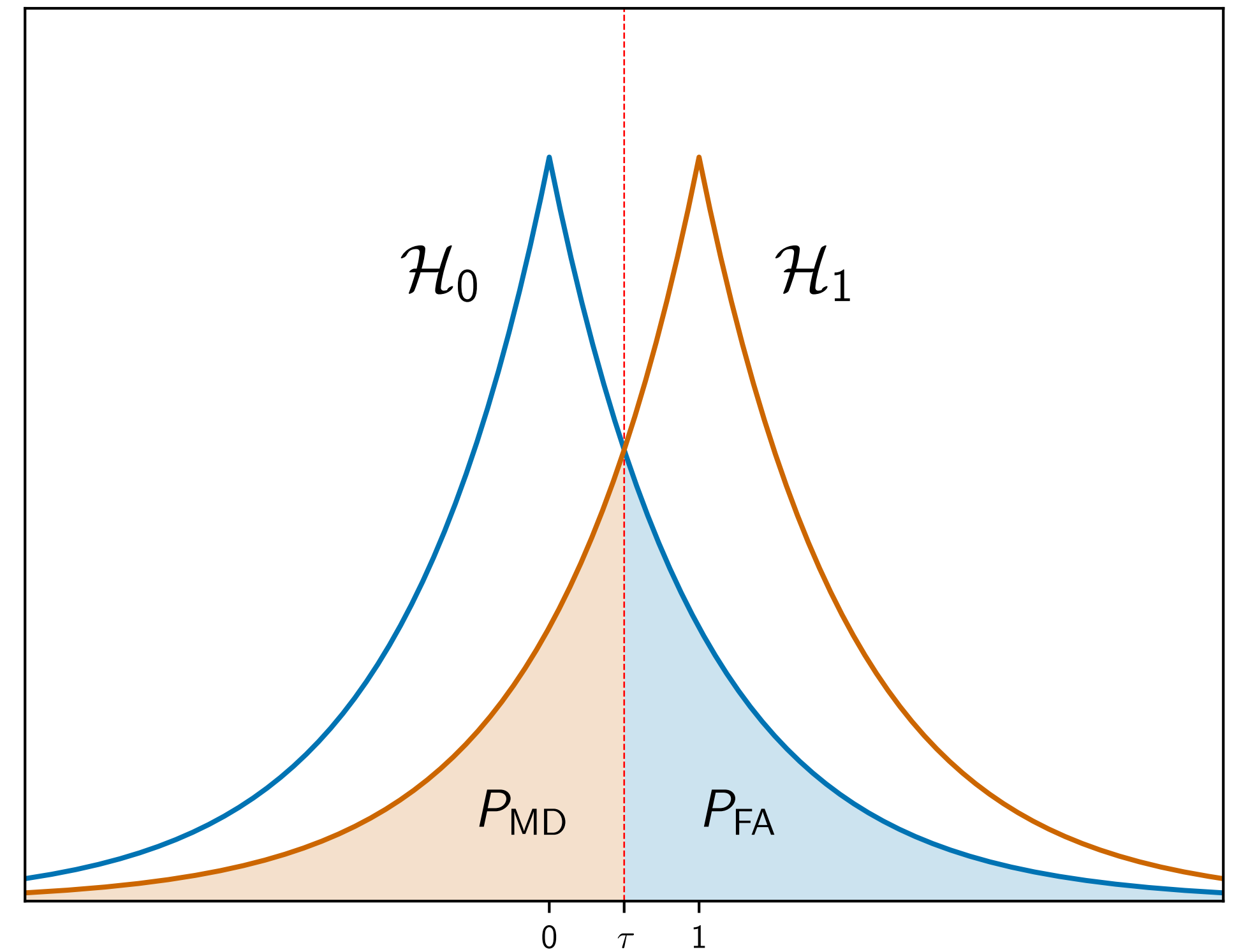
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- **The data itself is considered identifying:** no notion of some parts being personally identifiable information (PII) and others not.

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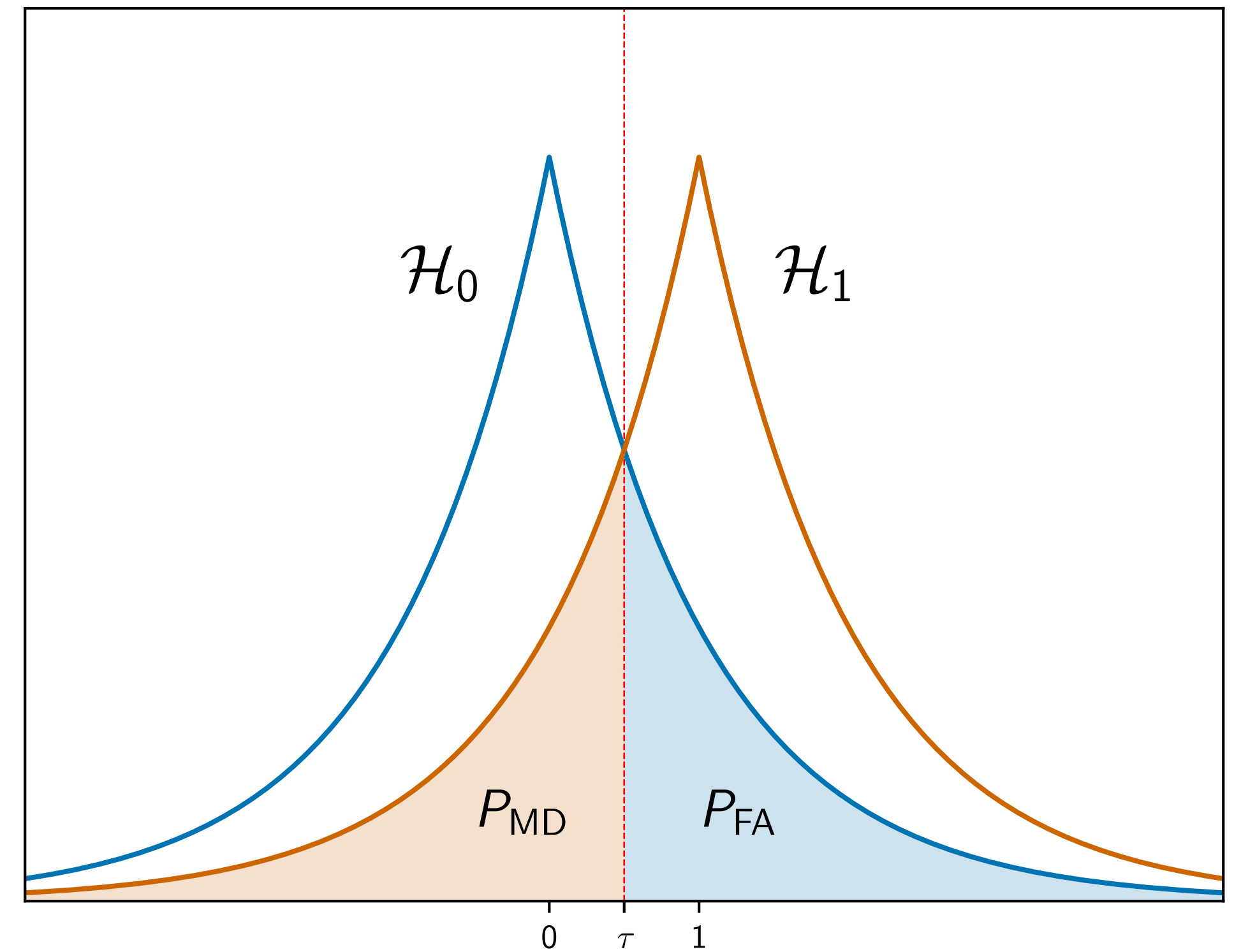
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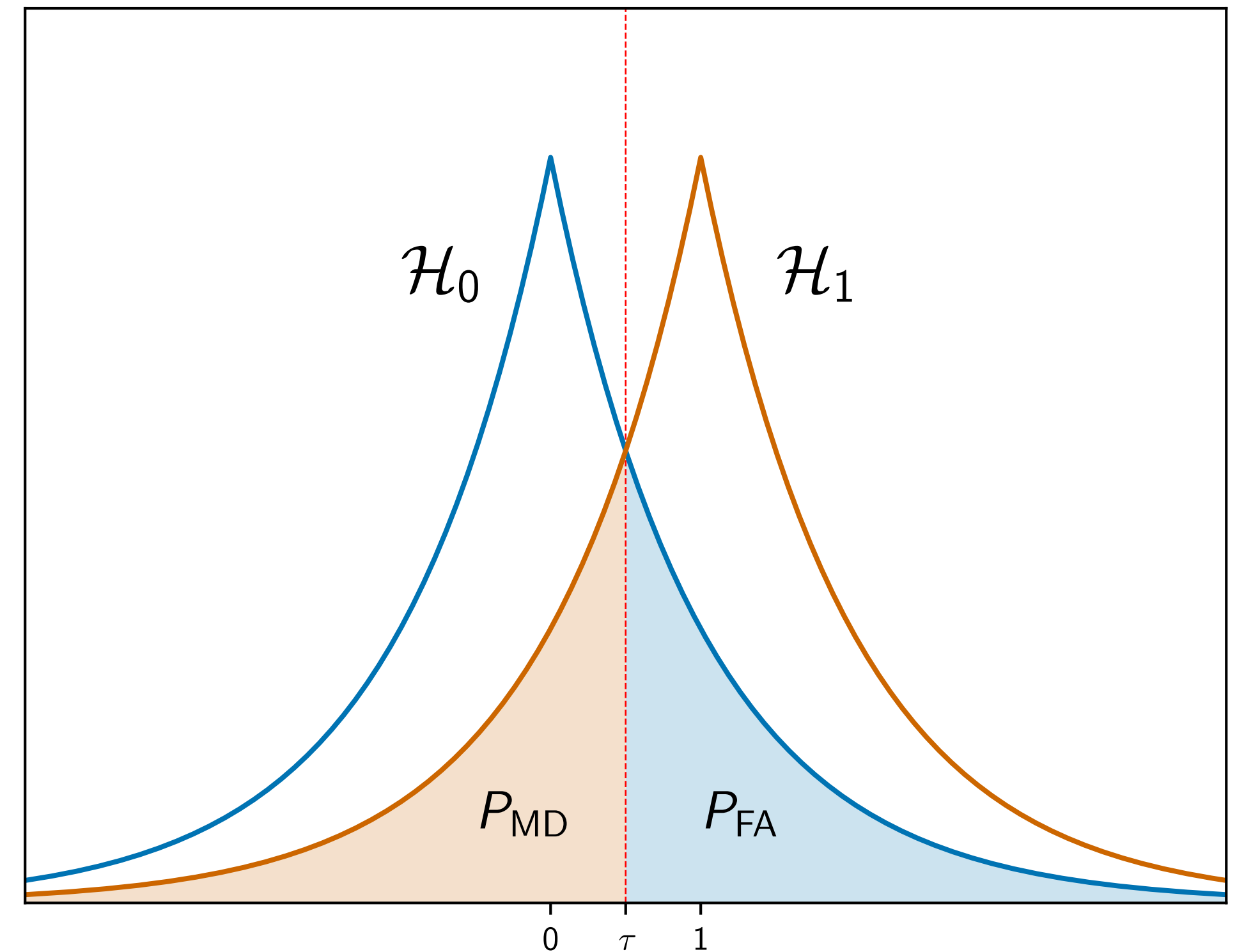


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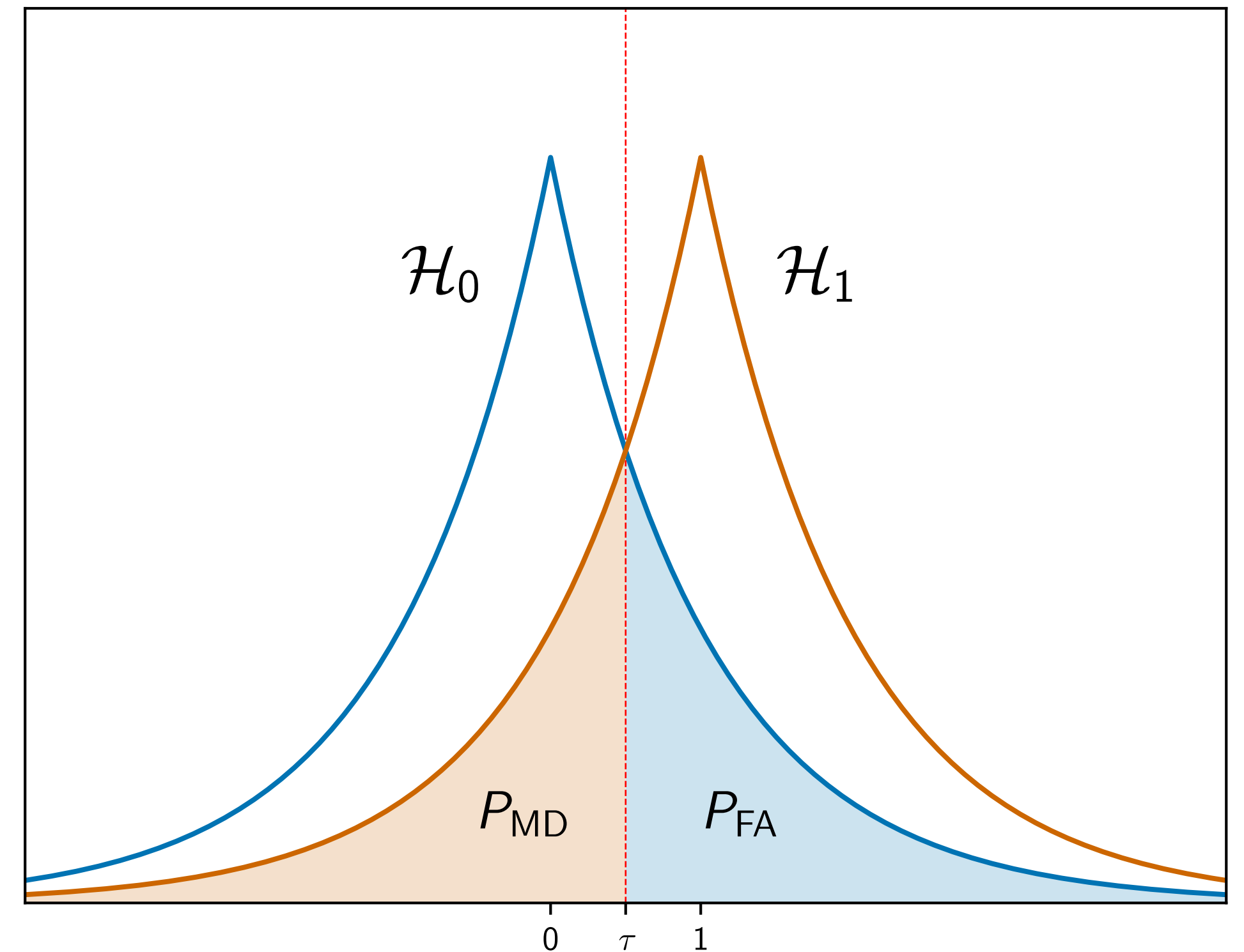
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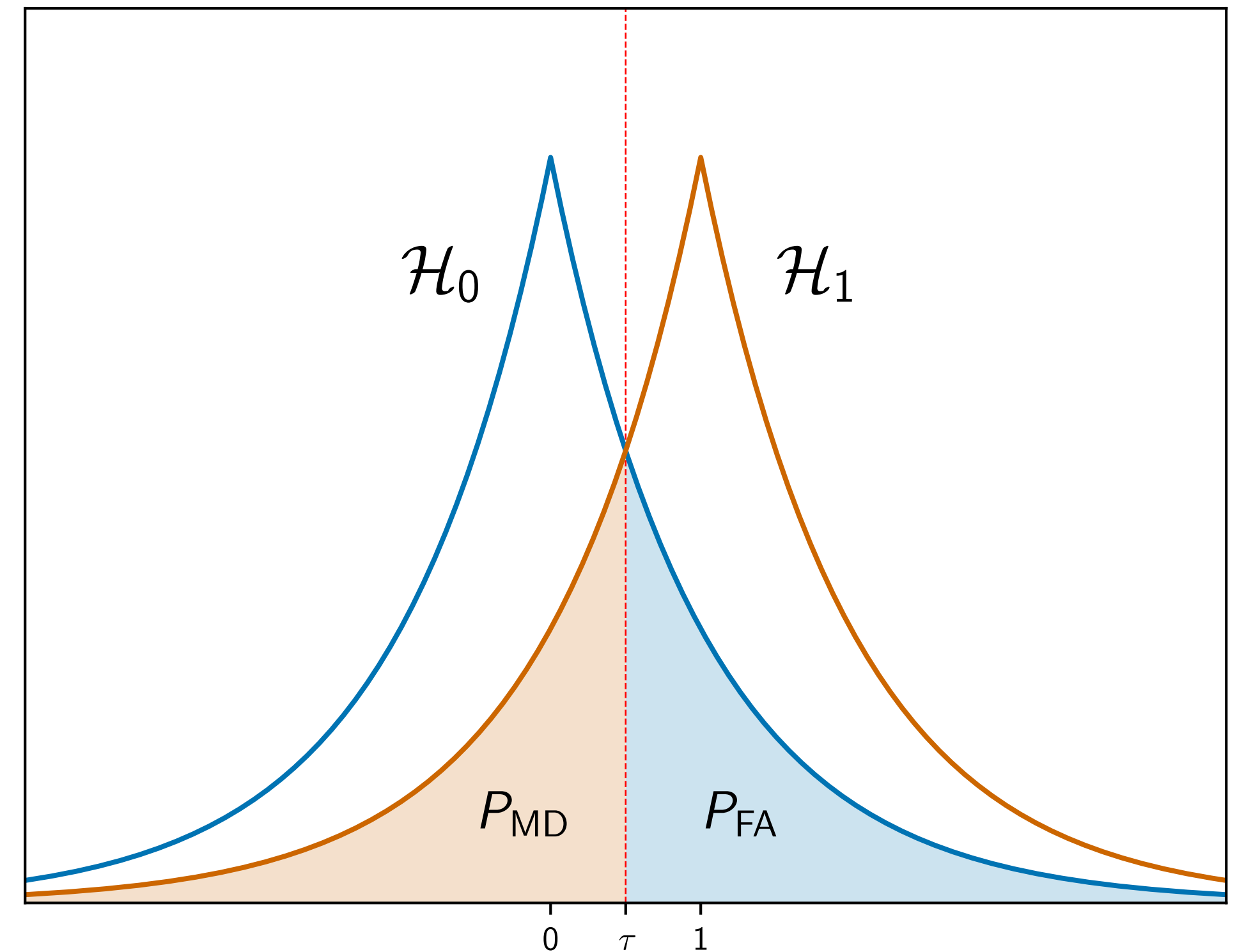
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I hate Laplace noise!

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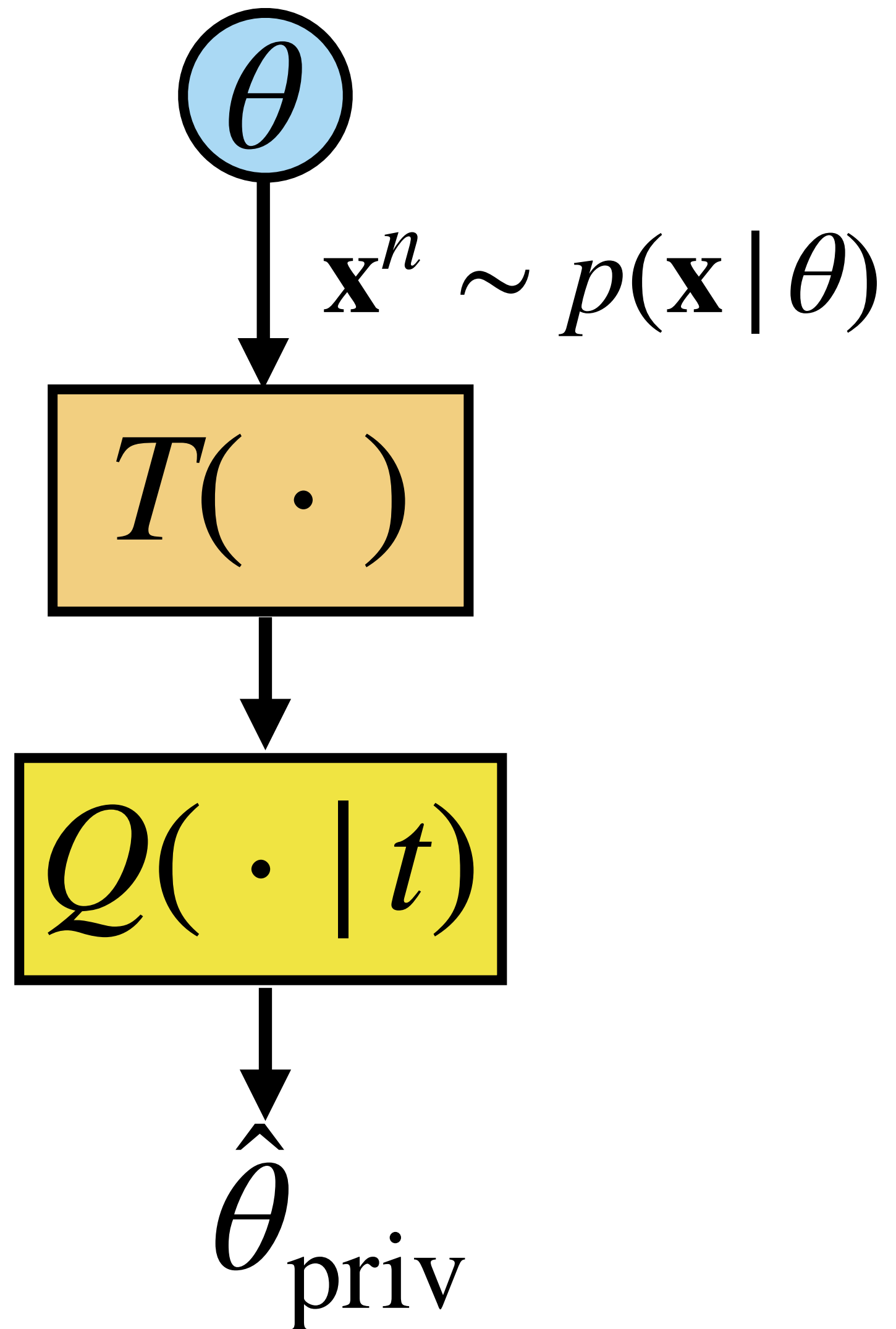
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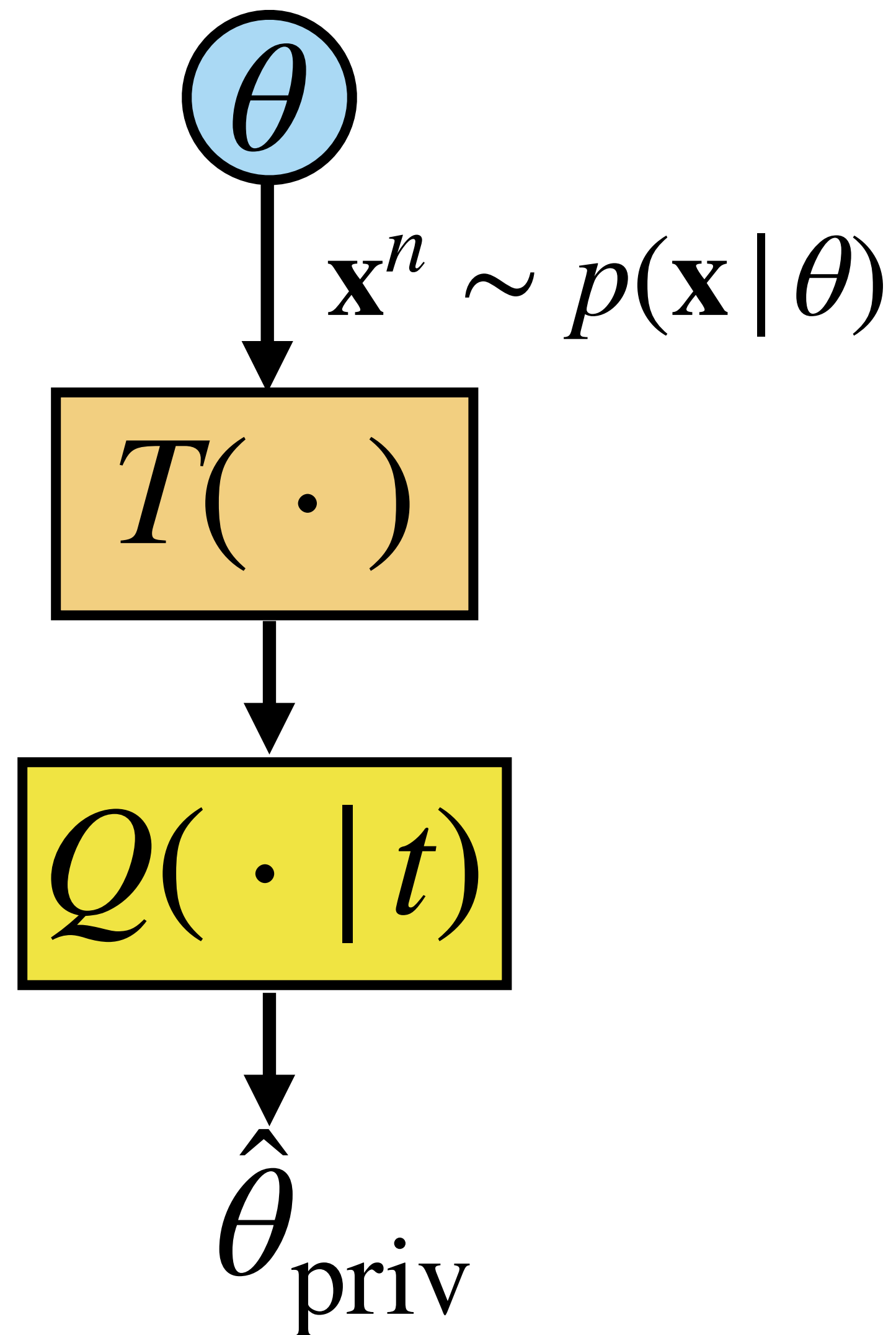
# Point estimation with differential privacy

Adding noise to sufficient statistics



# Point estimation with differential privacy

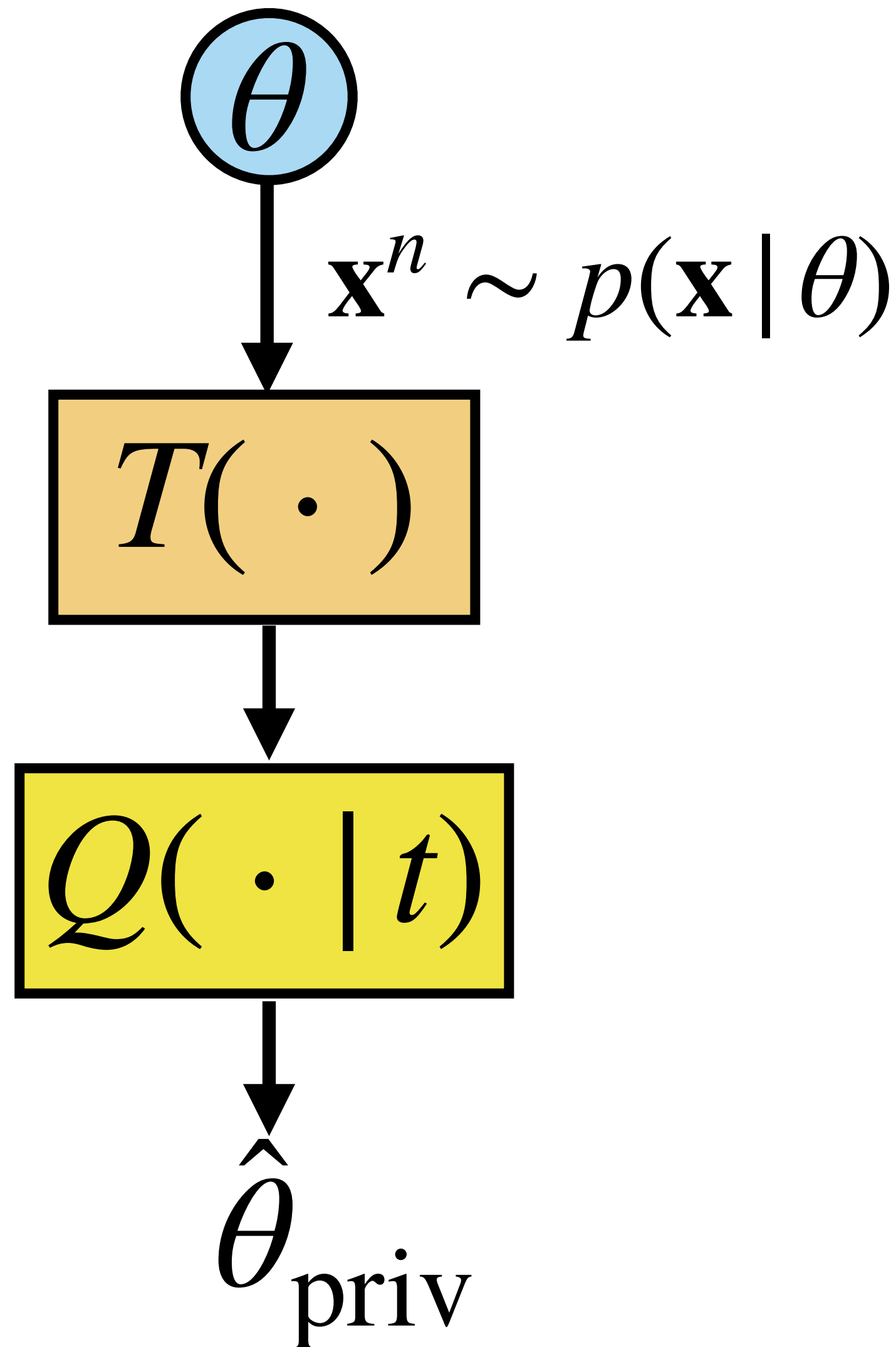
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(Smith 2009):

# Point estimation with differential privacy

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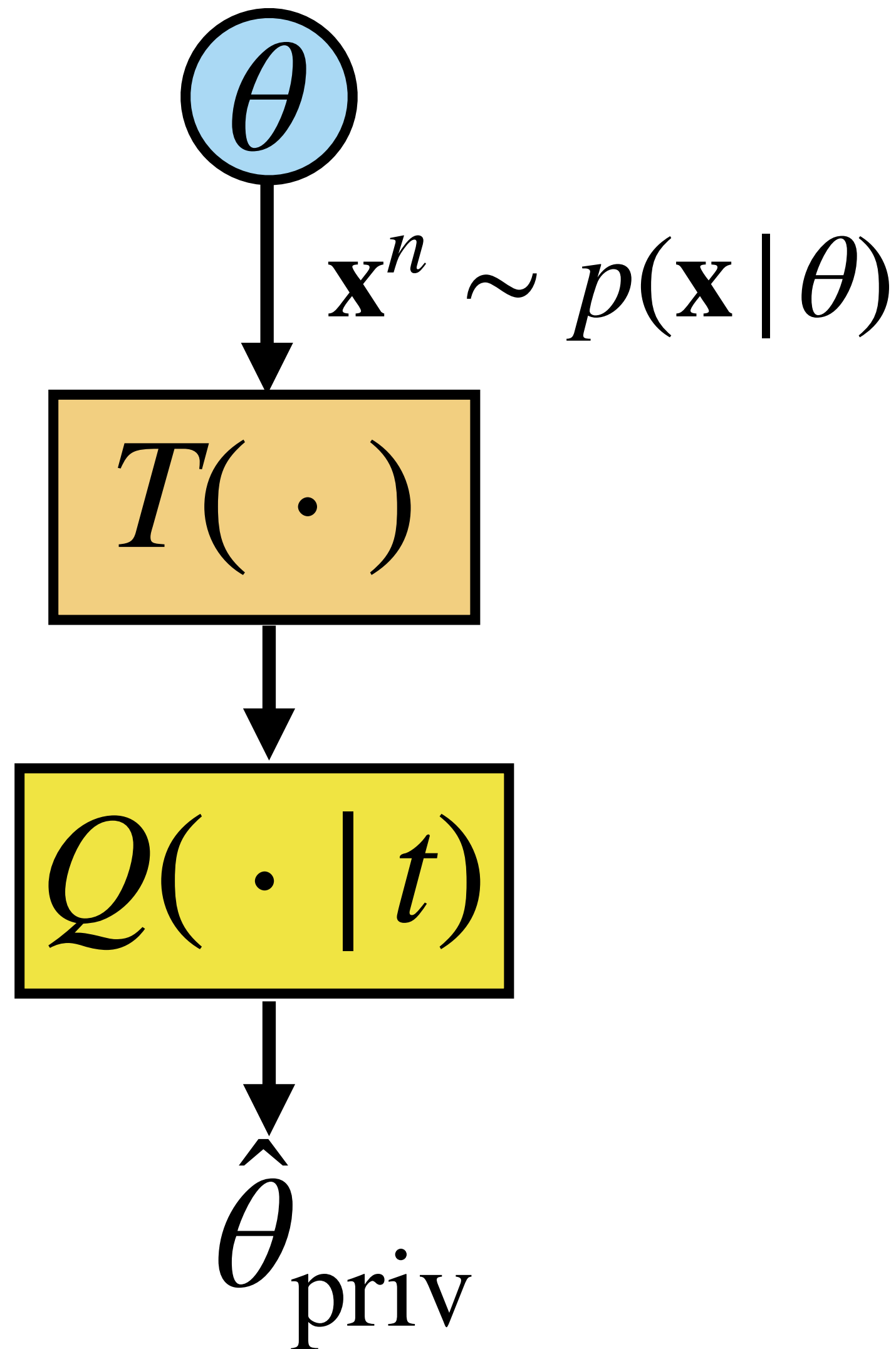


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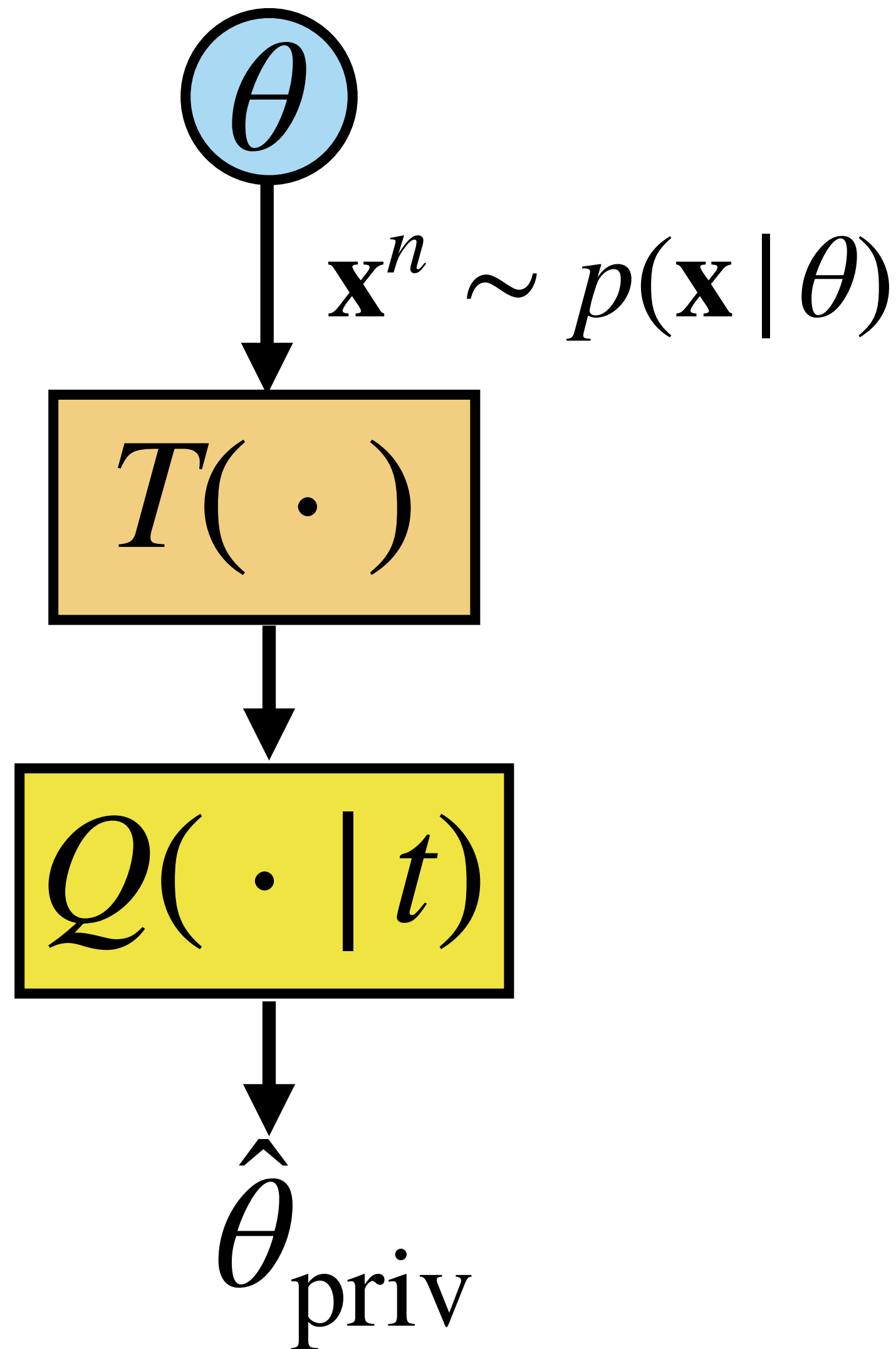
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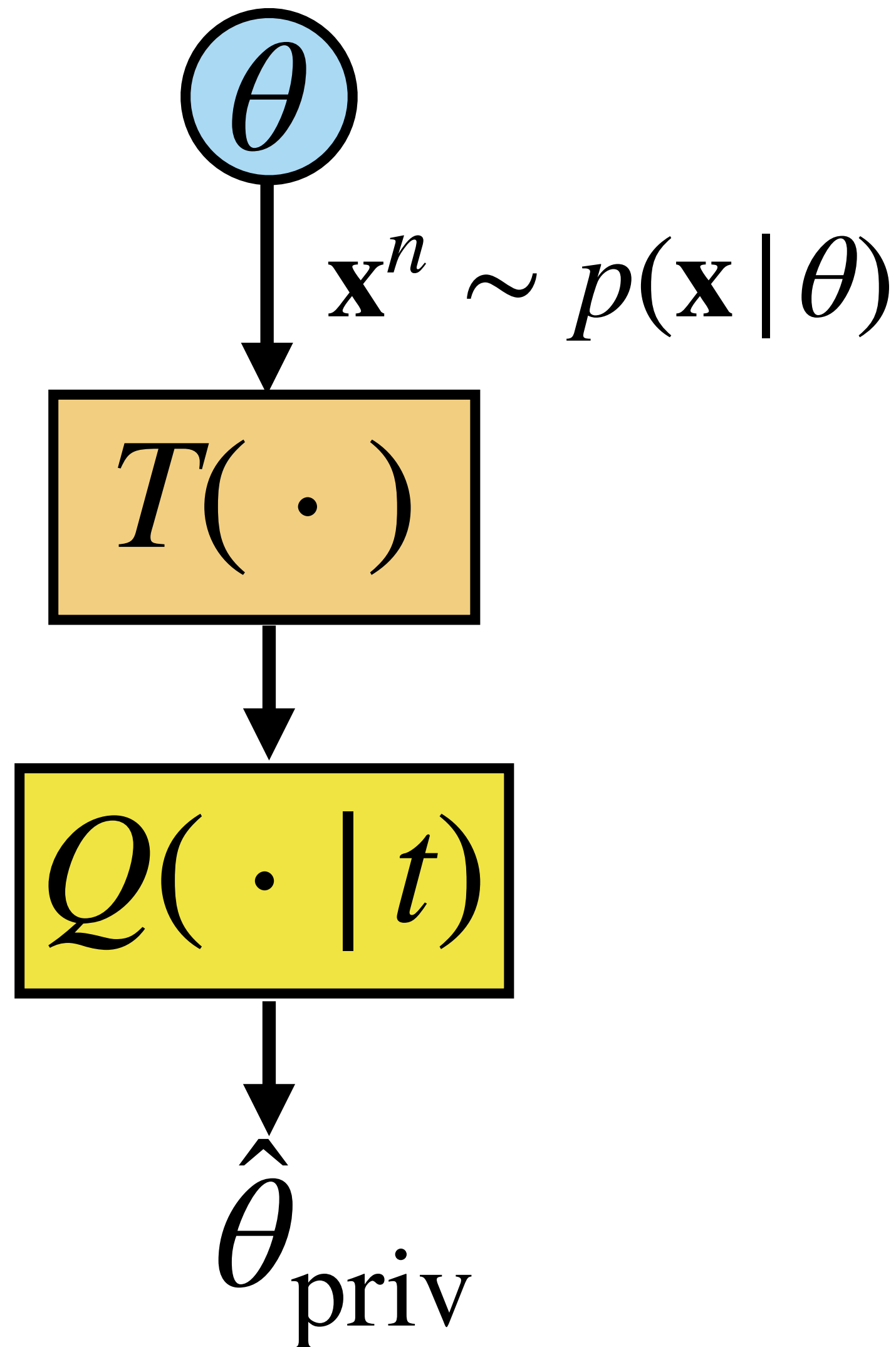


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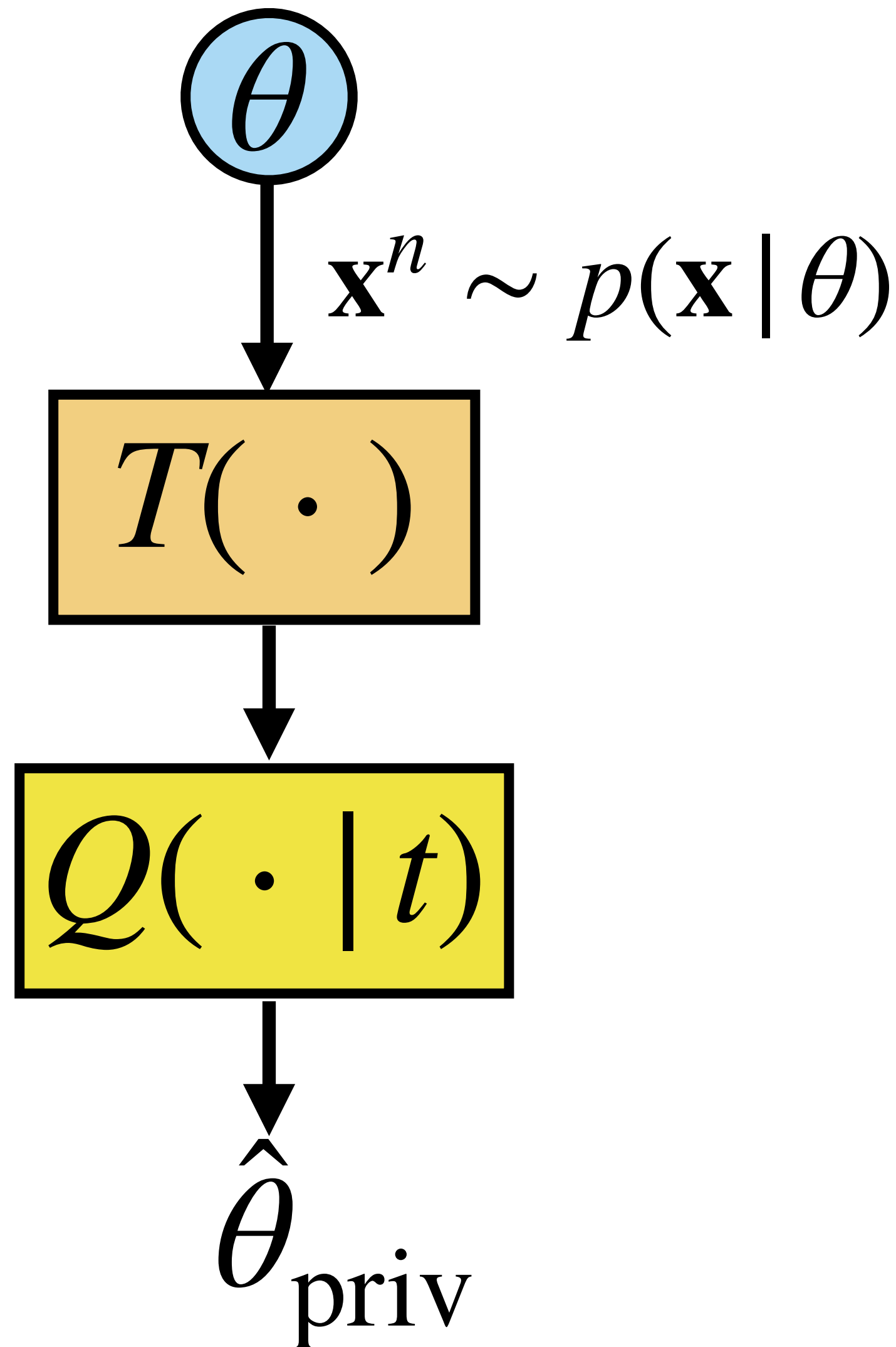


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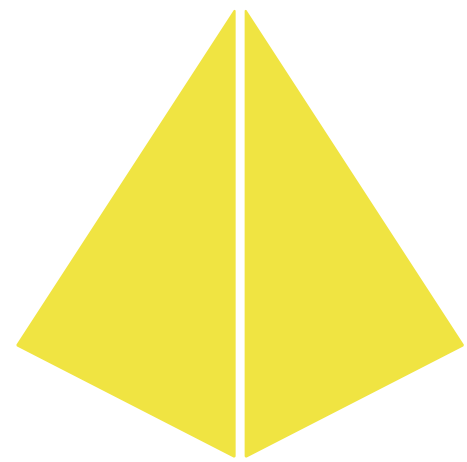
**All we need to know is the sensitivity of  $T(\cdot)$ .**

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## Variations on geometric noise

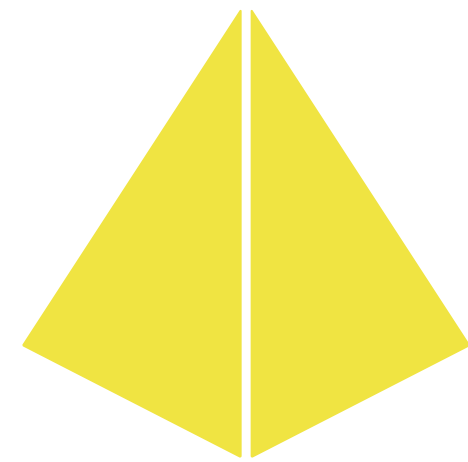
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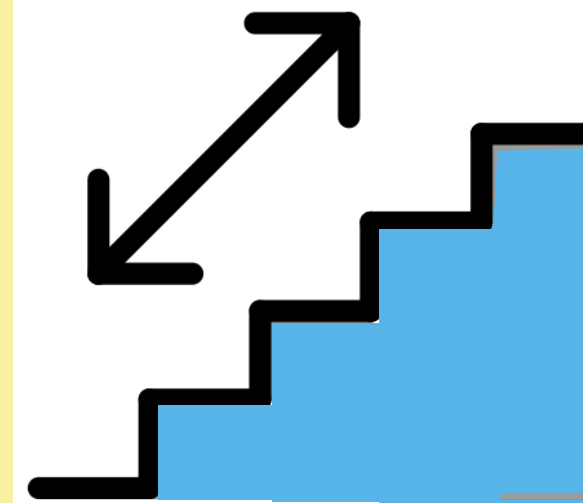


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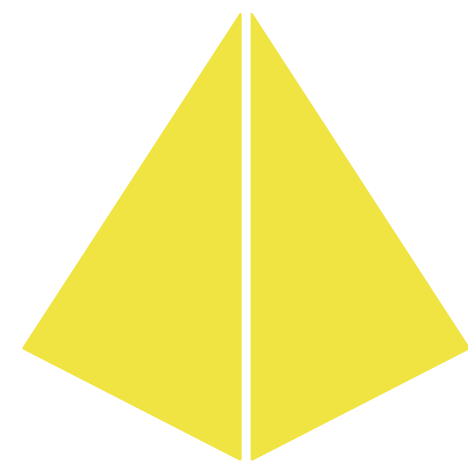
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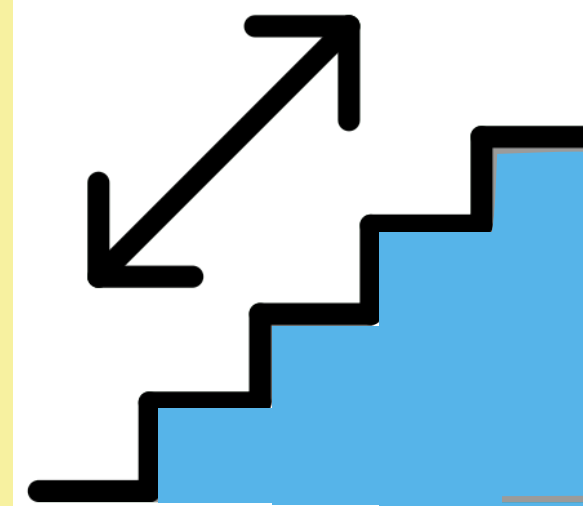


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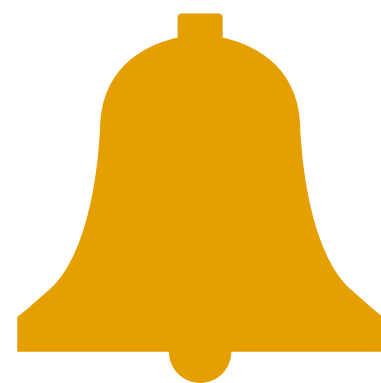
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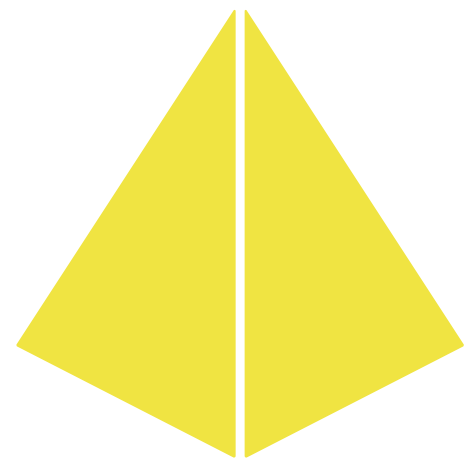
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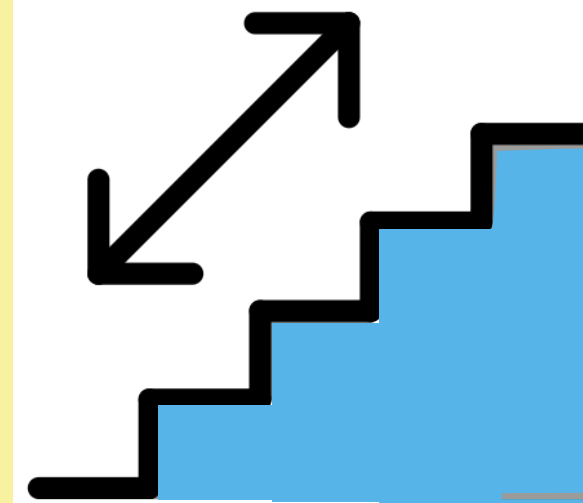


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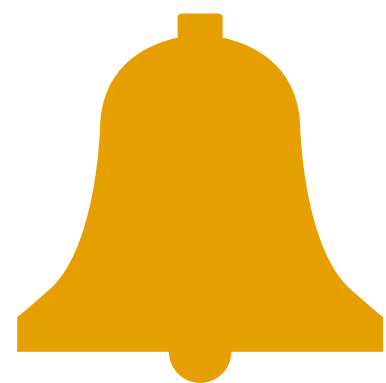
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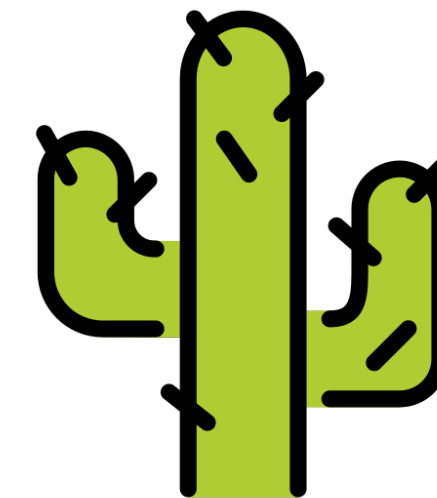
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## "Other"

Geng, Ding, Guo, Kumar (2019/2020)

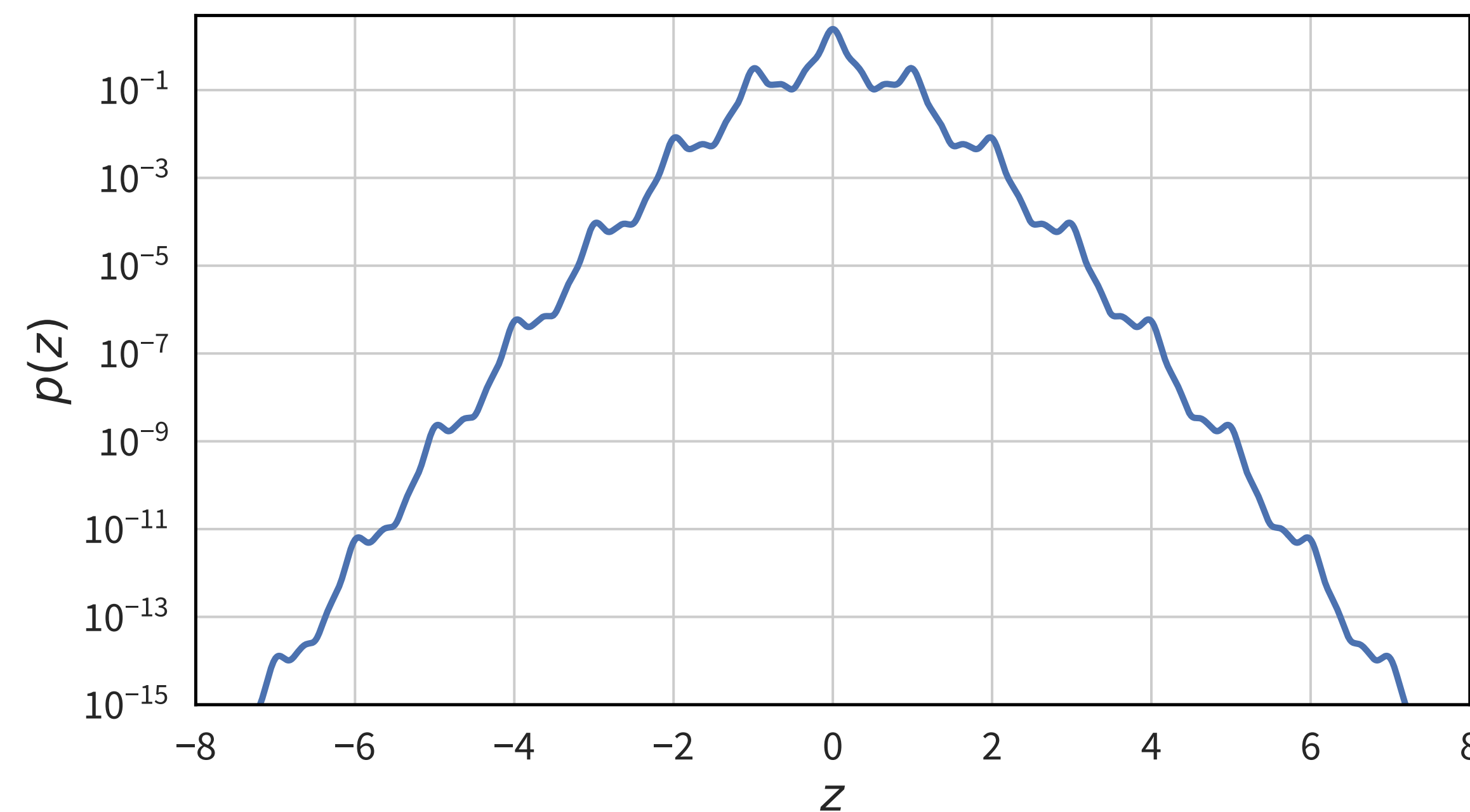
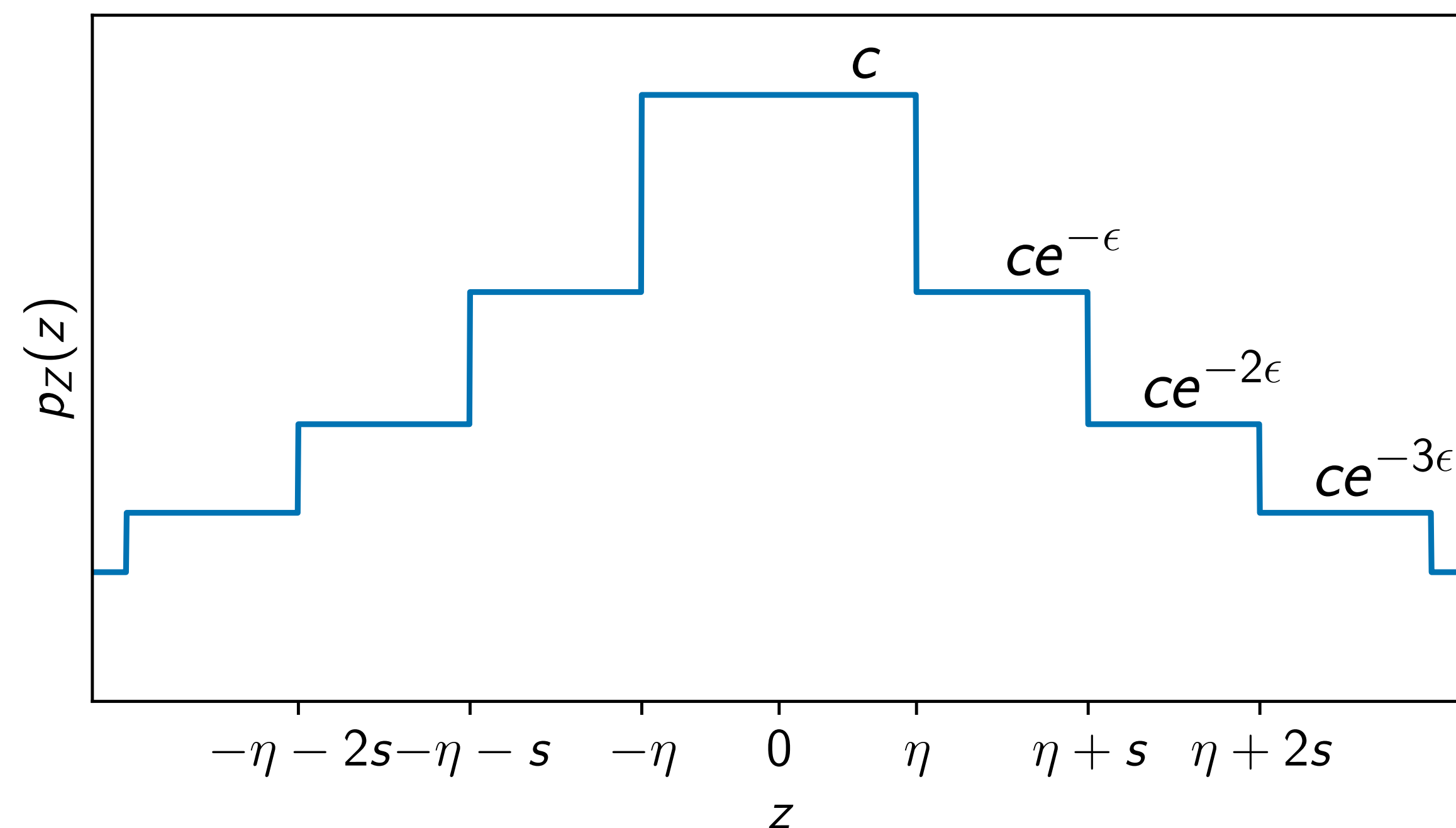
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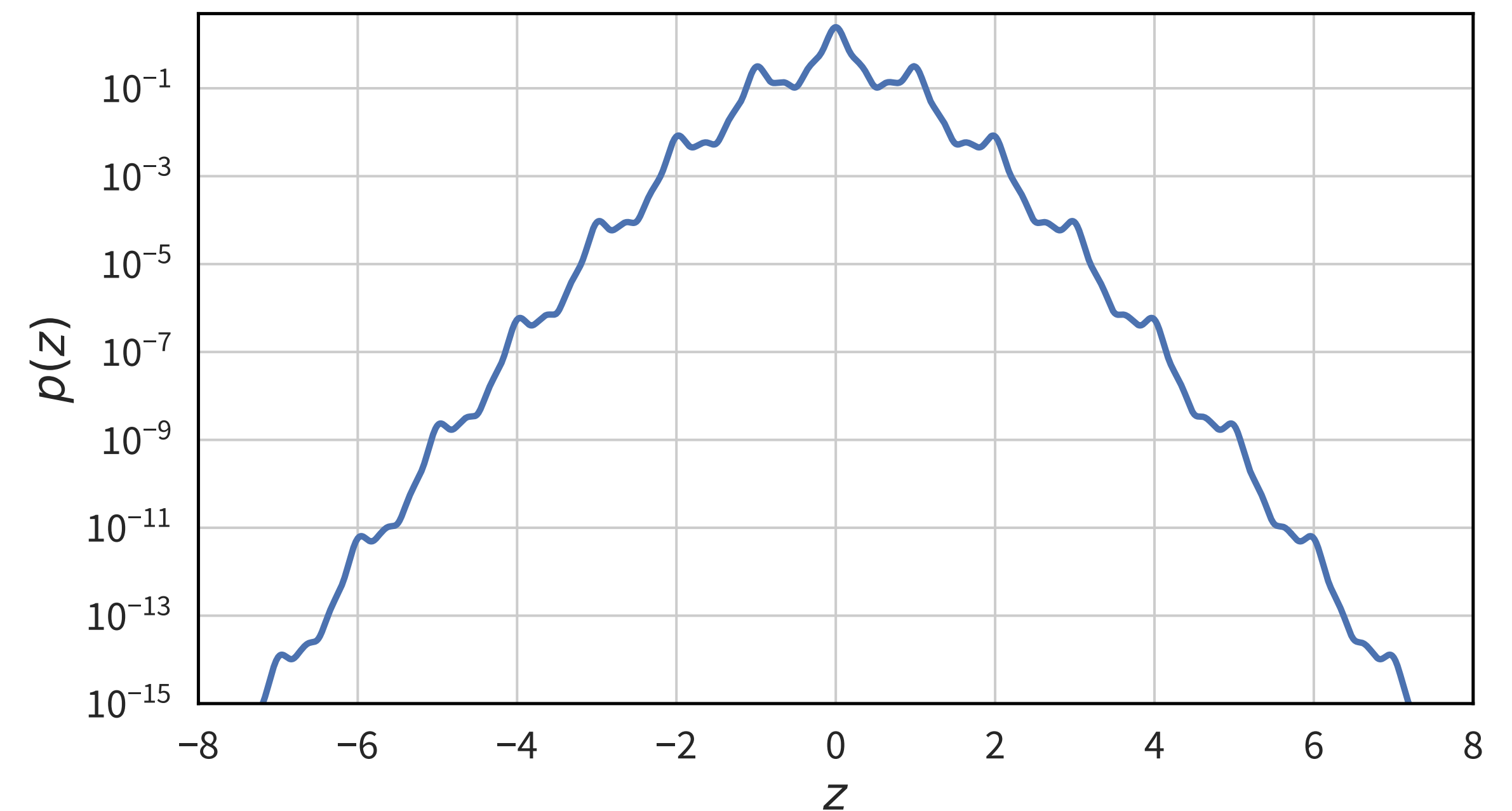
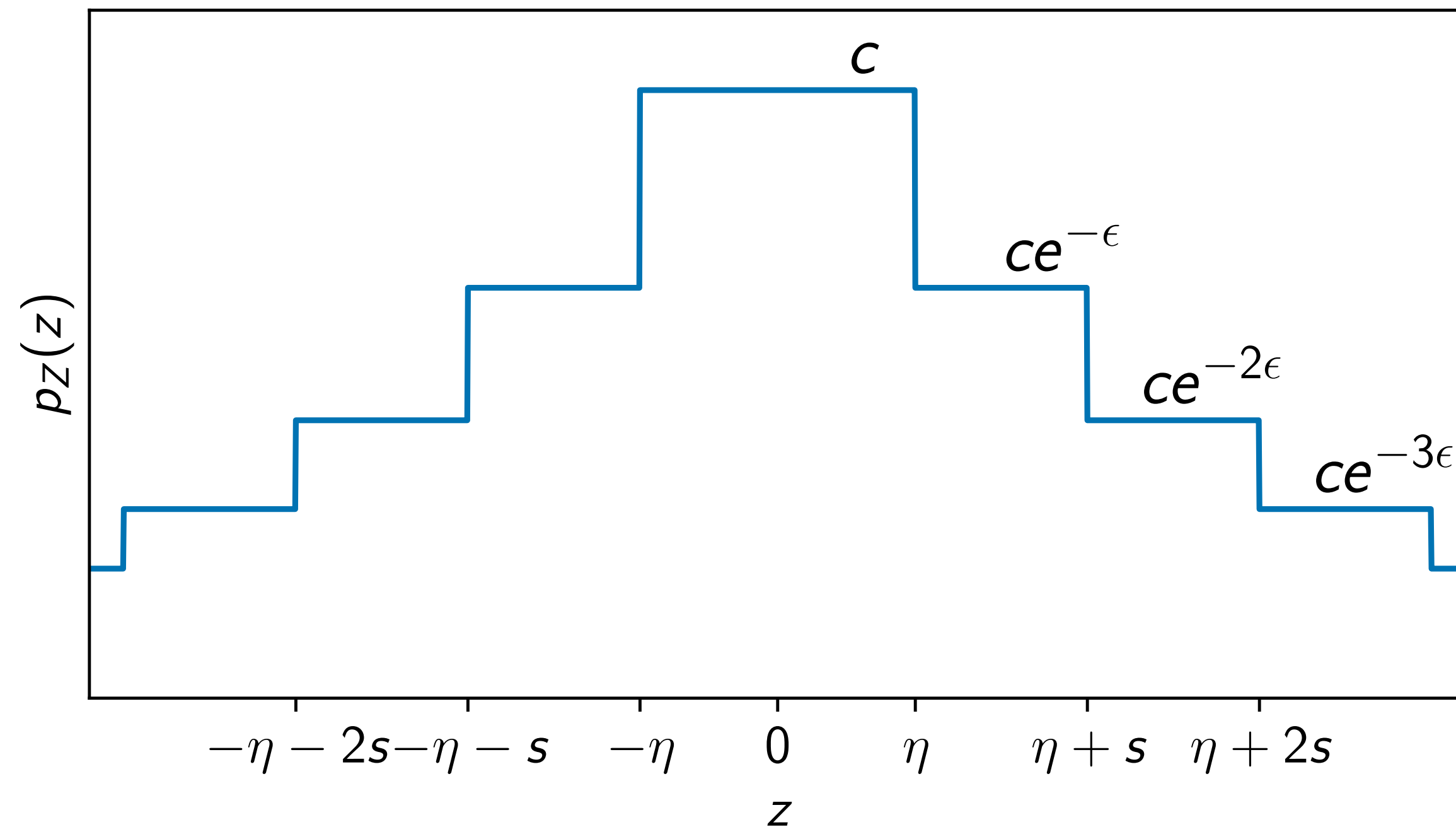
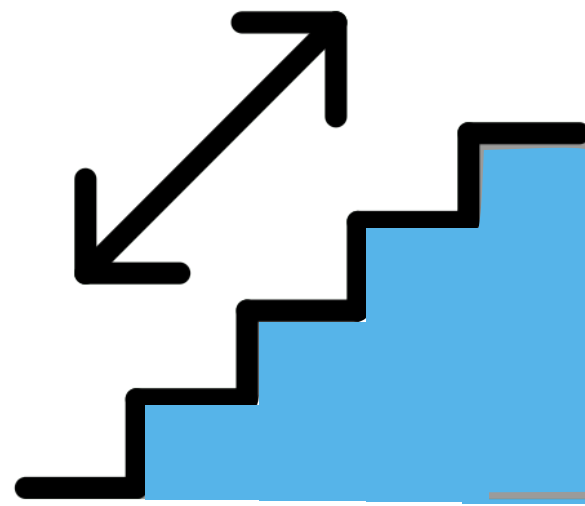
# “Optimal” noise distributions

Beyond Gaussian and Laplace



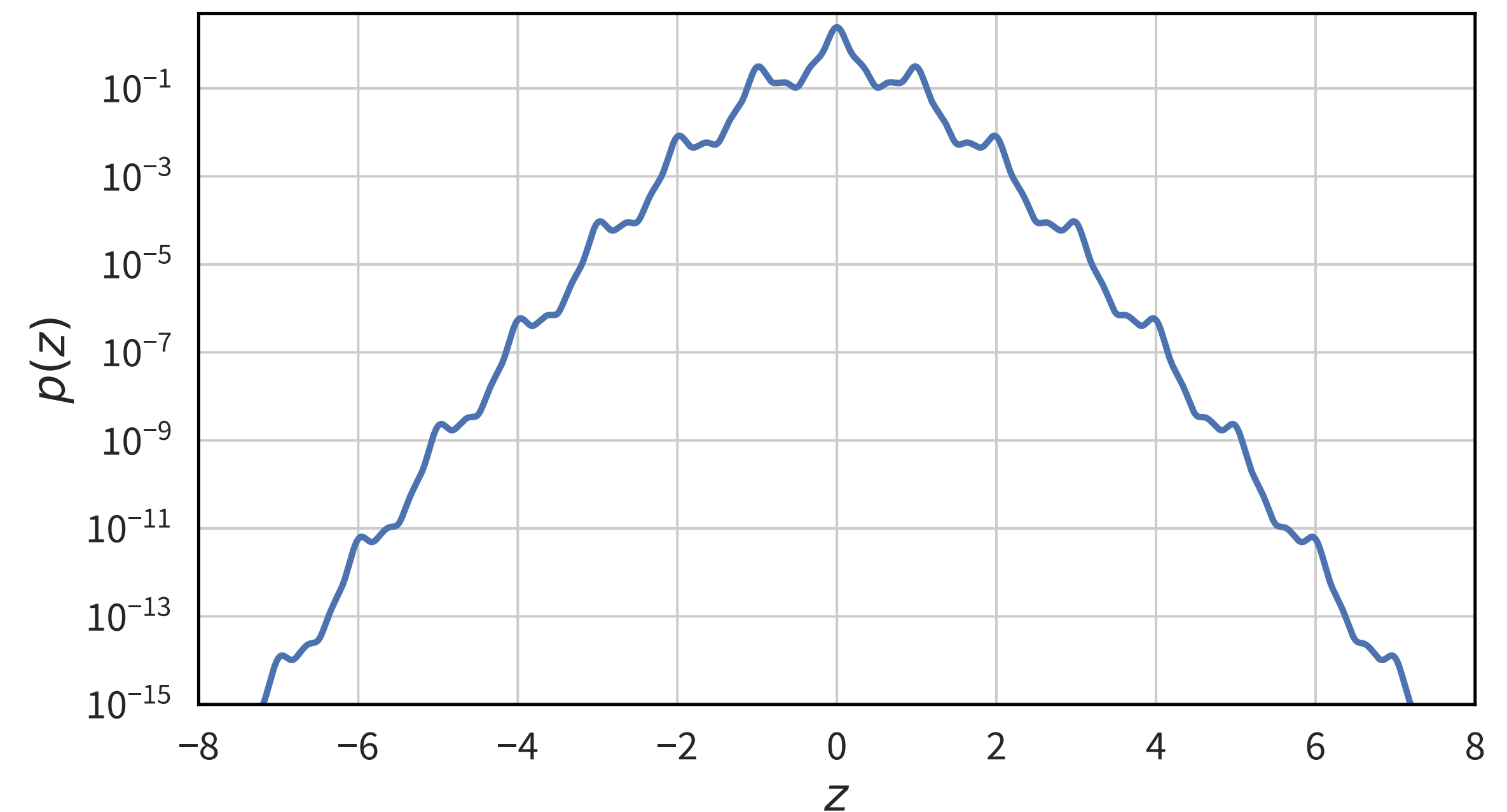
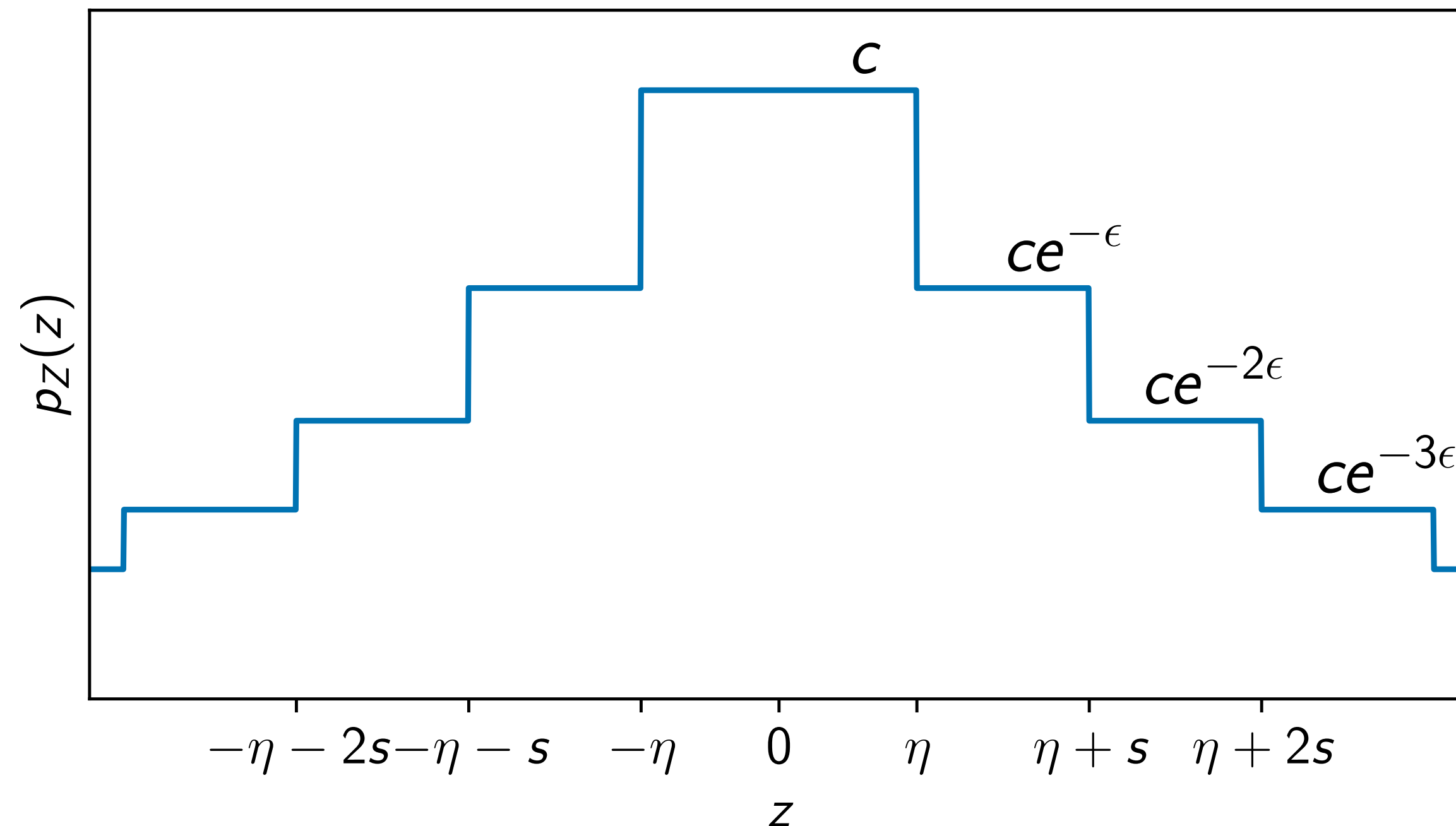
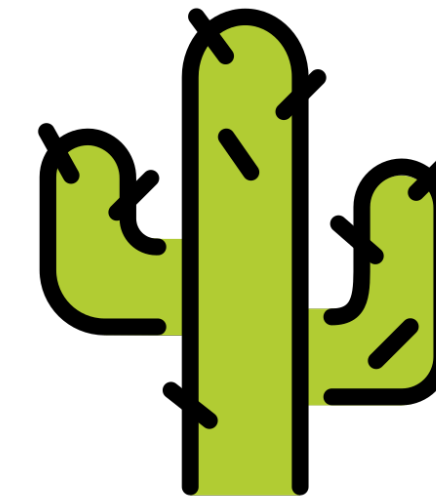
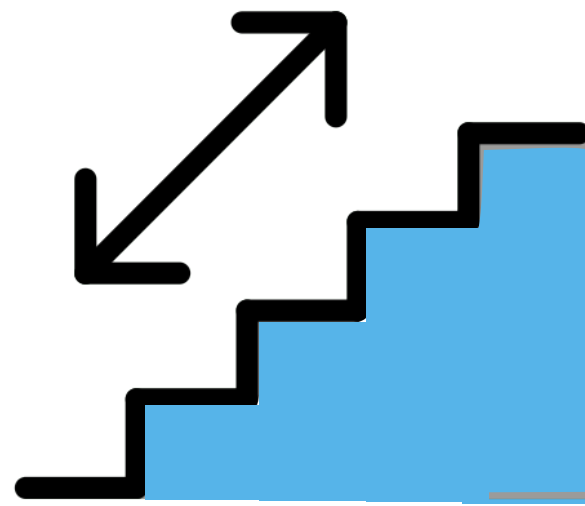
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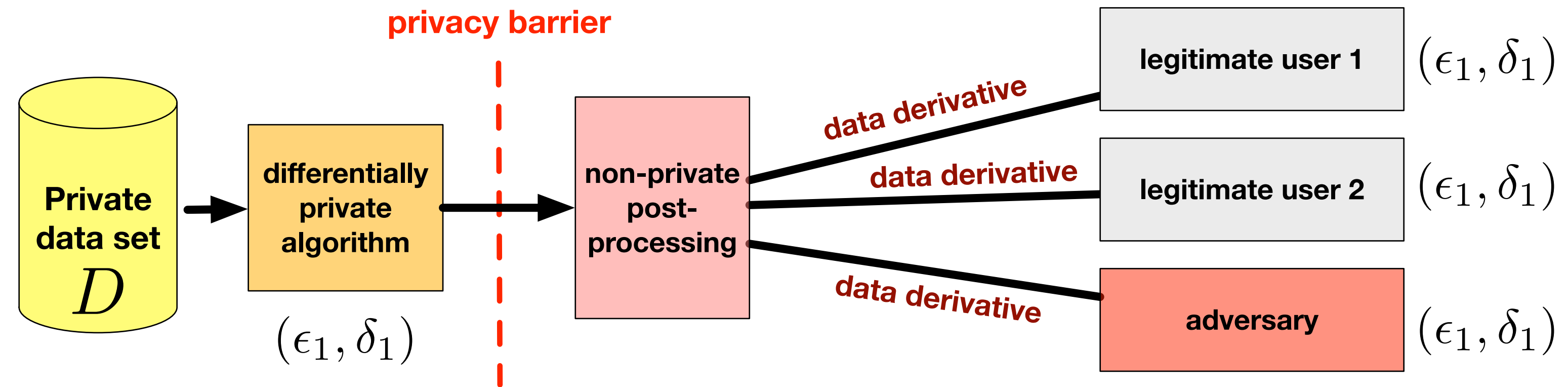
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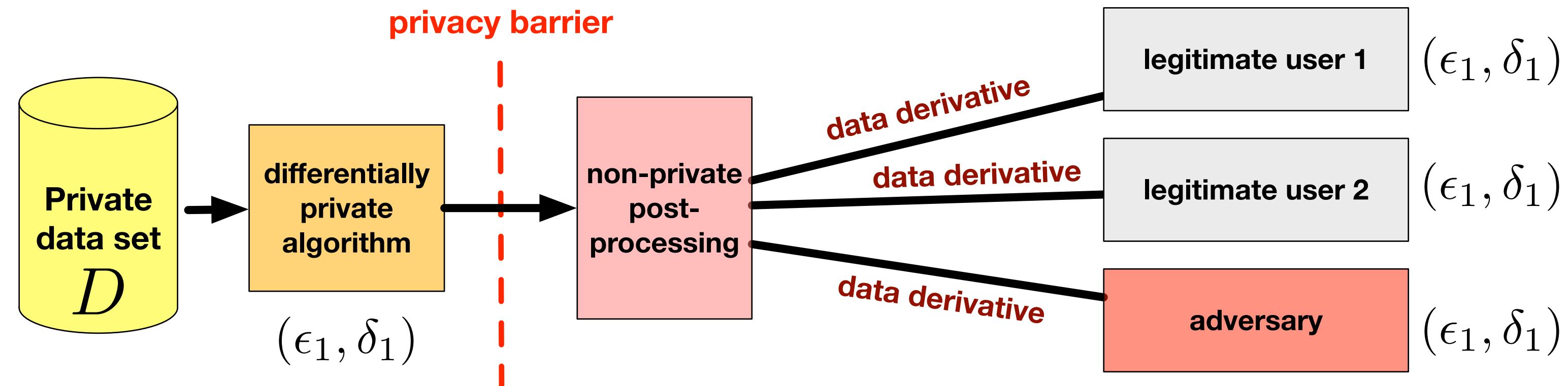
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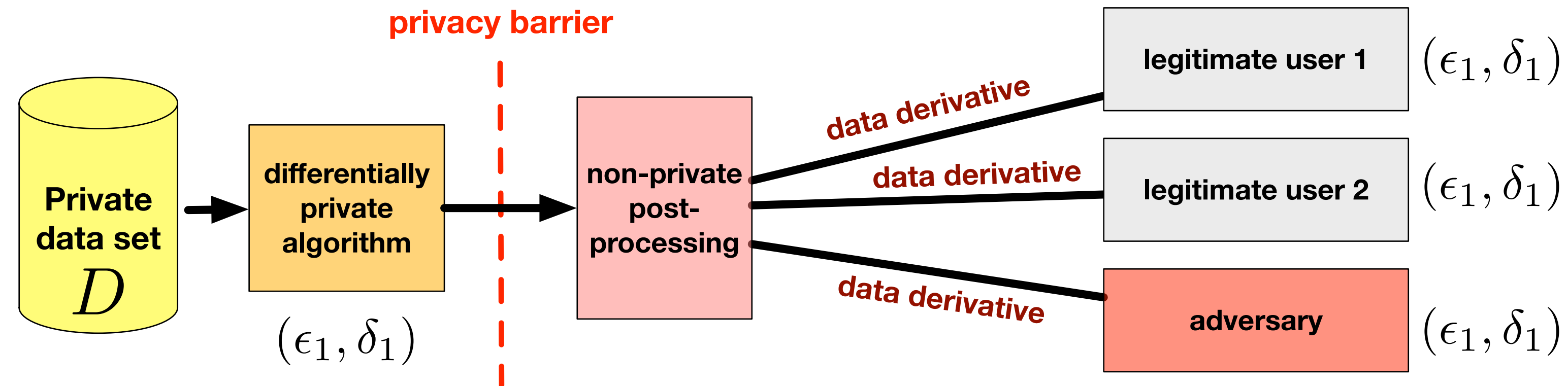
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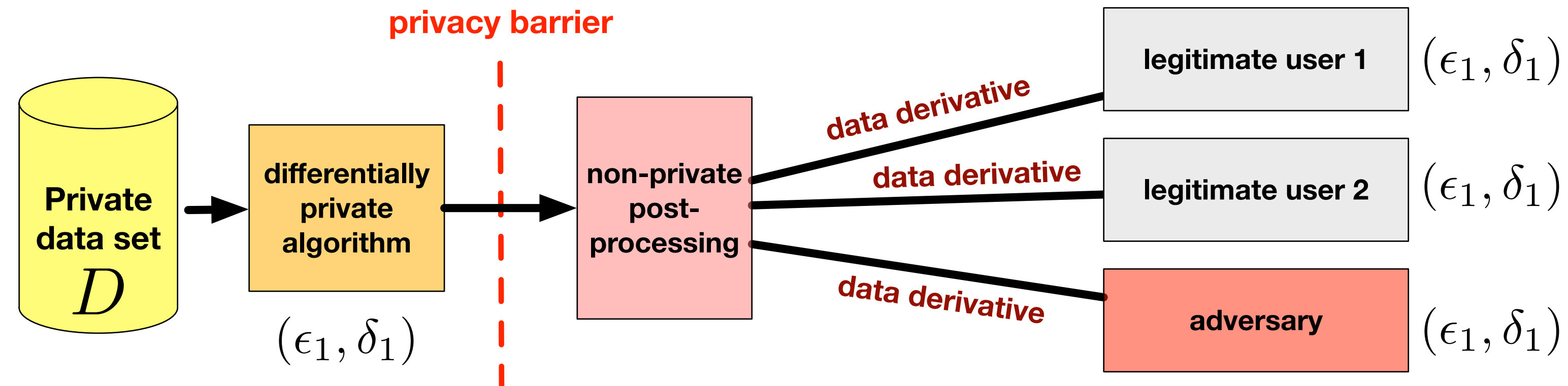


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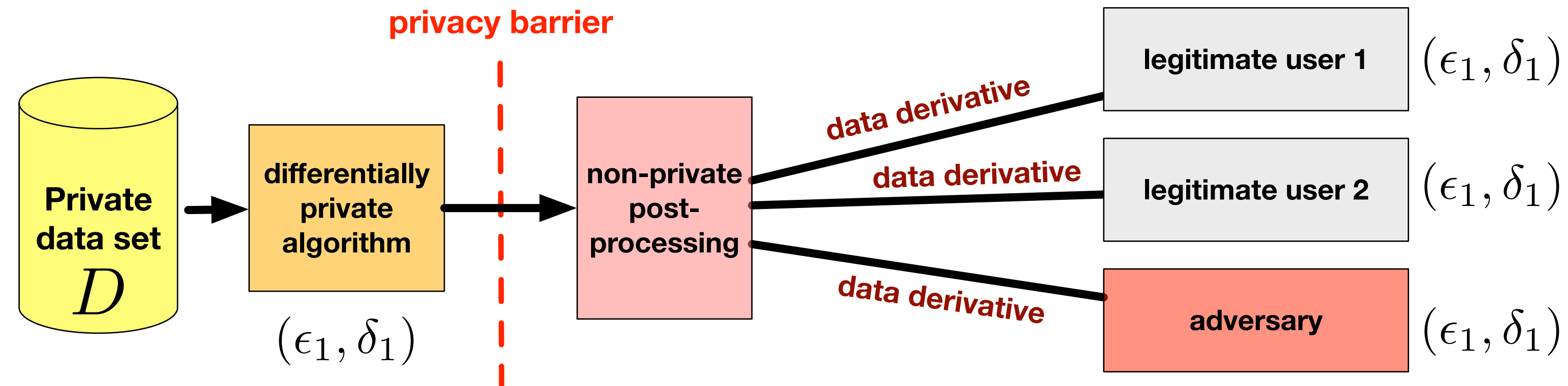
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Umezawa in Sagami  
Province

相州梅沢庄

Soshū Umezawanoshō

Vista 3

*f*-divergences/composition



# Privacy loss random variables

**Characterizing the distribution of the LLR**

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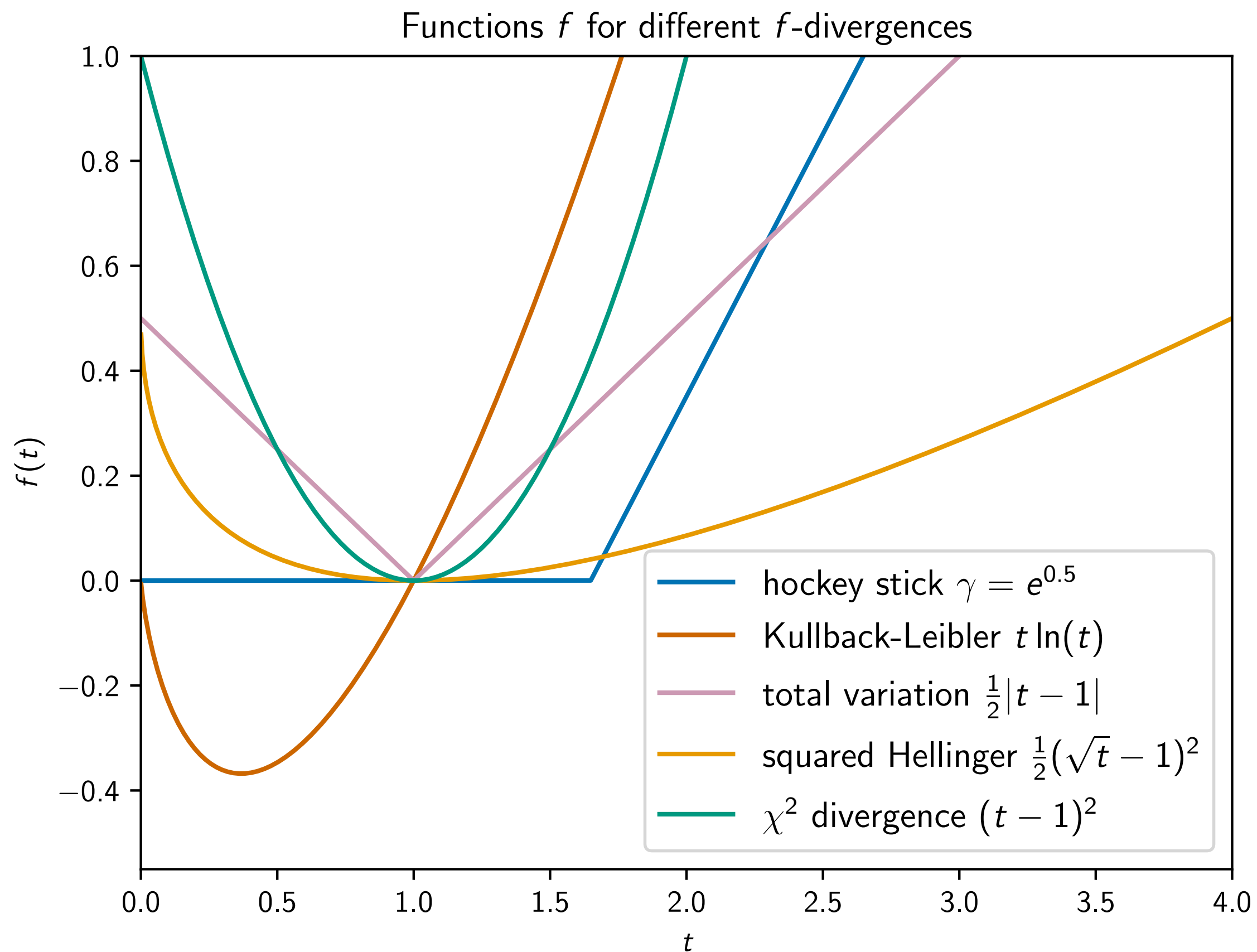
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**A challenge:** this is defined for a single pair of inputs  $(x, x')$ . We would like to only deal with the “worst case” pair of inputs.

# Generalized divergences and the 🏒 divergence

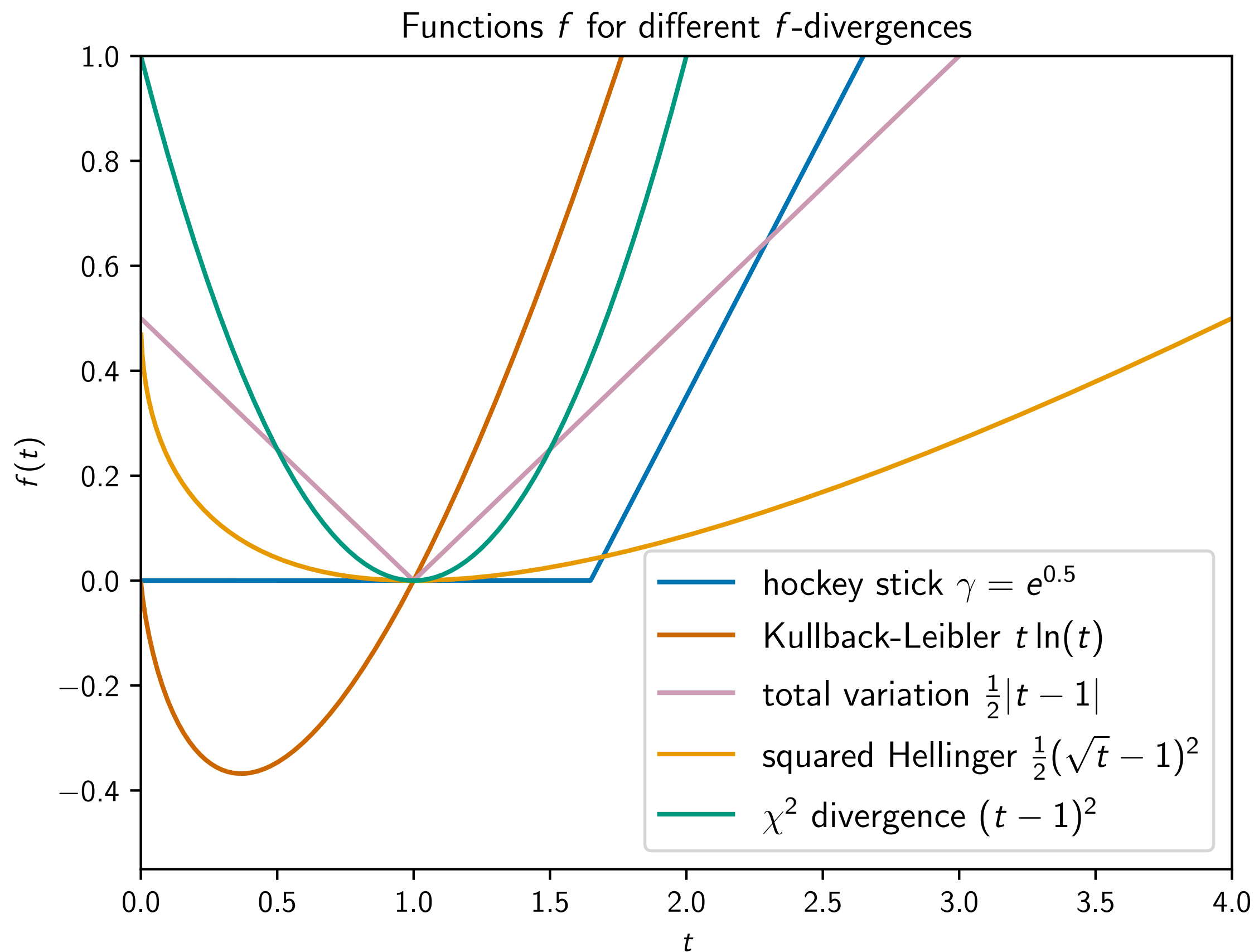
How different are these two distributions?



Rényi (1961), Csiszár (1963), Morimoto (1963), Ali, Silvey (1966), Csiszár (1967), Polyanskiy, Poor, Verdu (2010), Balle, Barthe, Gaboardi, Geumlek (2019)

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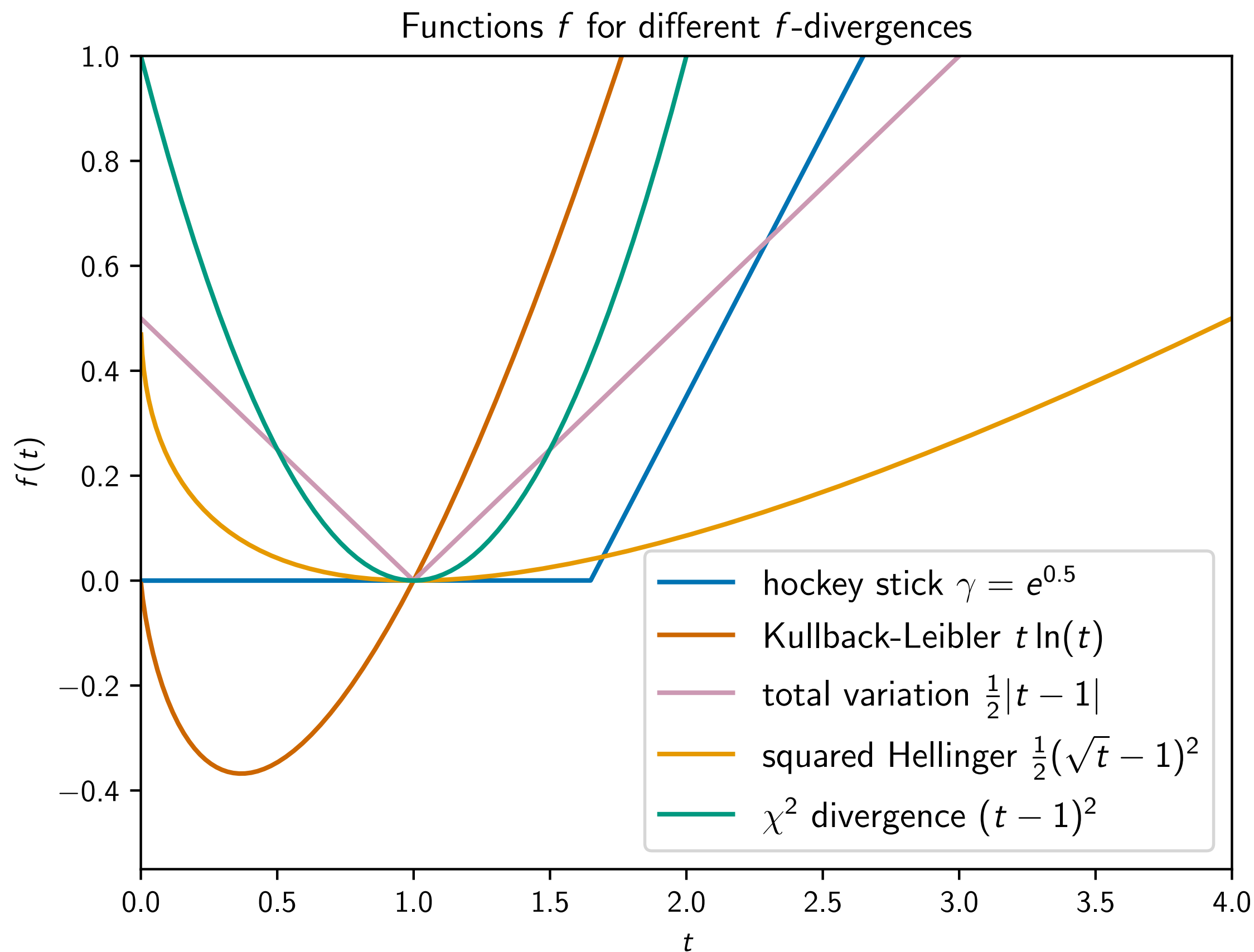


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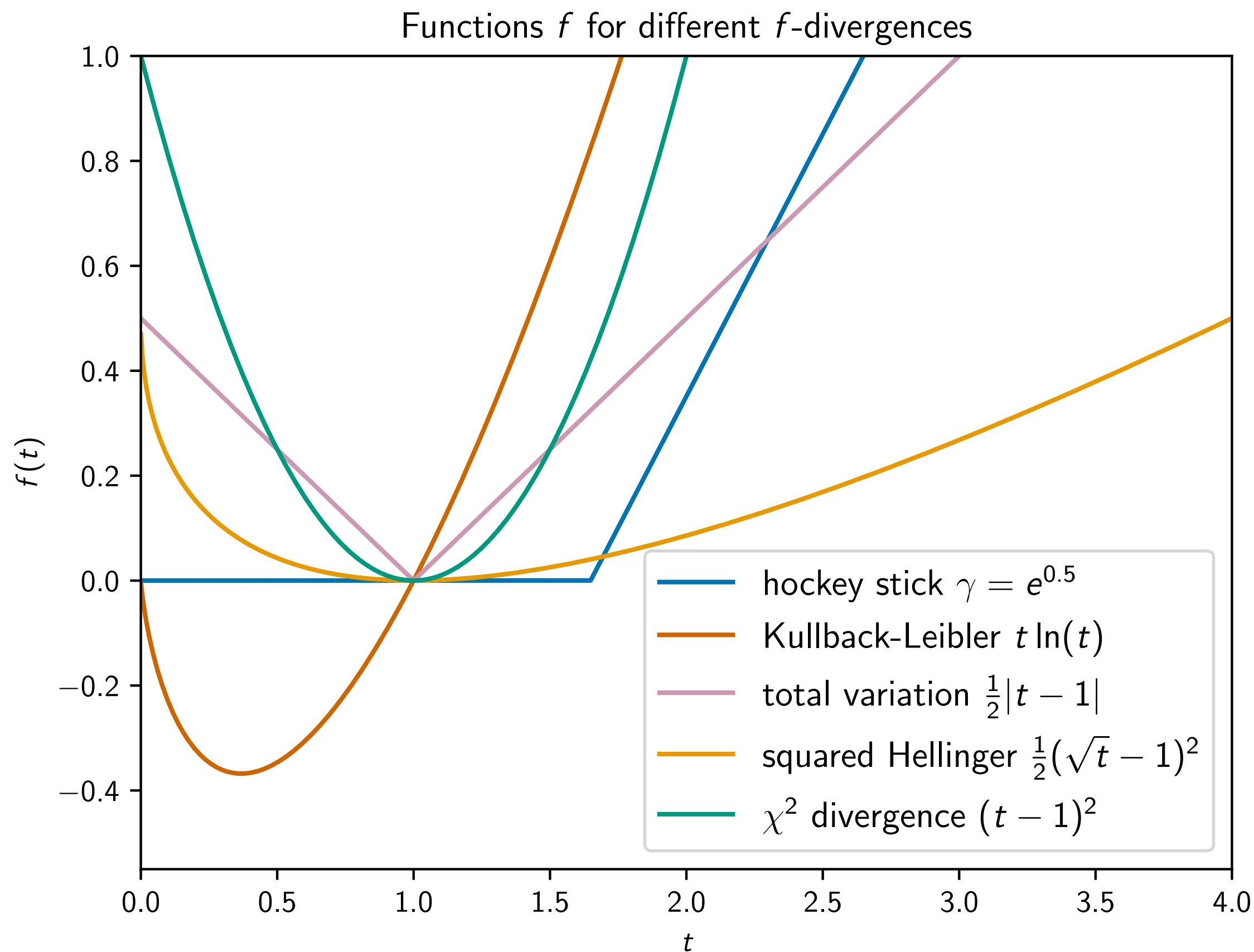
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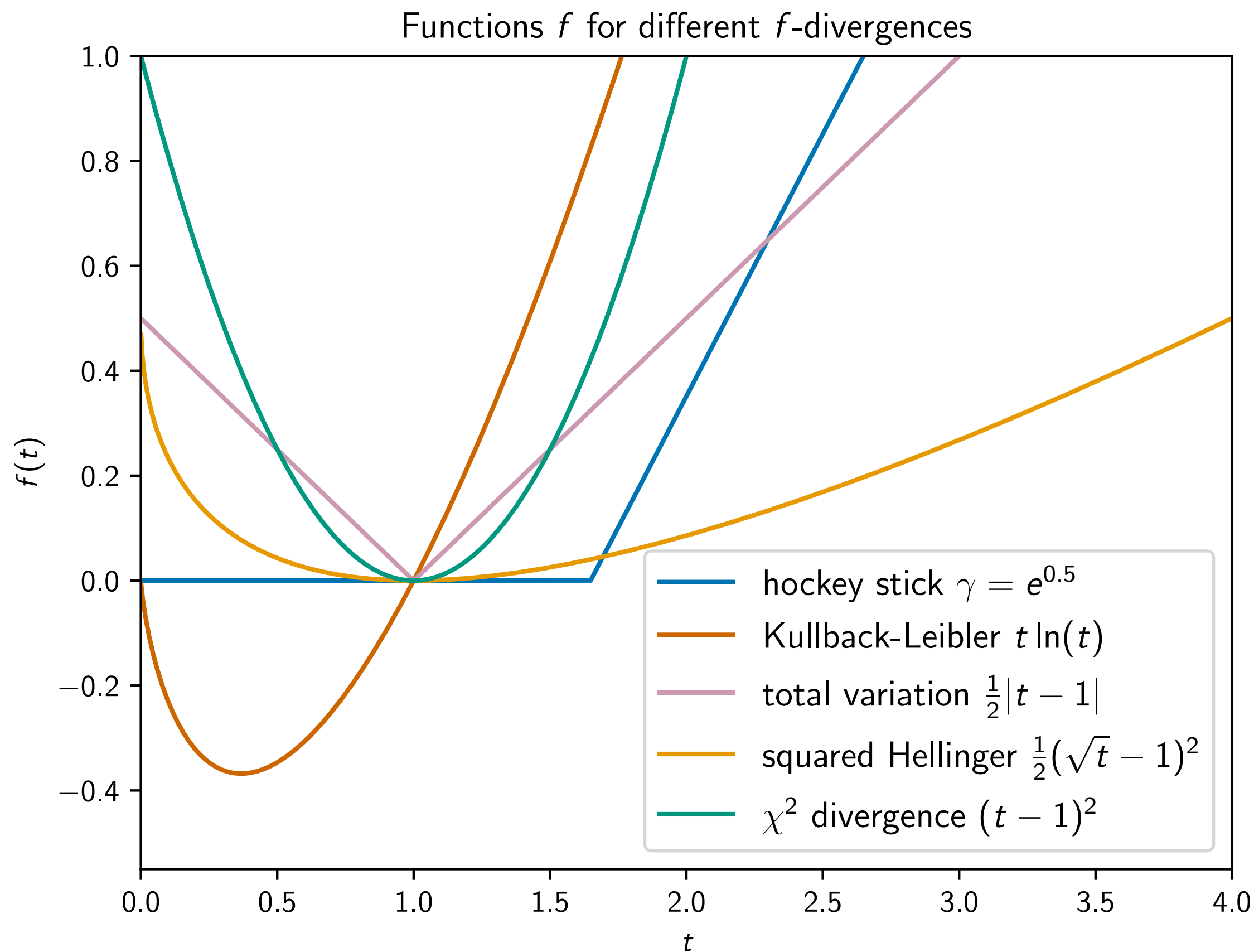
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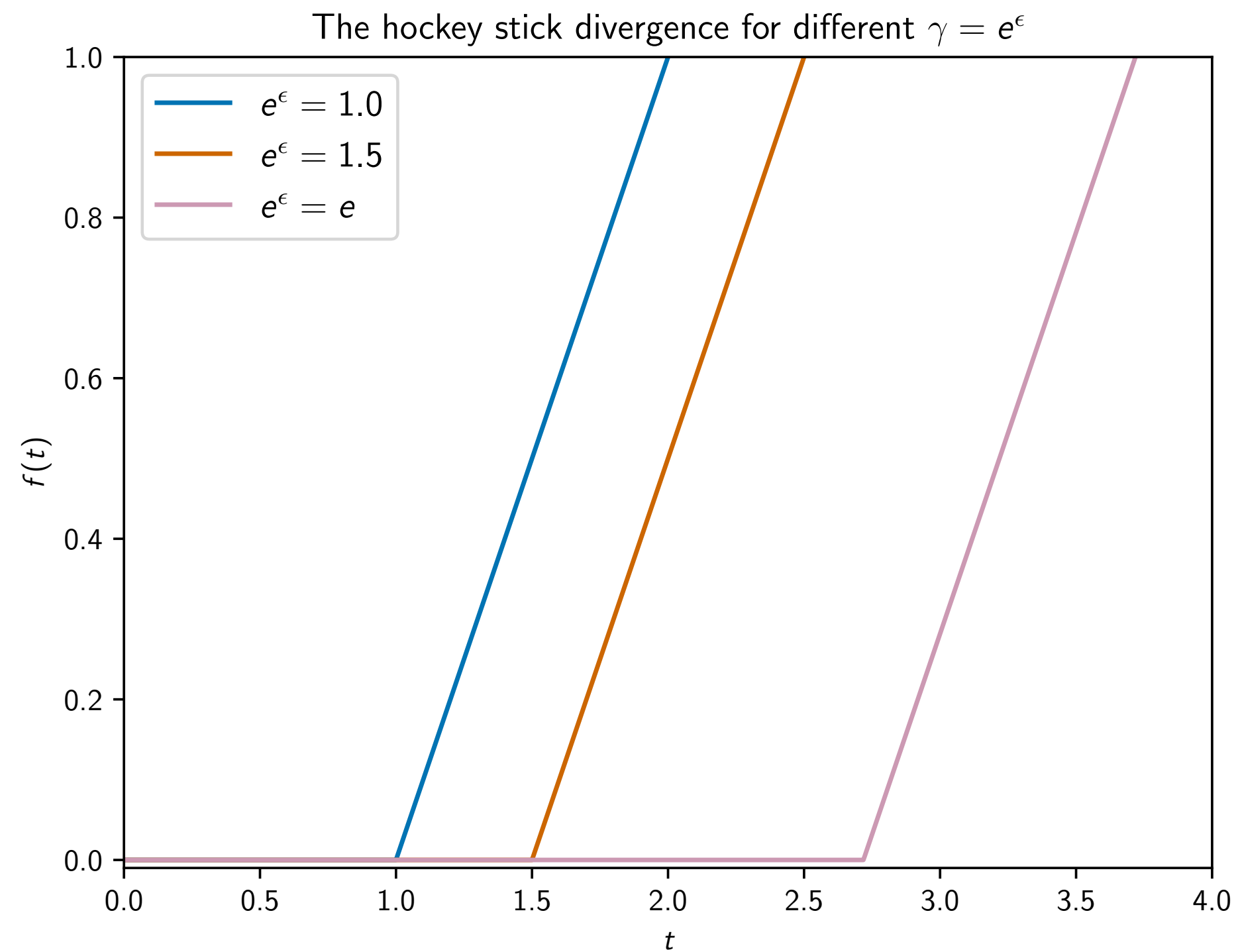
$$E_\gamma(\mu \parallel \nu) = \int_\Omega \left( \frac{d\mu}{d\nu} - \gamma \right)^+ d\nu = \sup_A [\mu(A) - \gamma \nu(A)]$$

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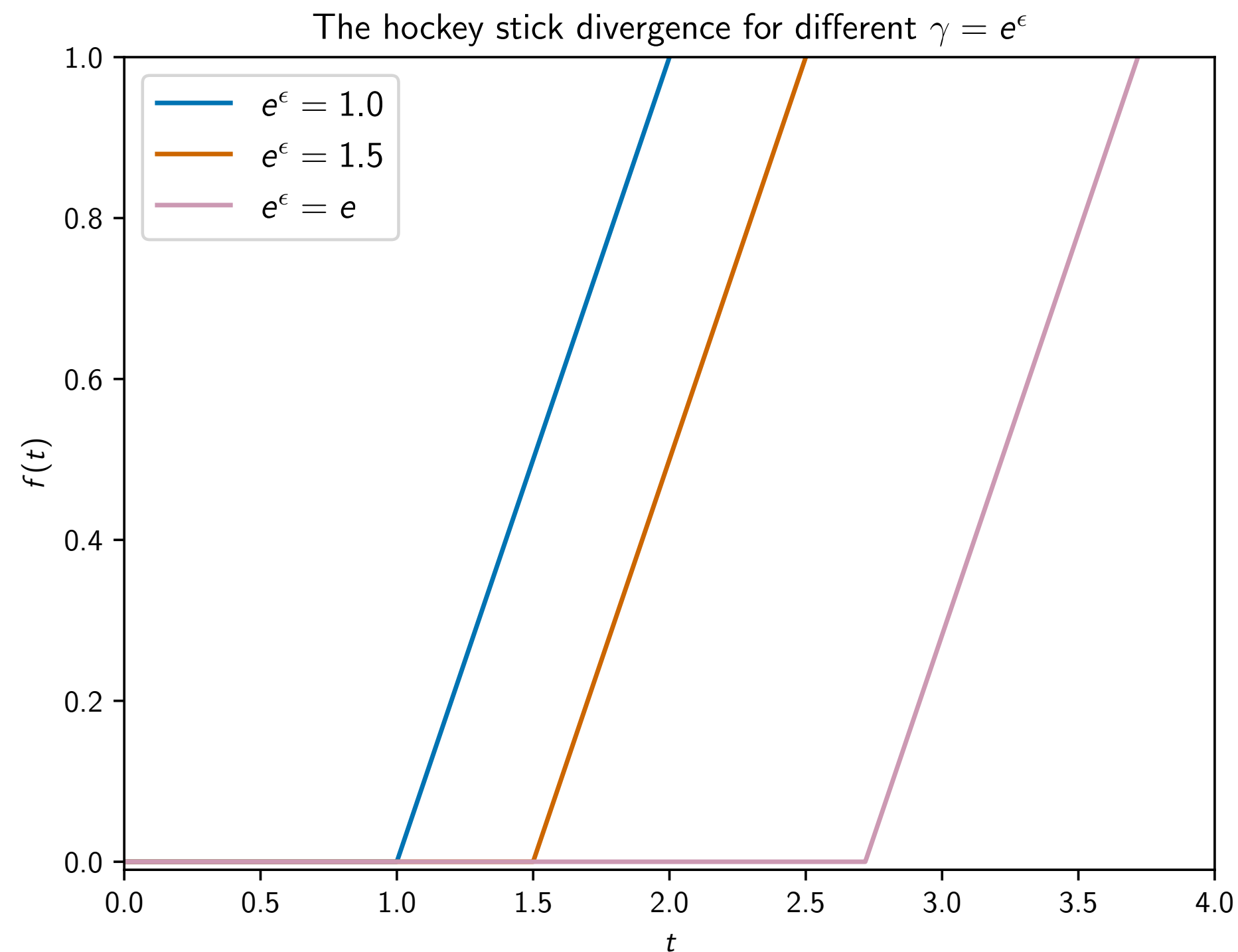
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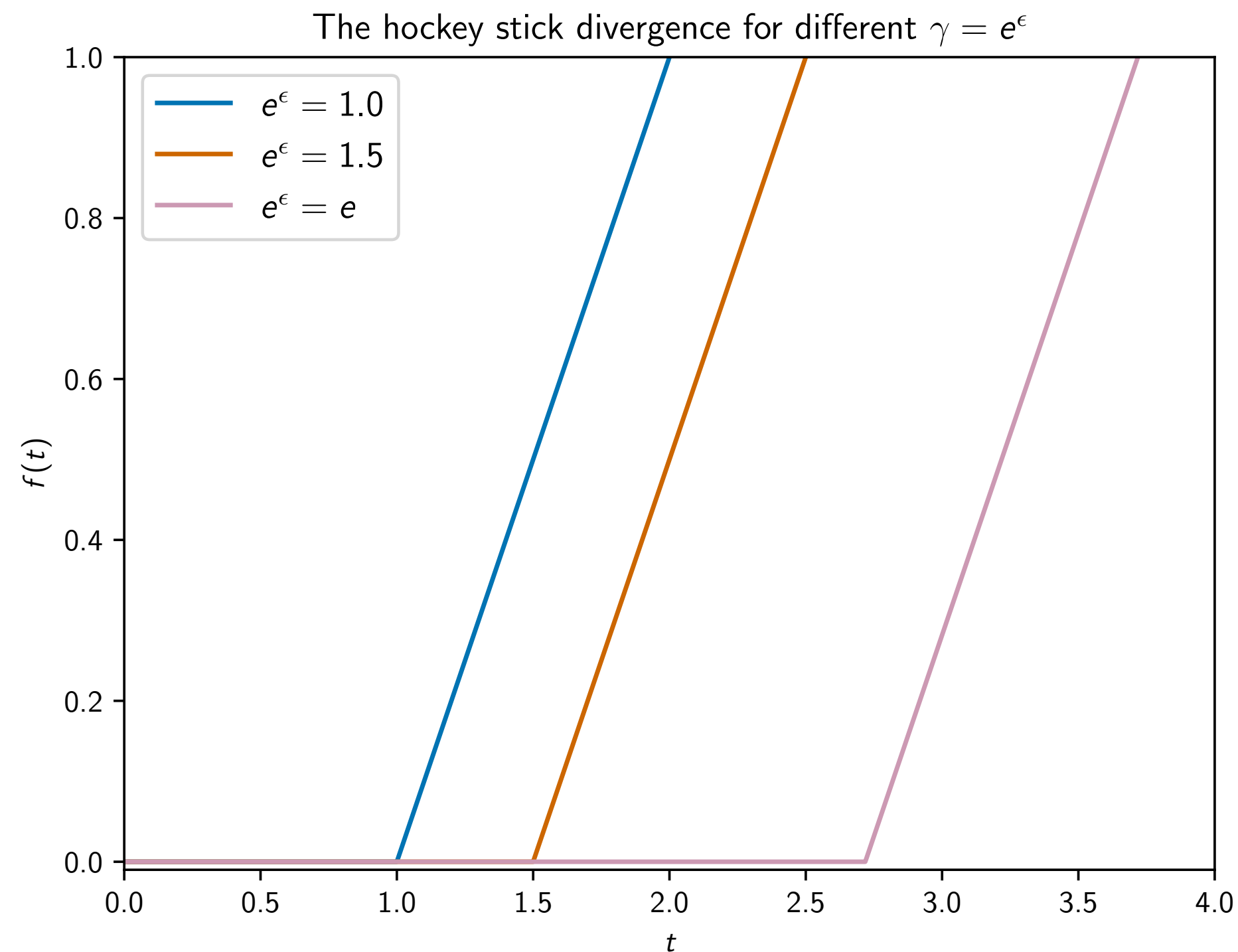
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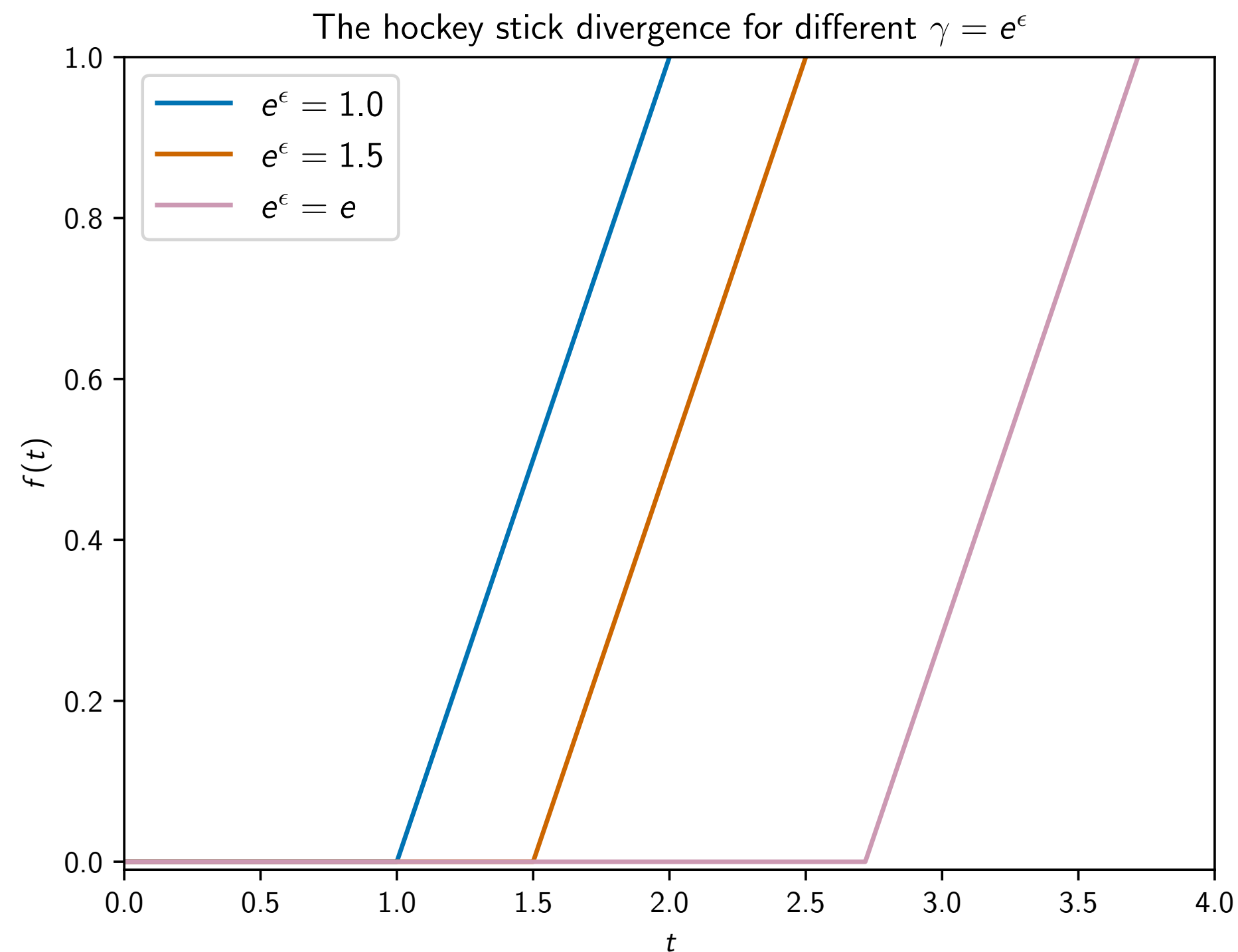


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Where  $L$  is the PLRV corresponding to  $(\mu, \nu)$ .

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- We can use these dominating pairs to bound the loss for compositions.

# Composition and divergences

**Adding things up**

With thanks to Flavio Calmon, Oliver Kosut, and Shahab Asodeh!

# Composition and divergences

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# **A taxonomy of composition analyses**

**(on non-interactive settings, also non-exhaustive)**

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## Large deviations

Bun and Steinke (2016)

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(on non-interactive settings, also non-exhaustive)

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## Warning about subsampling!

Lebeda, Regehr, Kamath, Steinke (2024)

Chua, Ghazi, Kamath, Kumar, Manurangsi, Sinha,  
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# Tilting a distribution

Maintaining exactness for composition

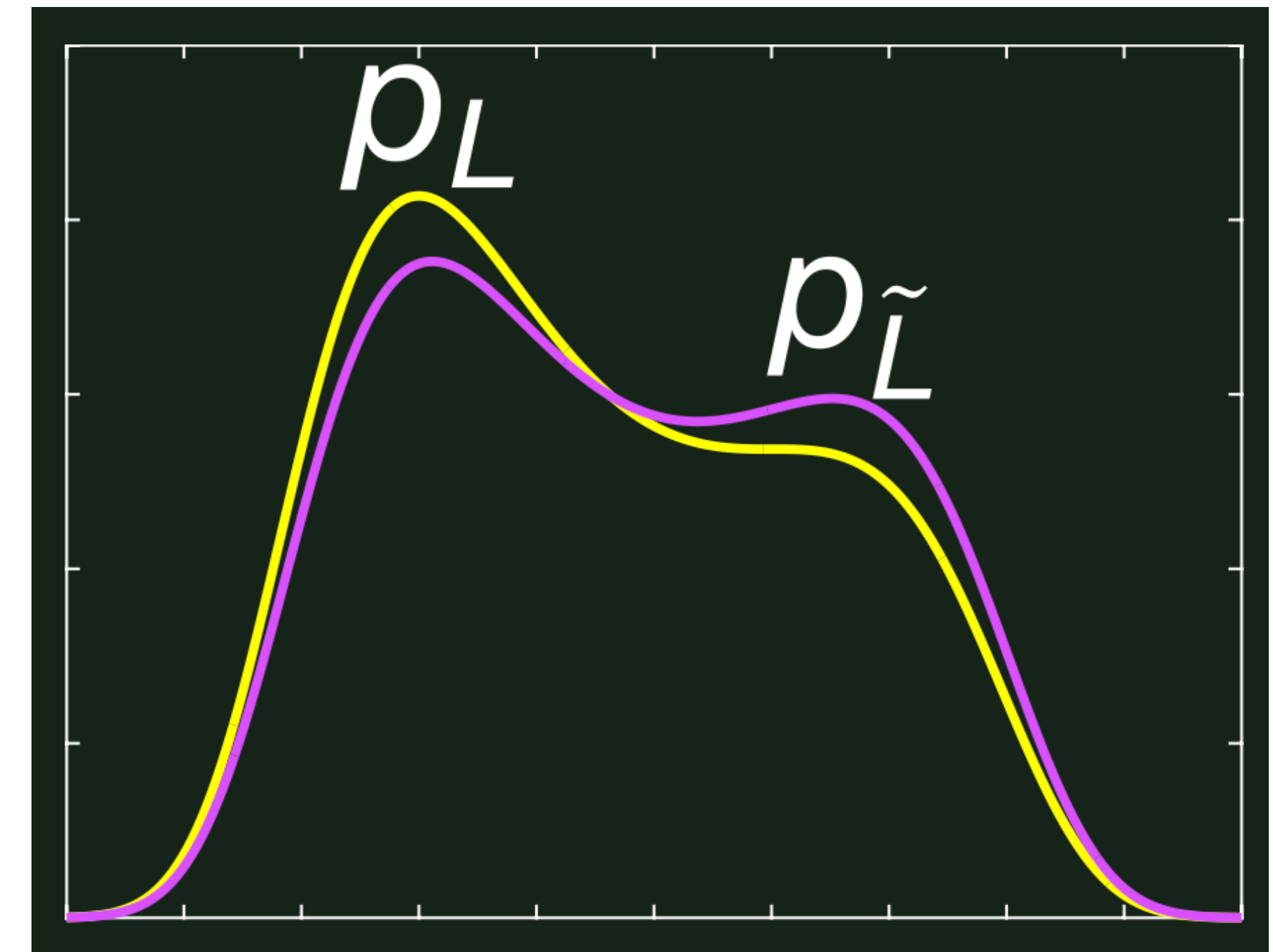


Figure: Oliver Kosut

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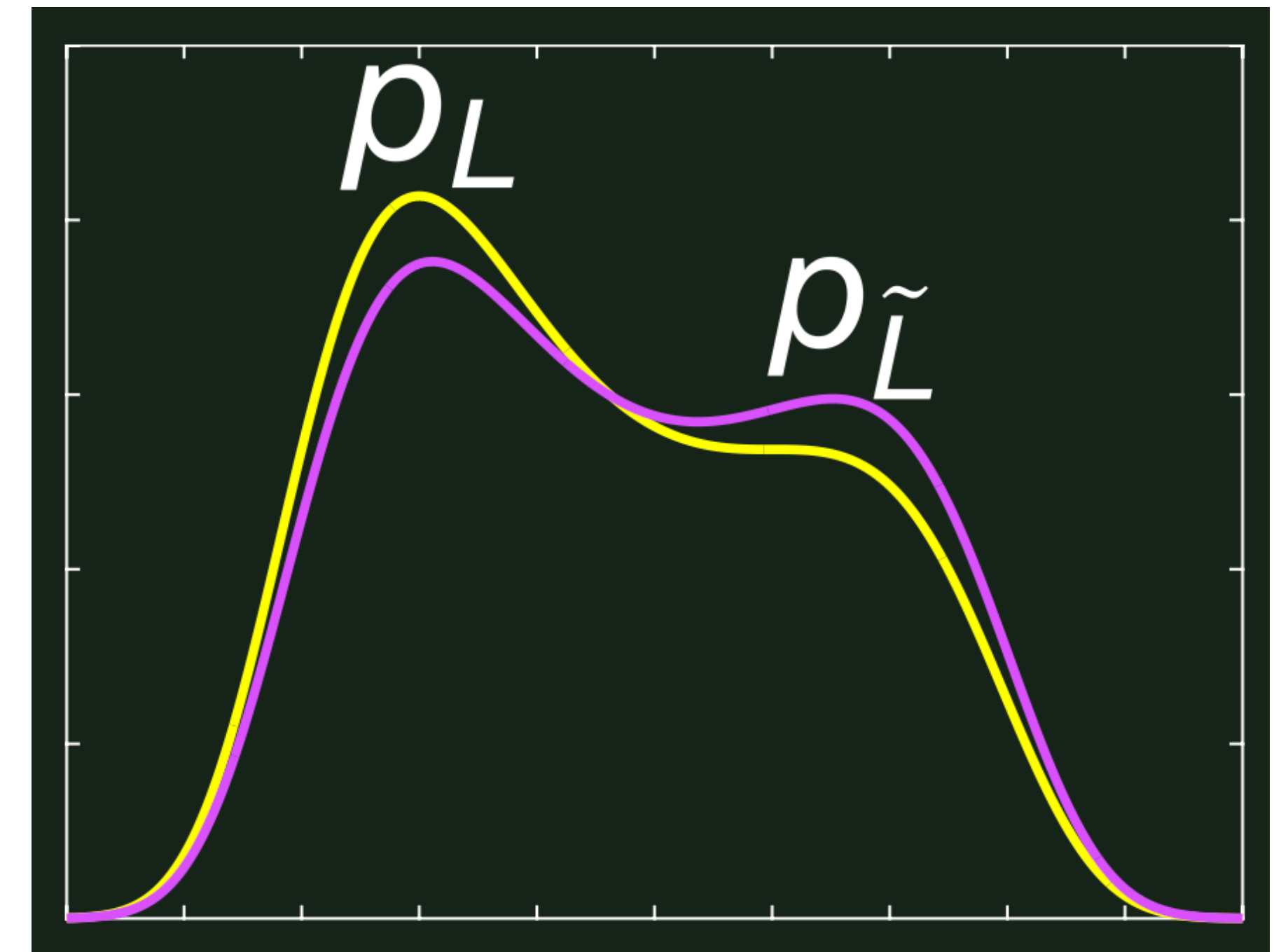


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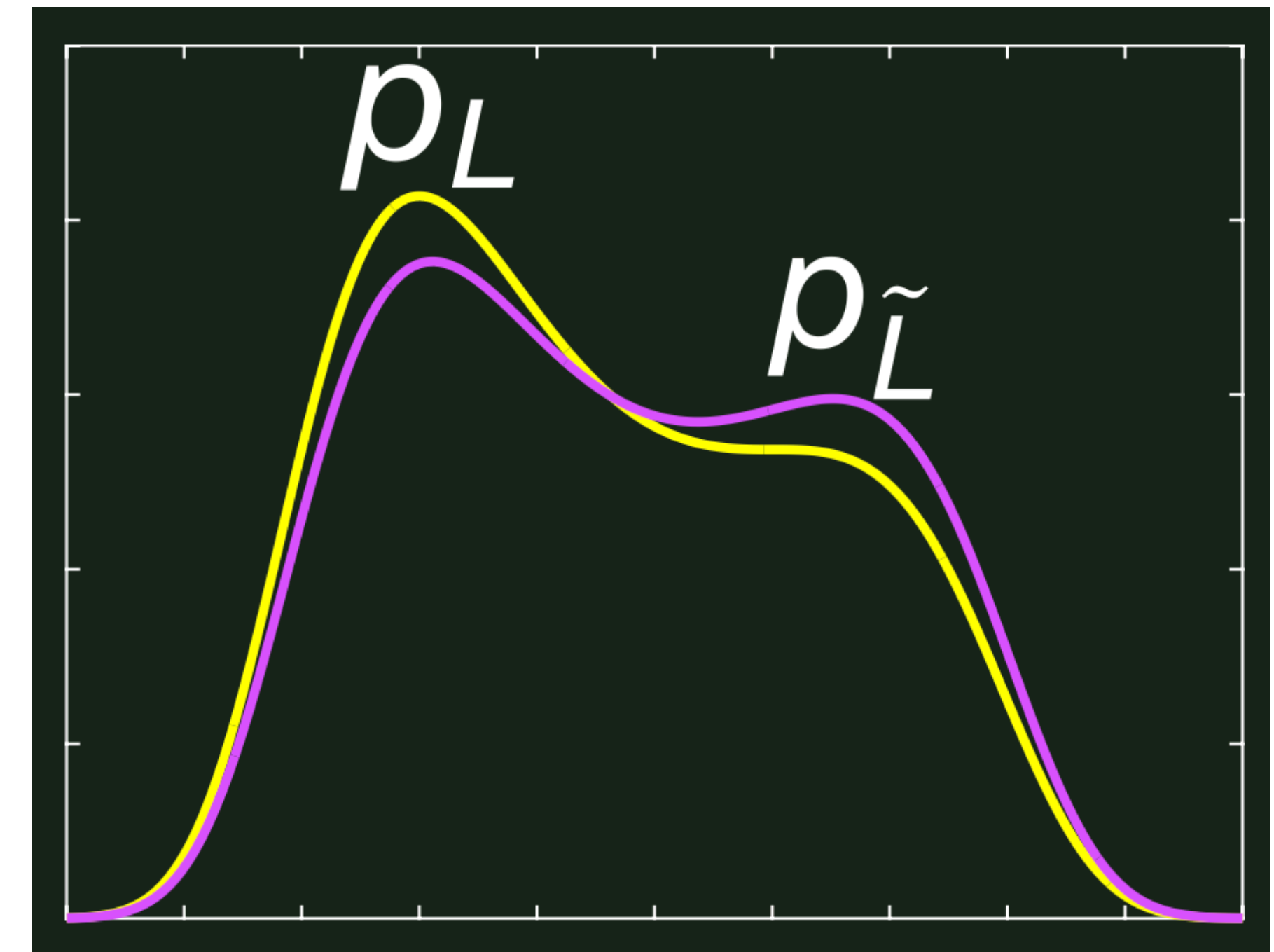


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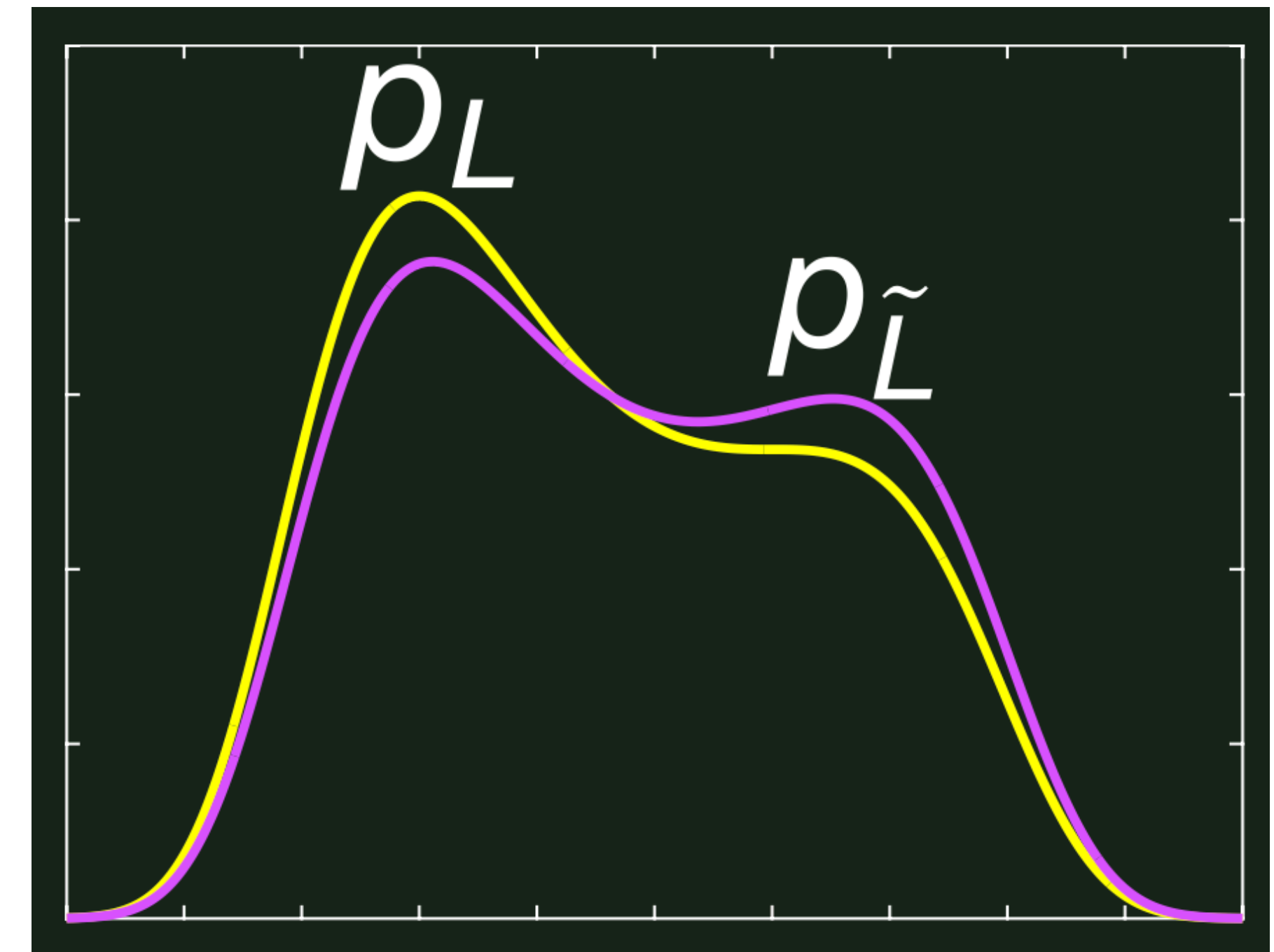


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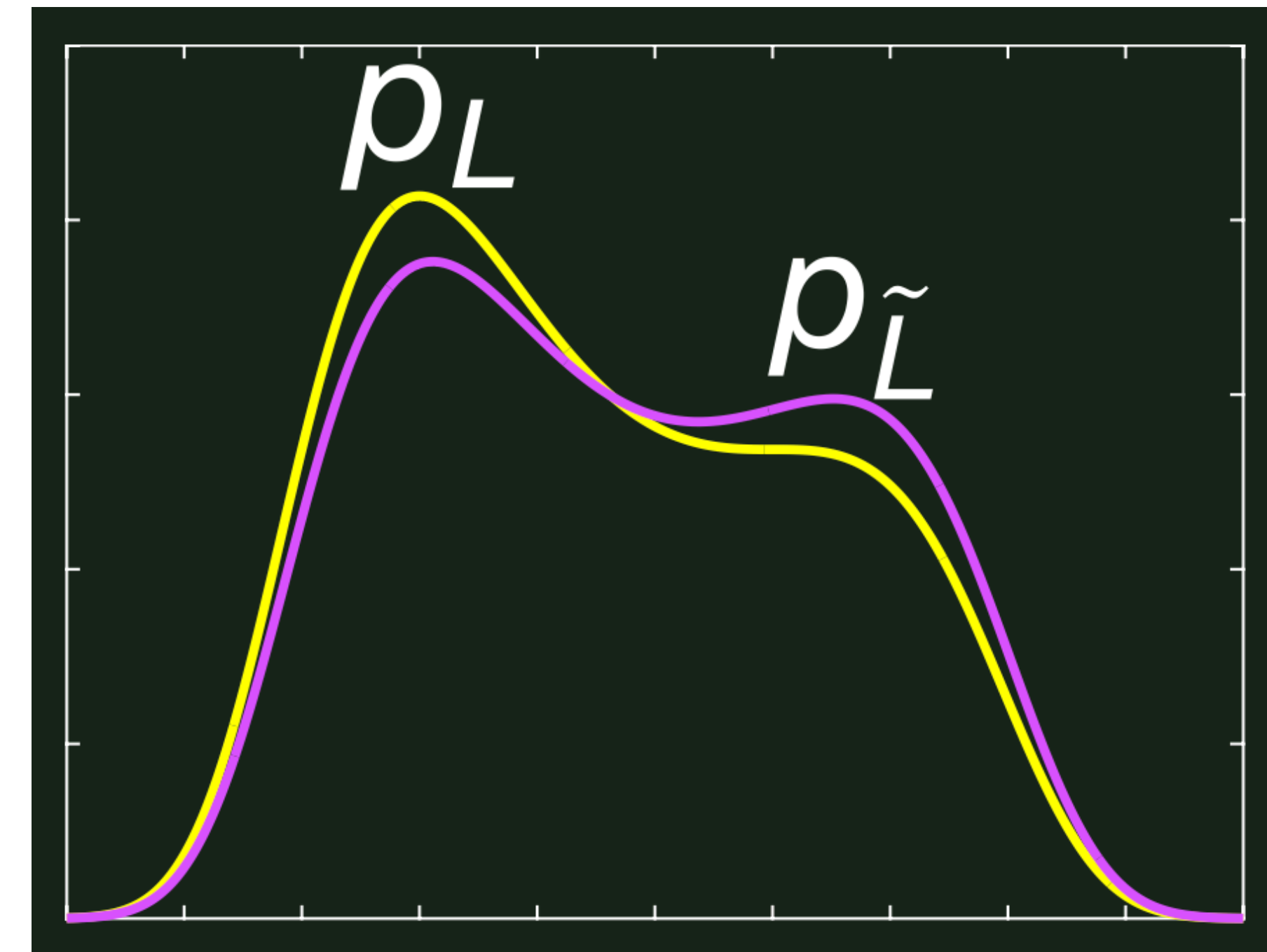


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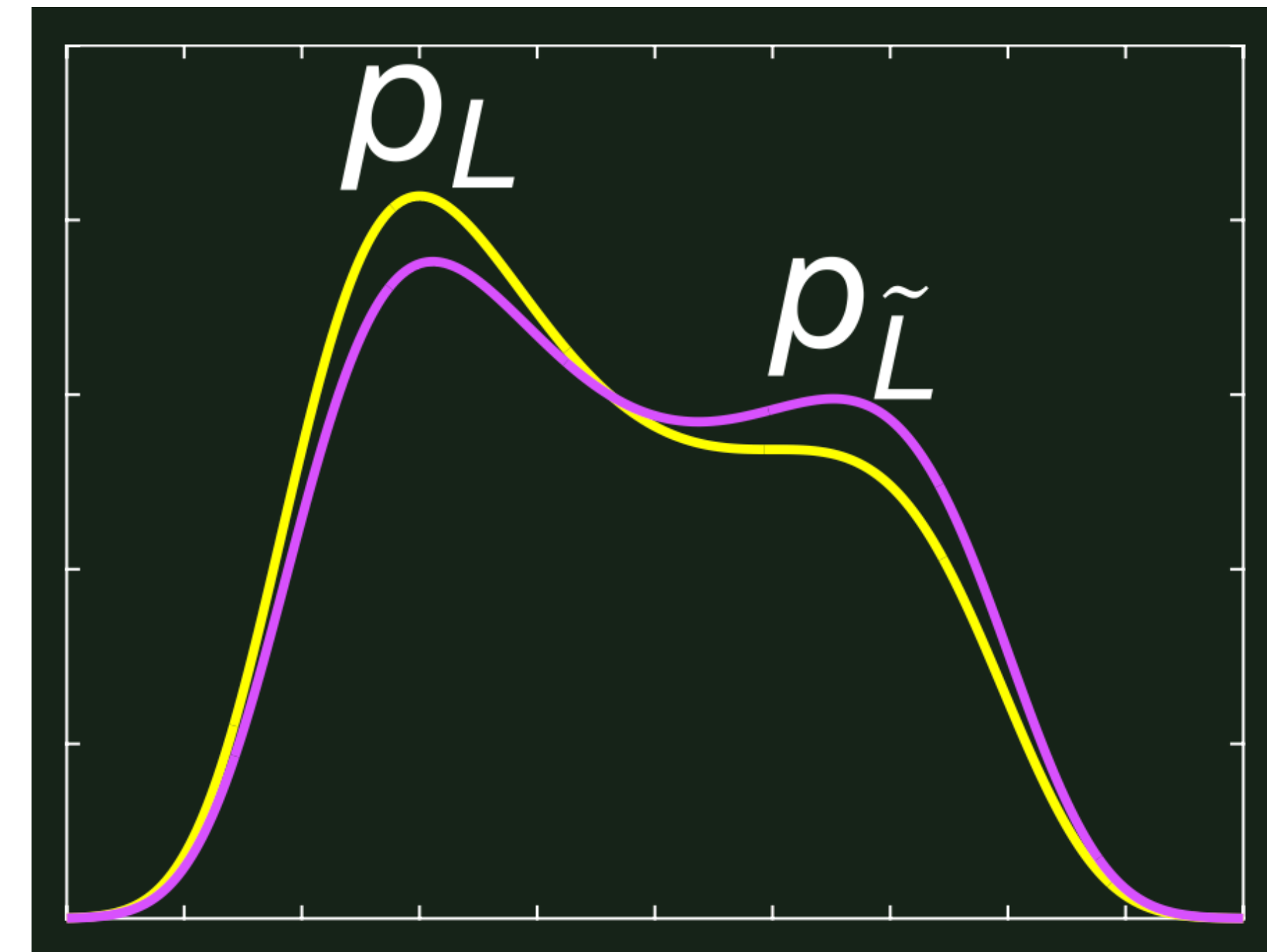


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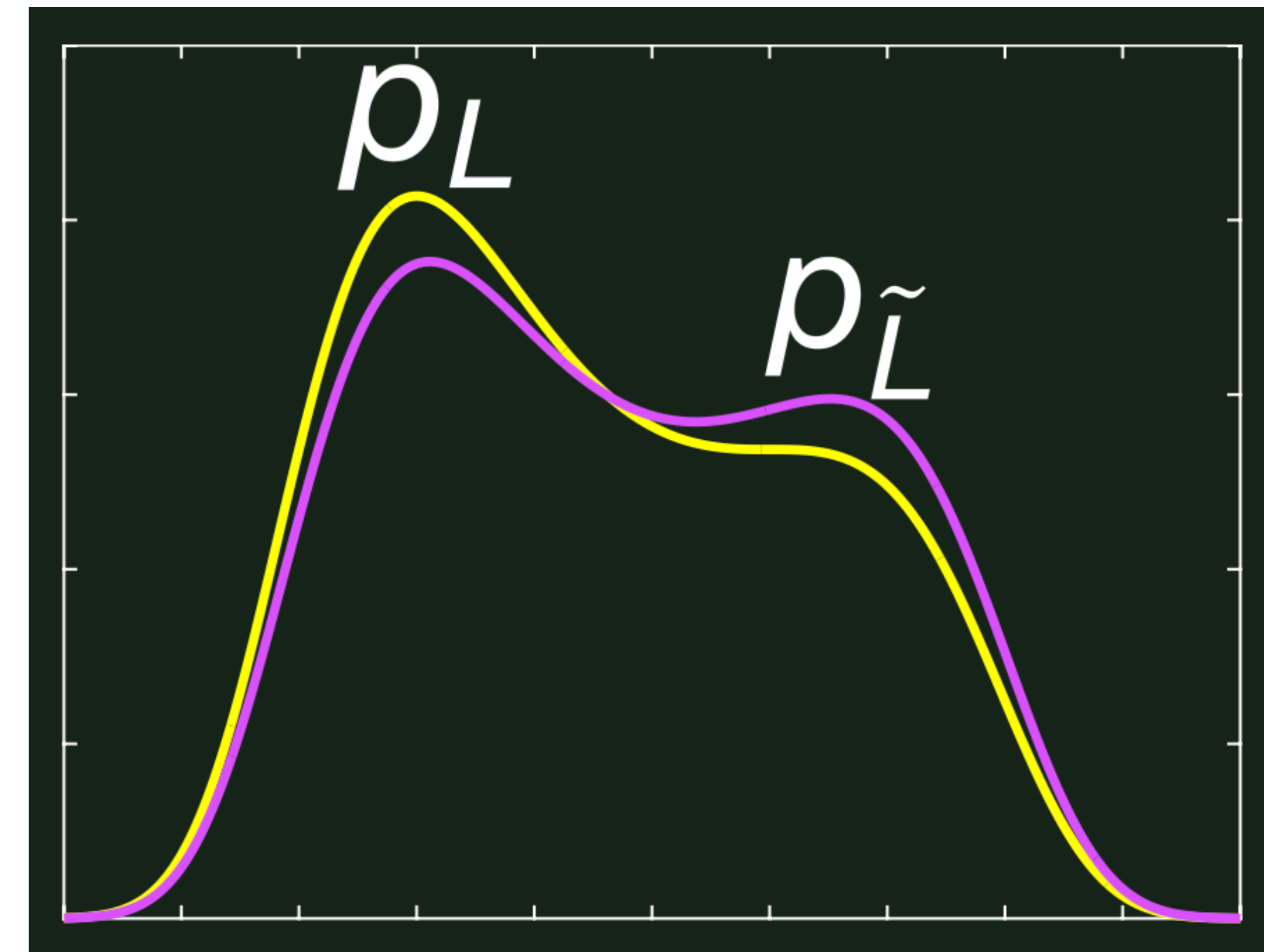


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Can use this to derive a “**saddle-point**” accountant in terms of the exponent.

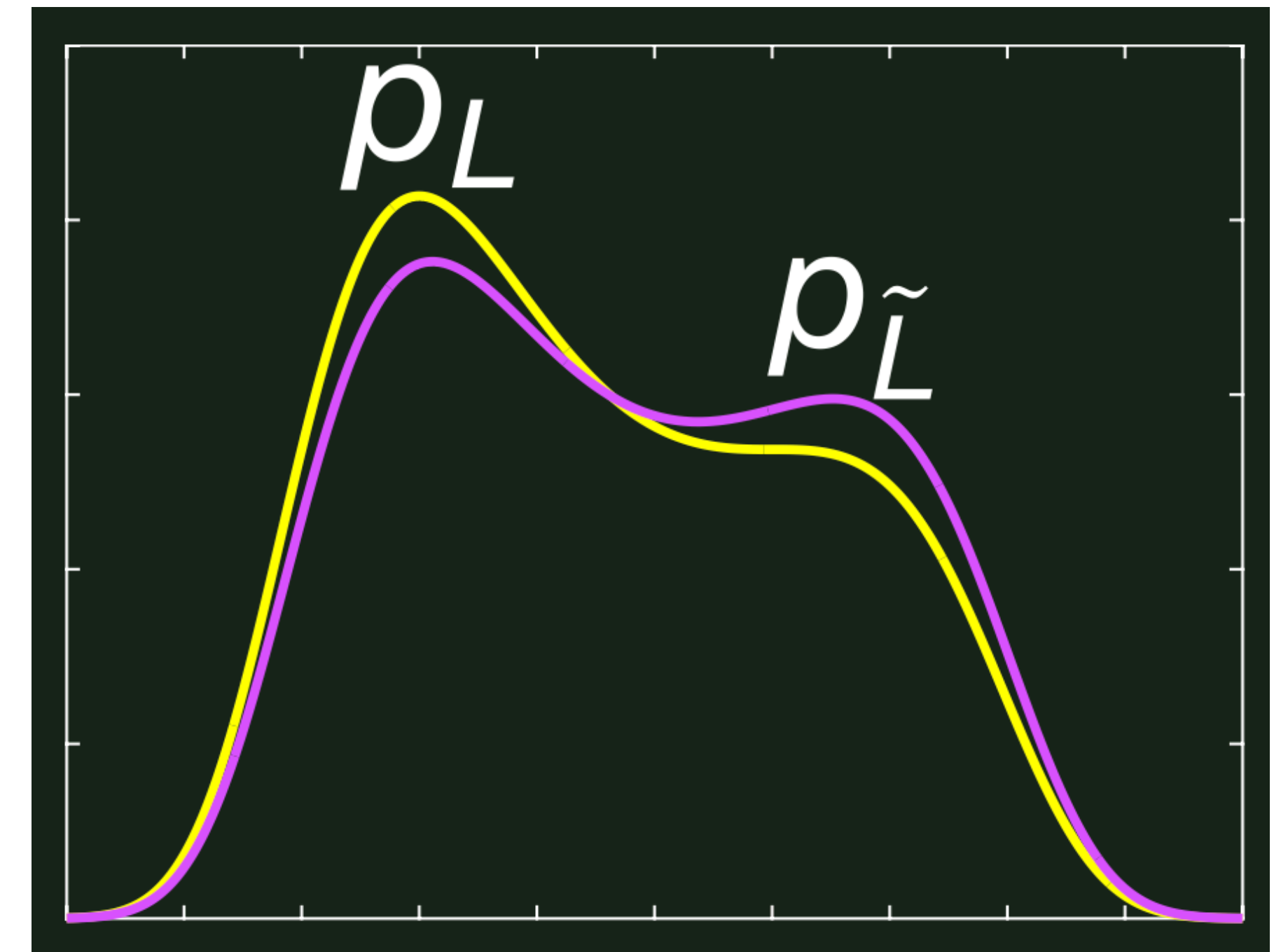


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# Tilting in other contexts

(Beyond *Don Quixote*)



A tiling (not tilting) by  
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Perhaps of interest to folks here? Botev (2017) uses it to exact iid simulation from the truncated multivariate normal distribution.





Shichiri Beach in Sagami  
Province

相州七里浜

Soshū Shichiri-ga-hama

Vista 4

contraction coefficients/iteration



# Maximum likelihood and ERM

## Optimization and privacy

[Chaudhuri, Monteleoni, Sarwate 2011] [Zhang, Zhang, Xiao, Yang, Winslett 2012]

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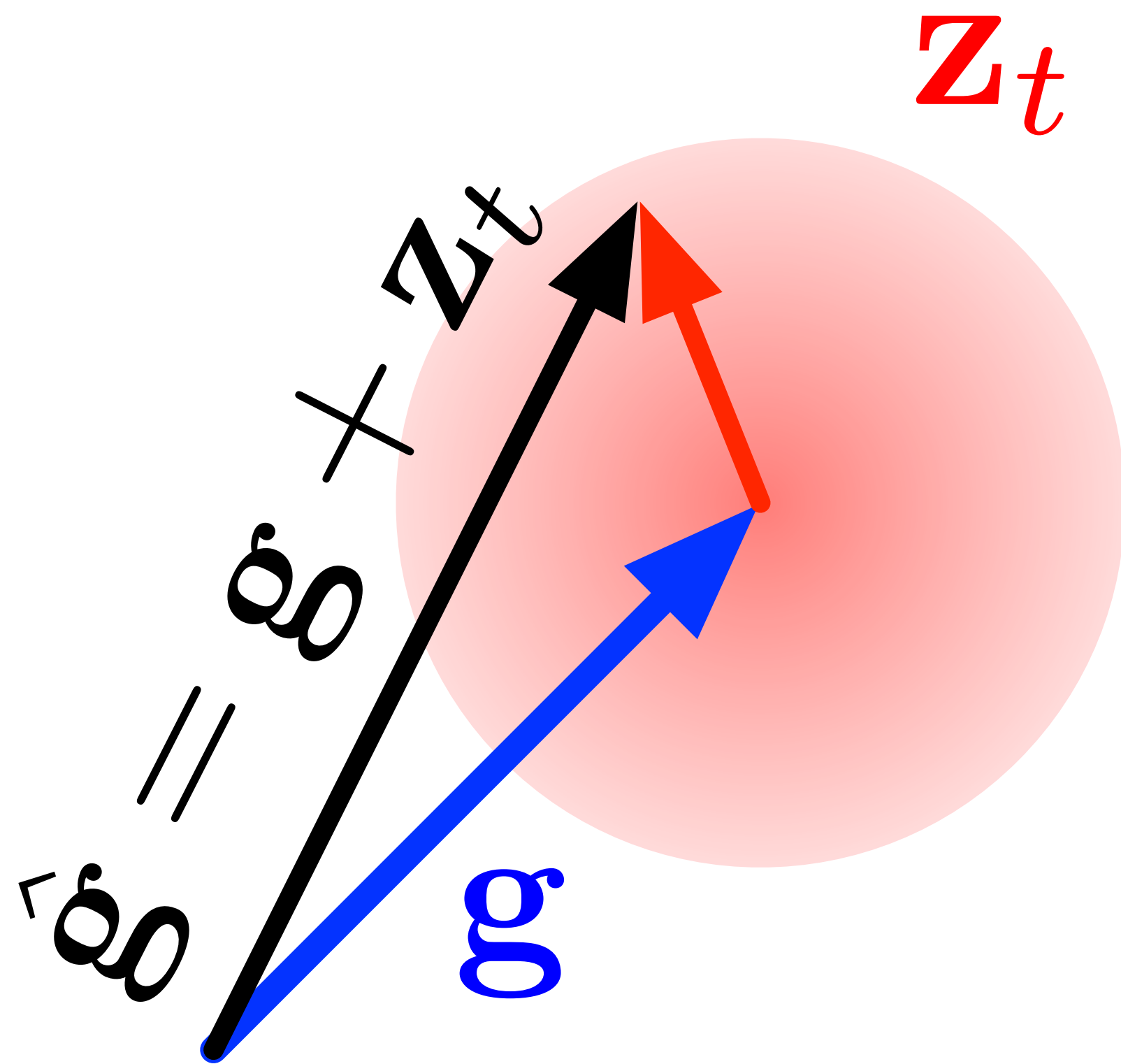
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- **“Functional mechanism”**: Add noise to an approximation of the loss function  $\ell(\cdot)$ .

# Deep Learning and DP

Privacy for neural networks

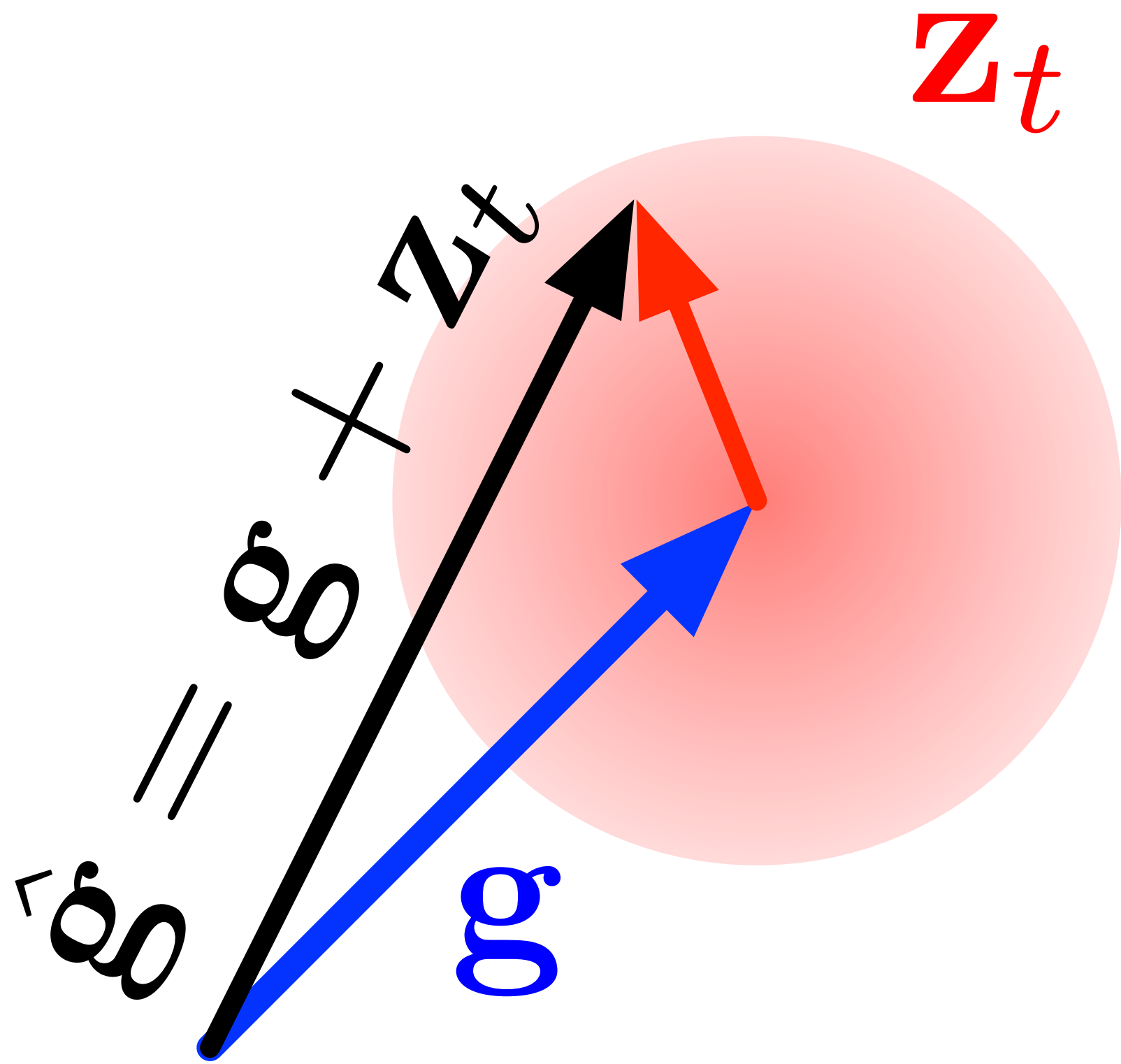


[Song et.al. 2013, Duchi et.al. 2014, Abadi et.al. 2016, Mironov 2017]

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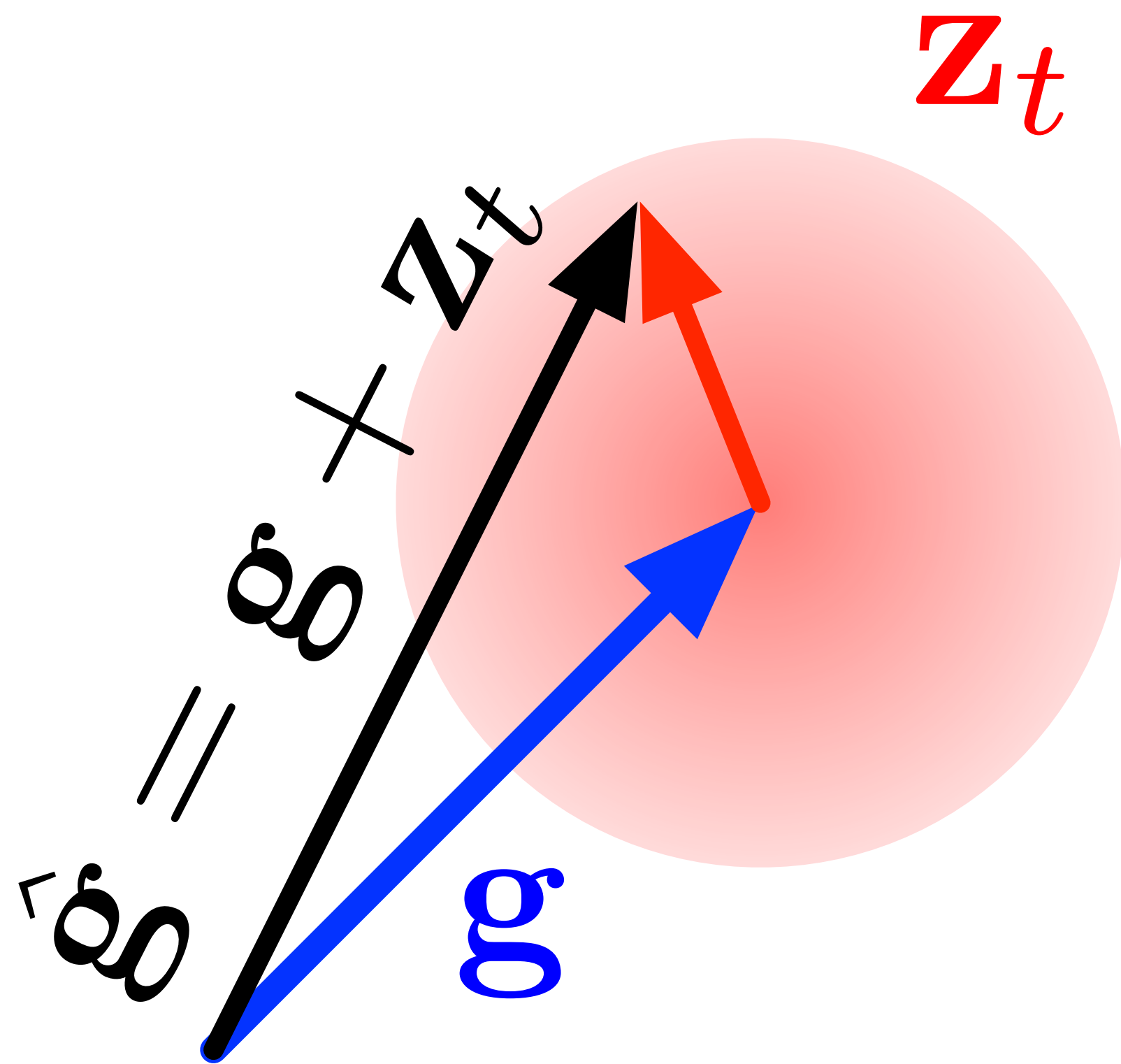
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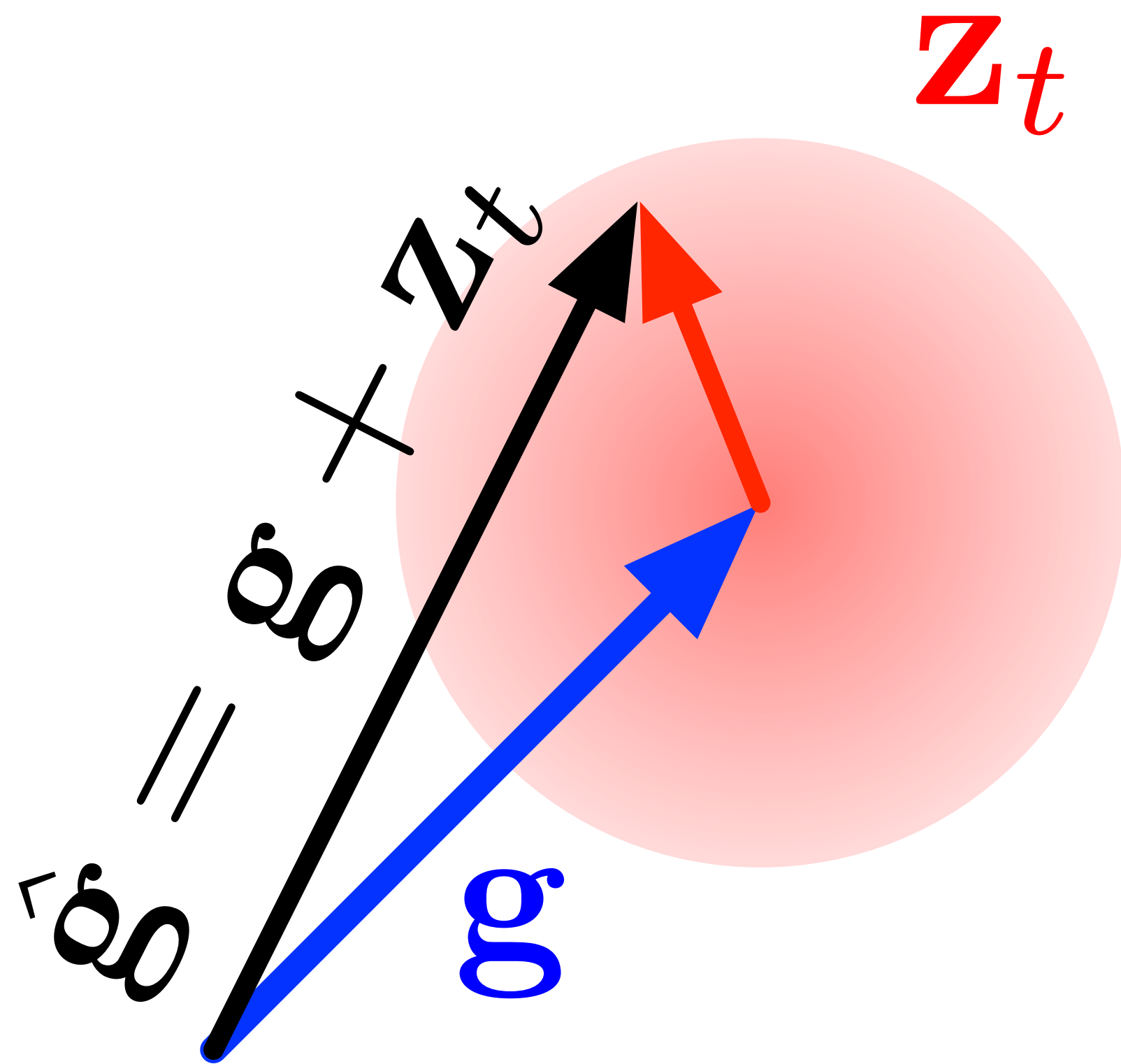


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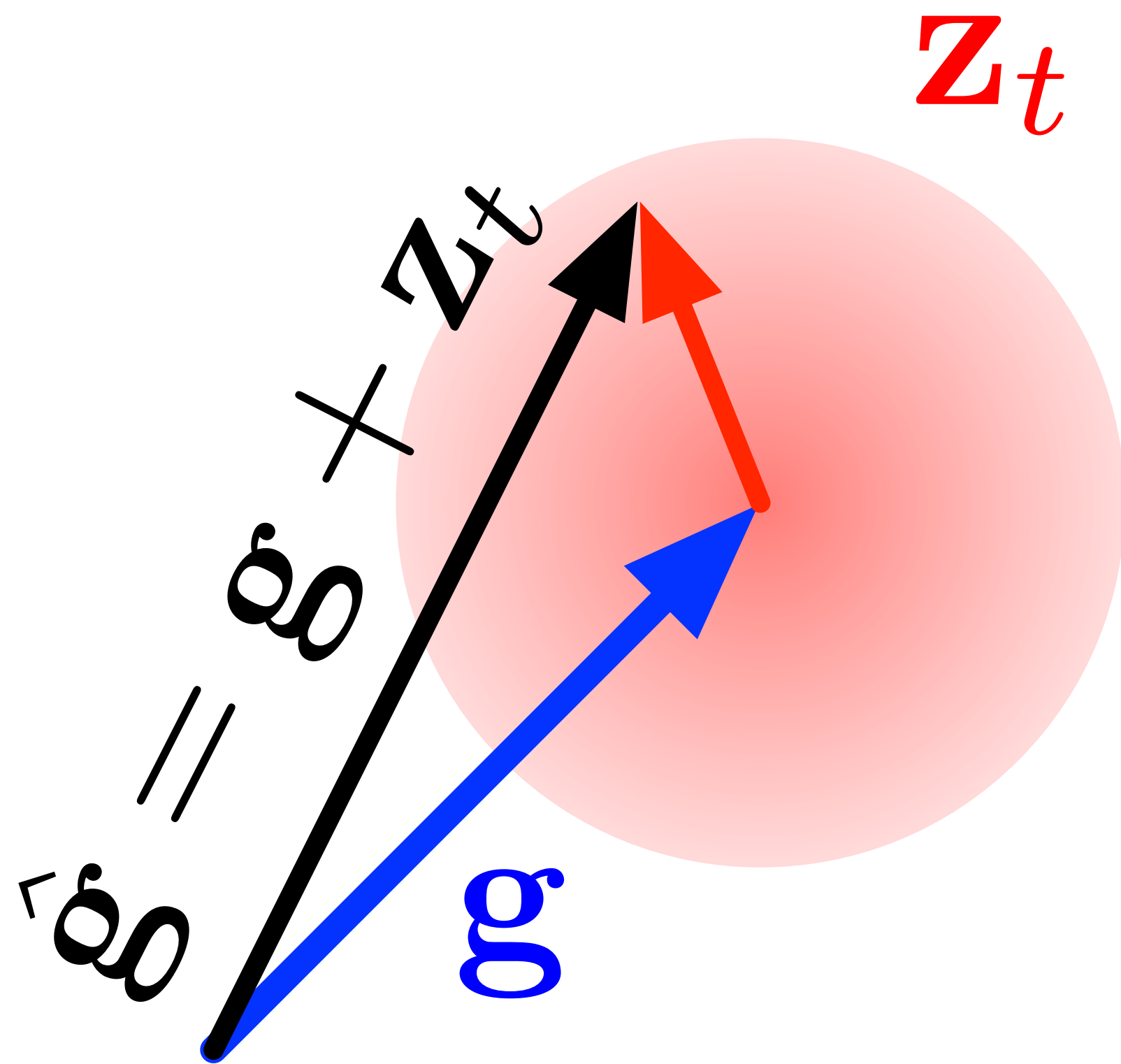
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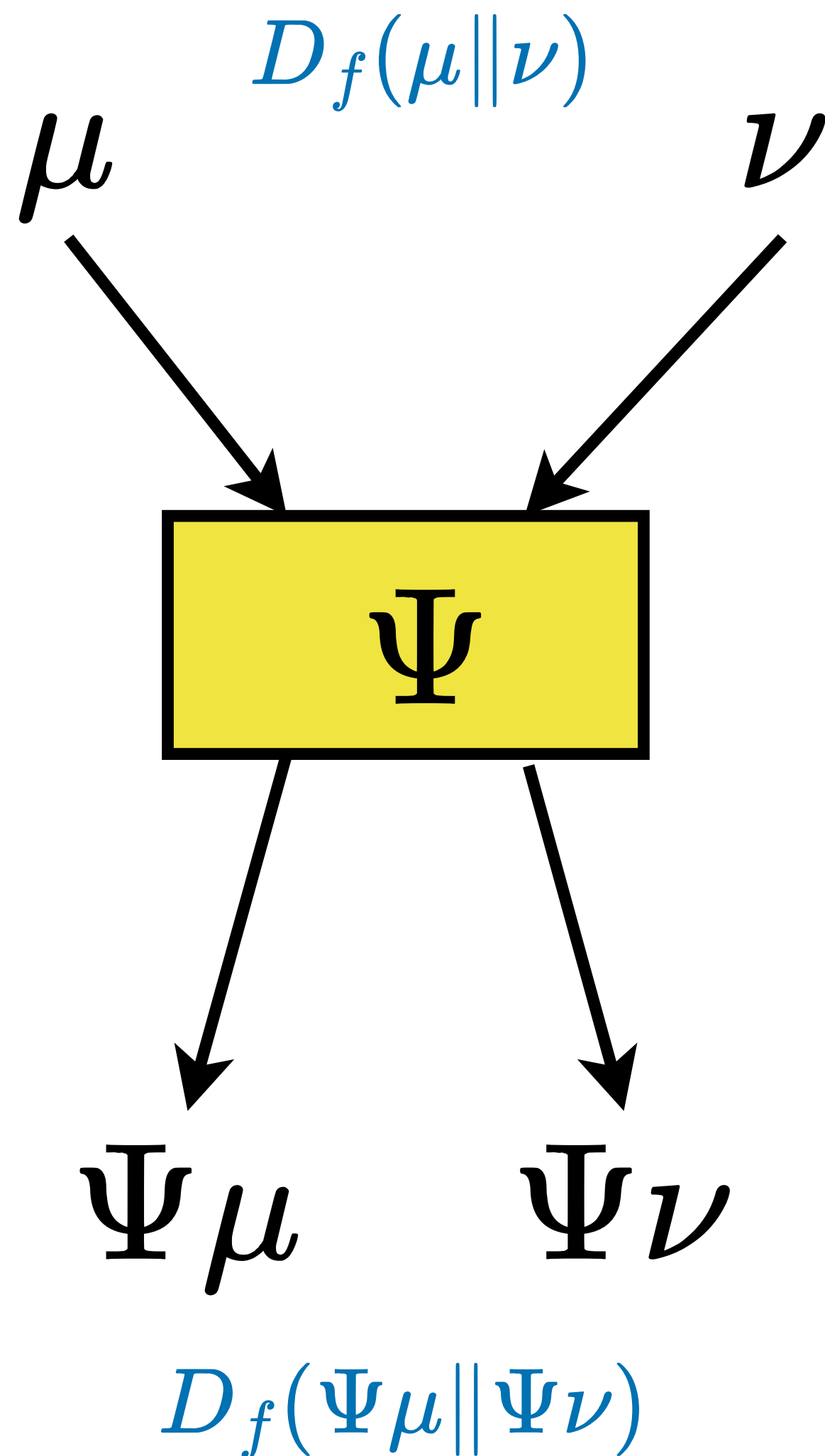
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**Privacy accounting** lets us track the overall privacy loss.



# Strong data processing inequalities

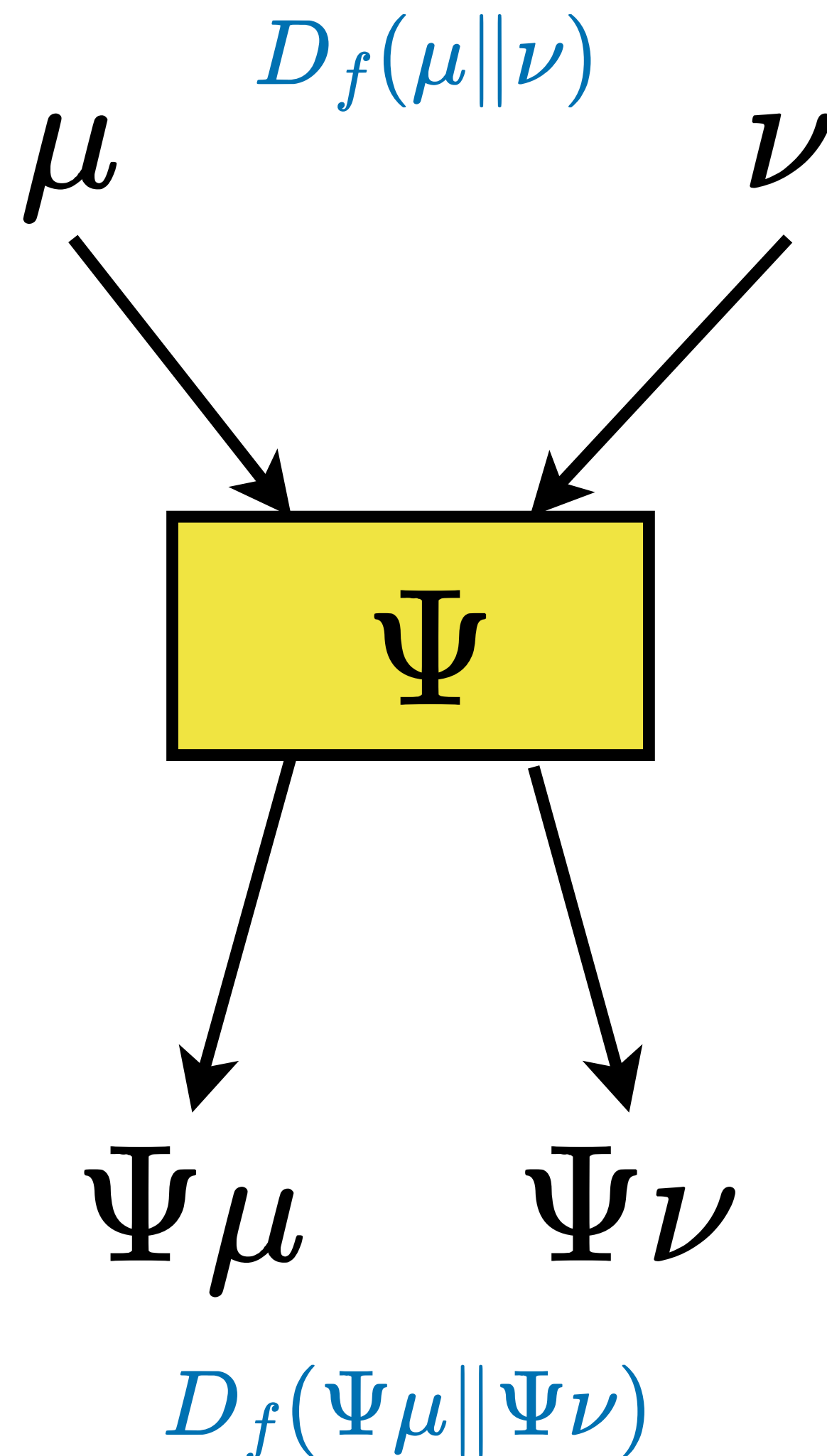
Quantifying the privacy gain from post-processing



Dobrushin (1956), Ahlswede, Gács (1976)

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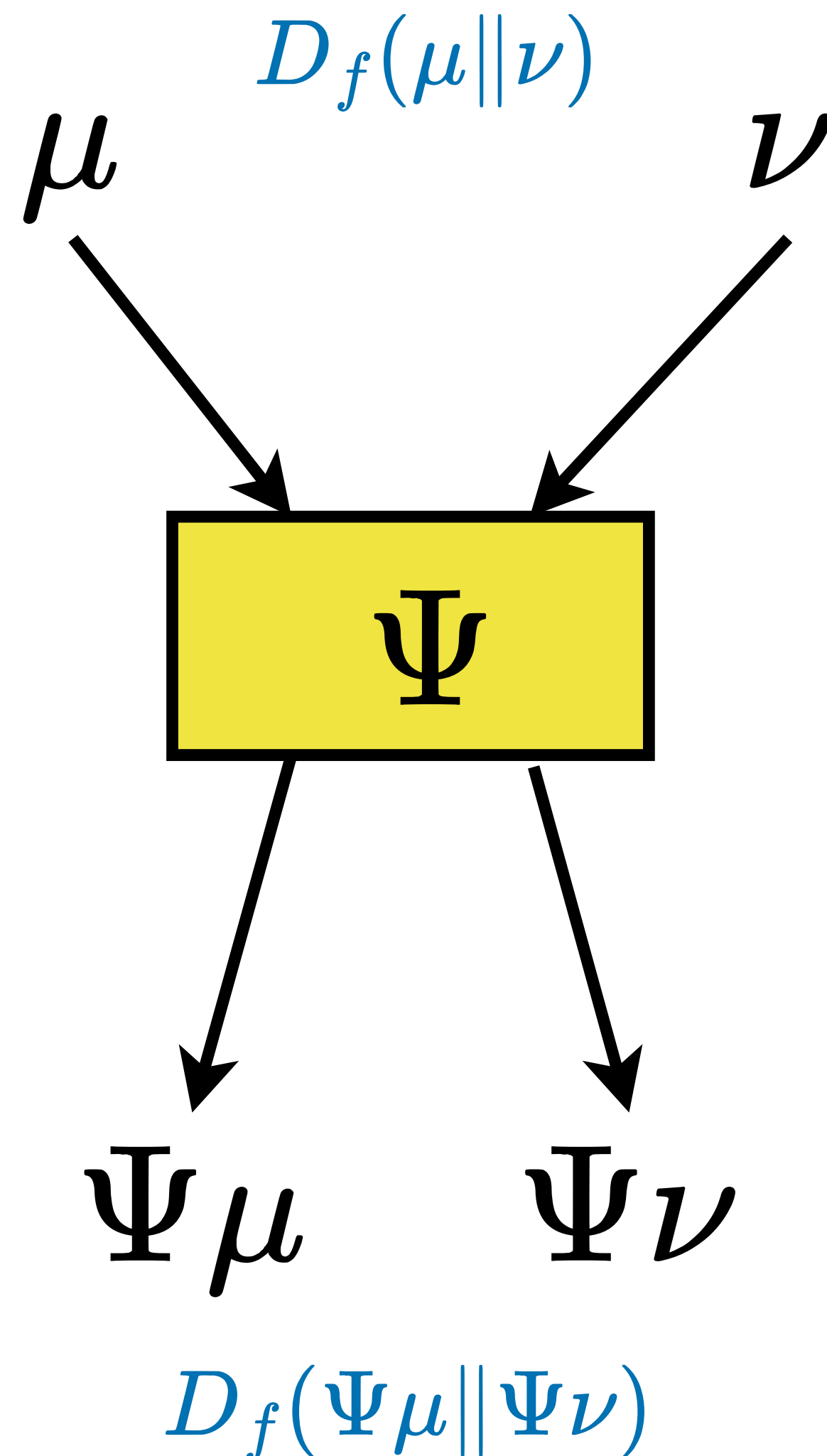
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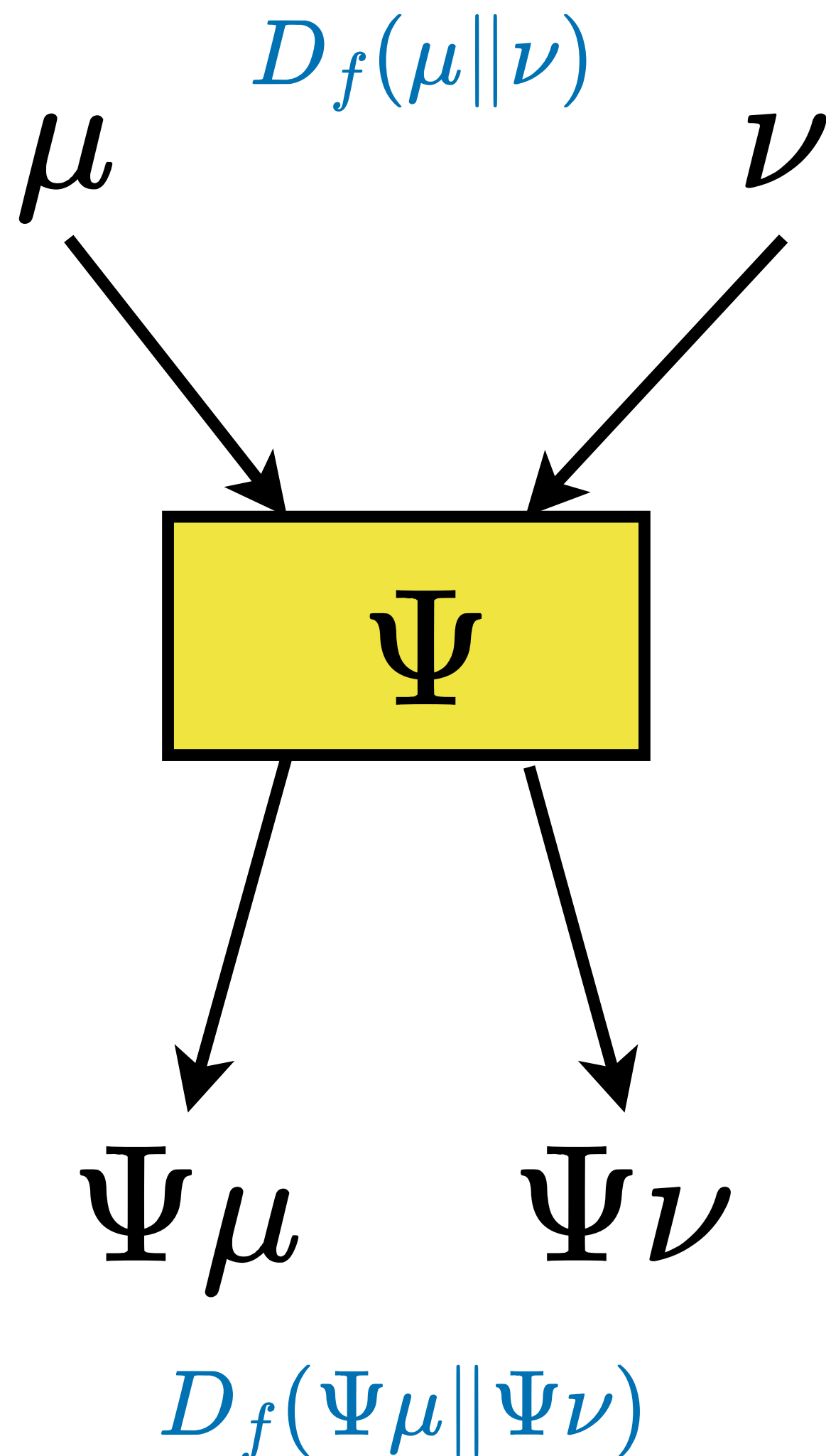


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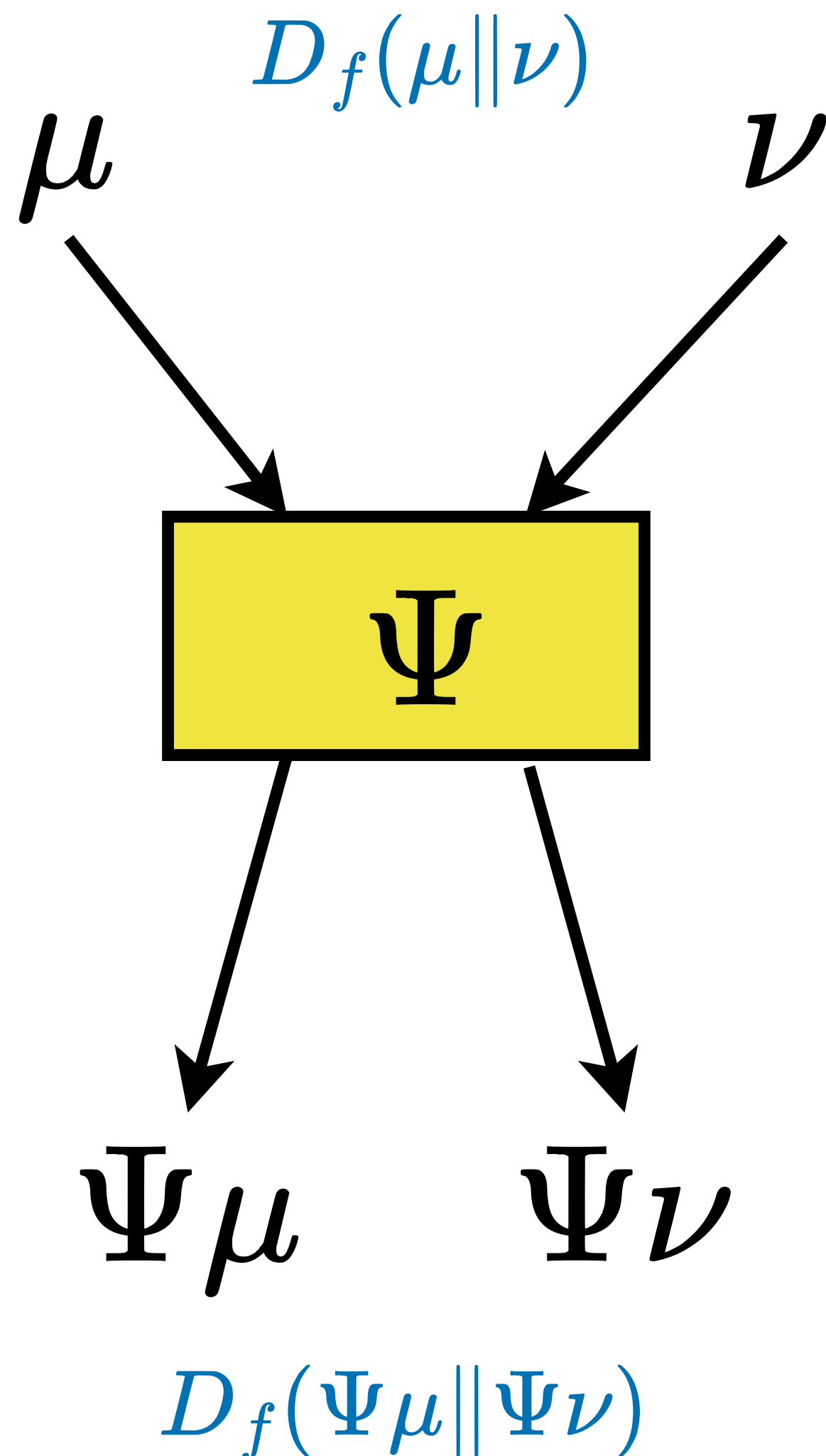
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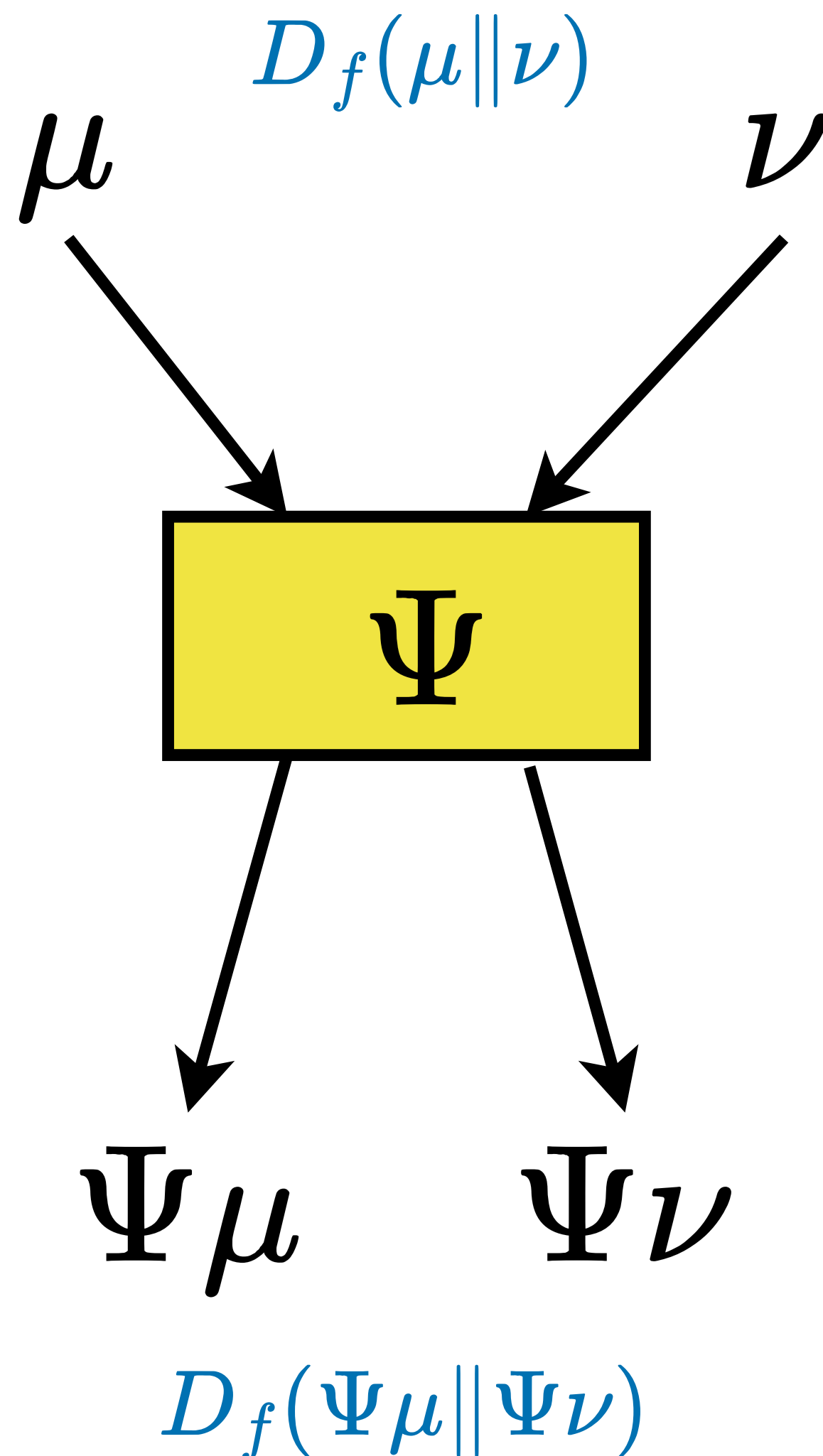
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If  $\eta_f(\Psi) > 0$  this is a **strong data processing inequality (SDPI)**.

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**Applications to DP-SGD and LDP**



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Idea: analyze privacy for the *last* iterate by using the contraction for the  $E_\gamma$  divergence.

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**An abbreviated timeline**

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- Asodeh, Diaz (2024) - use data processing inequalities to remove convexity and smoothness assumptions for projected DP-SGD and regularized DP-SGD.

# Contraction and Bayesian estimation

**Focusing on the local model**



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$\theta$  is a secret, the loss  $\ell$  is a negative gain, and we look for the maximally leaky channel subject to an  $(\varepsilon, \delta)$  constraint... (Is this right?)





**Morning After a Snowfall  
at Koishikawa**

**礫川雪の旦**

**Koishikawa yuki no  
ashita**

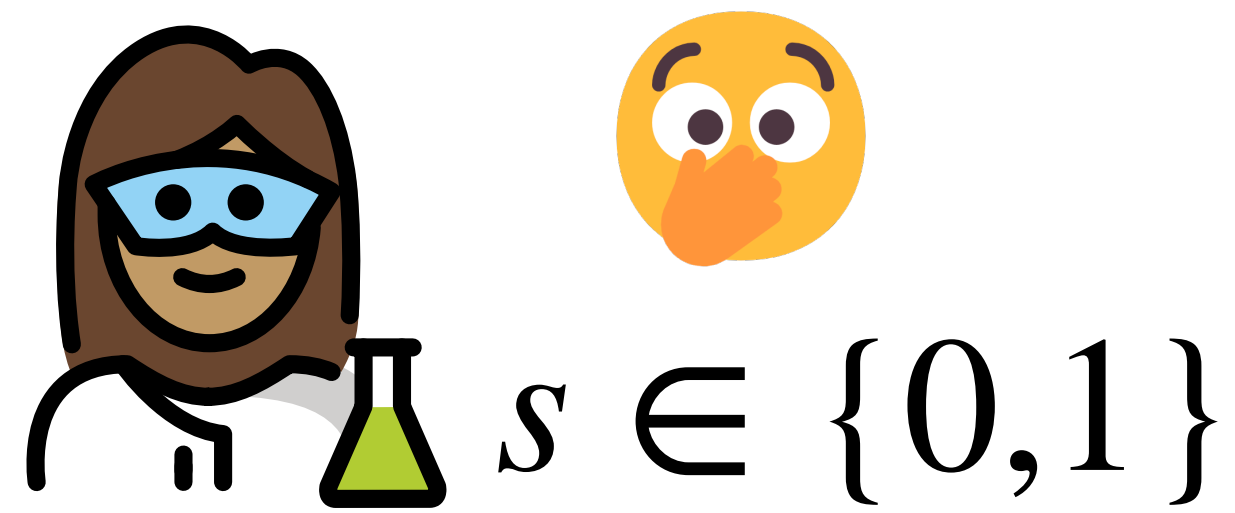
**other destinations**



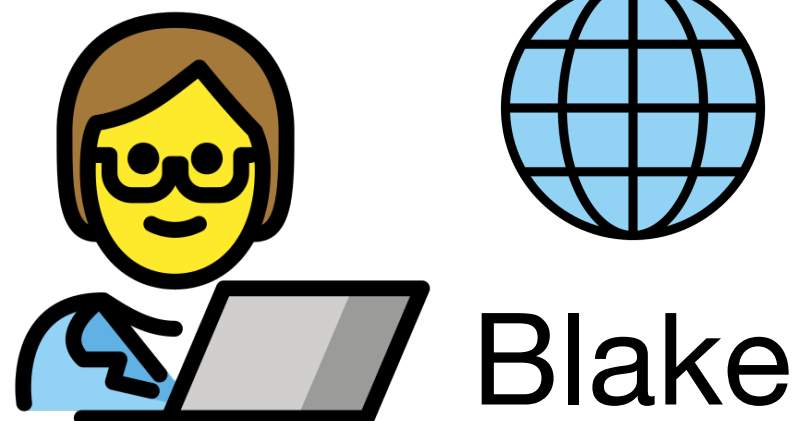
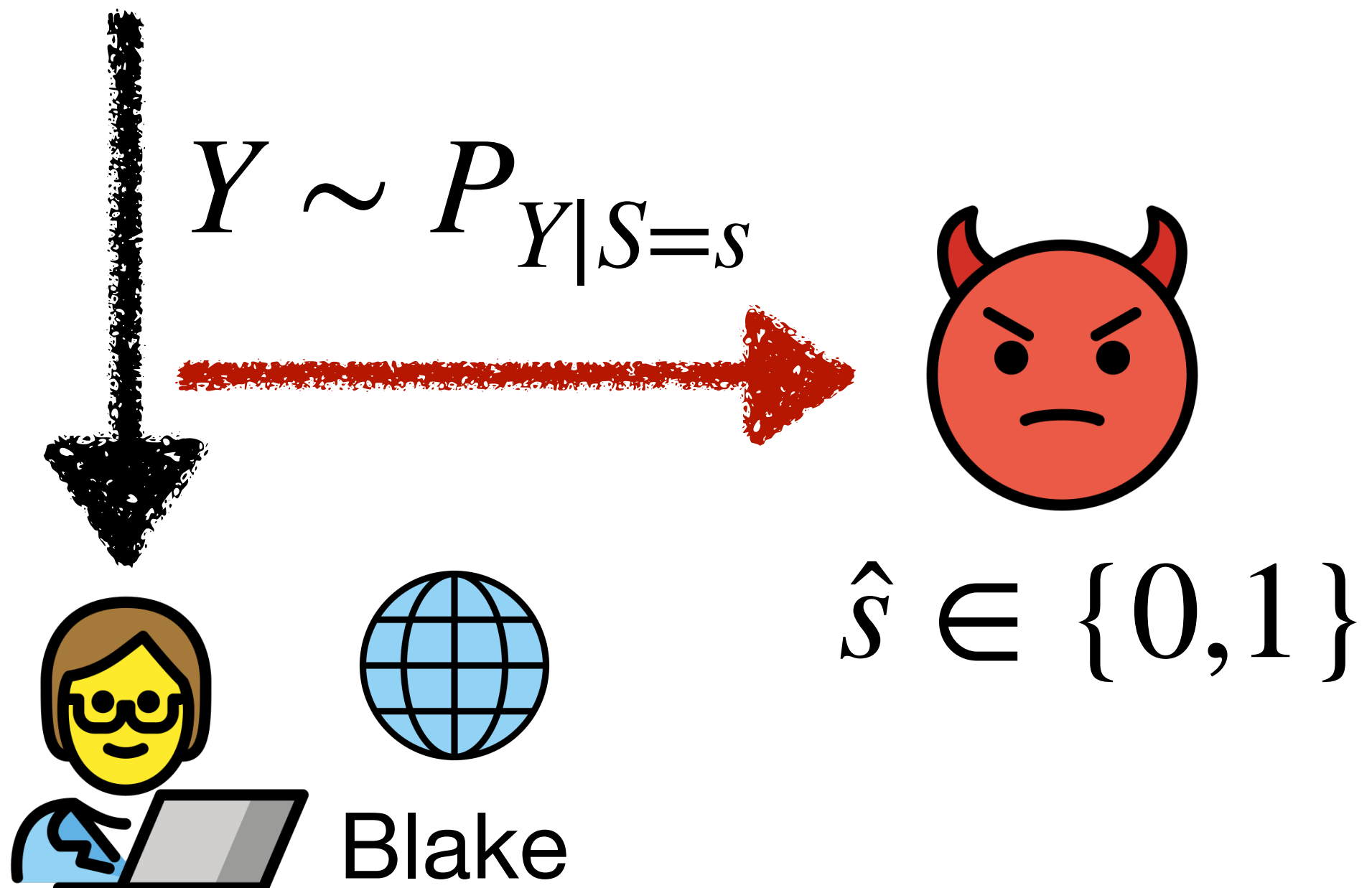
# What we've seen so far

Let's start simple

Sasha



$$Y \sim P_{Y|S=s}$$

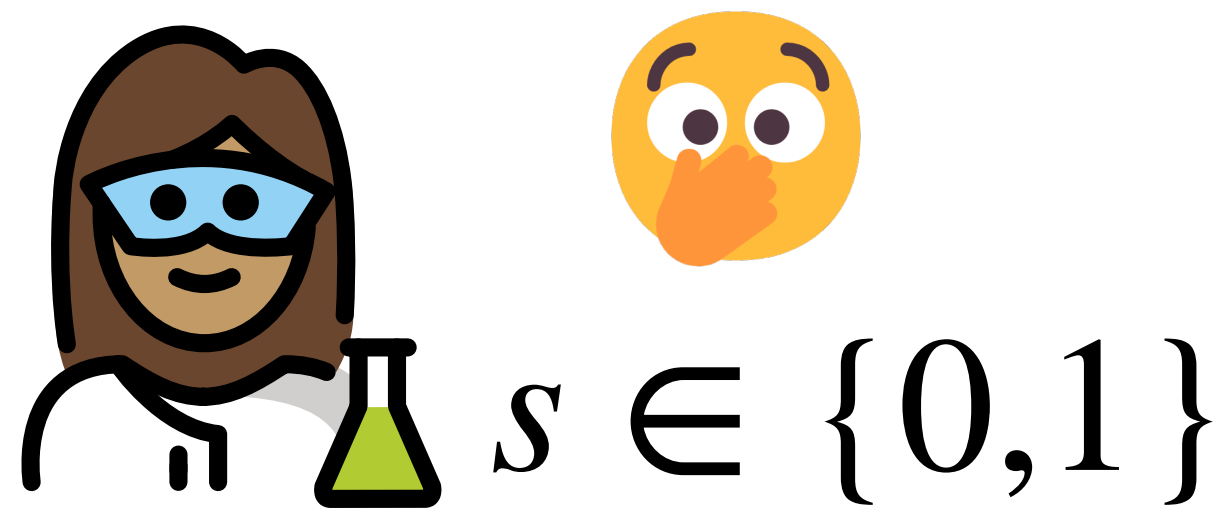




# What we've seen so far

## Let's start simple

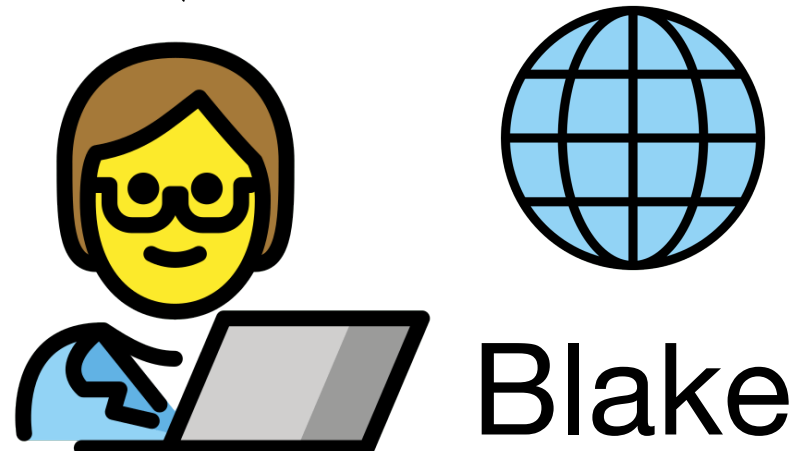
Sasha



$$Y \sim P_{Y|S=s}$$



$$\hat{s} \in \{0,1\}$$



We started out with a simple story: protecting a single bit.

- Differential privacy both is and is not just as simple as hypothesis testing.
- Taking an information-theoretic view opens the door to better analyses.
- The gap between algorithms and analysis is shrinking.
- The gap between algorithms and applications is still large.

# The gap between theory and practice

**It's wider than you might think**



# The gap between theory and practice

**It's wider than you might think**

There are lots of issues with implementing differential privacy in practice:

- Approximate versus exact sampling (and side channels)
- Approximate versus exact optimization
- “Privacy amplification” and its implementation
- Numerical precision and floating points
- Managing privacy budgets





安野平兵衛  
市川市紅





Several interesting  
challenges left for:





Several interesting  
challenges left for:





Several interesting  
challenges left for:

maths





Several interesting  
challenges left for:

maths

computational stats







Several interesting  
challenges left for:

maths  
computational stats  
engineering





Several interesting  
challenges left for:

maths

computational stats

engineering

human-computer interaction





Several interesting  
challenges left for:

maths

computational stats

engineering

human-computer interaction

technology policy





**The Great Wave off  
Kanagawa**

**神奈川沖浪裏**

**Kanagawa oki nami ura**

**Thank you!**